Developments in the NVM Programming Model

Walt Hubis

Vice Chair, Solid States Storage Initiative Hubis Technical Associates January 20, 2015



Need for A New Model





■ NVM Tread ■ NVM xfer ■ Misc SSD ■ Link Xfer ■ Platform + adapter ■ Software

With Next Generation NVM, the NVM is no longer the bottleneck

- Need optimized platform storage interconnect
- Need optimized software storage access methods

Application Access to Non-Volatile Memory

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Disk-like non-volatile memory

- Appear as disk drives to applications
- Accessed using disk stack

Memory-like non-volatile memory

- Appear as memory to applications
- Applications store variables directly in RAM
- No IO or even DMA is required

Memory-like non volatile memory is a type of persistent memory



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SNIA NVM Programming Model



Version 1 approved by SNIA in December 2013

- Downloadable by anyone
- Version 1.1 pending approval

Expose new block and file features to applications

- Atomicity capability and granularity
- Thin provisioning management

Use of memory mapped files for persistent memory

- Existing abstraction that can act as a bridge
- Limits the scope of application re-invention
- Open source implementations available for incremental innovation (e.g. Linux DAX extensions)

Programming Model, not API

- Described in terms of attributes, actions and use cases
- Implementations map actions and attributes to API's

The Four Modes



	Block Mode Innovation	Emerging NVM Technologies
	Atomics	
	Access hints	Performance
	 NVM-oriented 	Performance
	operations	 Perf okay, cost
	Traditional	Persistent Memory
User View	NVM.FILE	NVM.PM.FILE
Kernel Protected	NVM.BLOCK	NVM.PM.VOLUME
Media Type	Disk Drive	Persistent Memory
NVDIMM	Disk-Like	Memory-Like

Conventional Block and File Modes



Discovery of granularities

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Persistent Memory Modes



Use with memory-like NVM

NVM.PM.VOLUME Mode

- Software abstraction to OS components for Persistent Memory (PM) hardware
- List of physical address ranges for each PM volume
- Thin provisioning management

NVM.PM.FILE Mode

- Describes the behavior for applications accessing persistent memory Discovery and use of atomic write features
- Mapping PM files (or subsets of files) to virtual memory addresses
- Syncing portions of PM files to the persistence domain





Building on the Basic PM Model

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Remote access white paper

- Disaggregated memory
- RDMA direct to NVM
- High availability, clustering, capacity expansion use cases
- NVM PM Remote Access for High Availability V02R3 (Draft)

Atomic transactional behavior white paper

- Add atomicity and recovery to programming model
- Not addressed by current sync semantics
- Atomics Transactions Whitepaper V4b (Draft)

Open source contributions

- Linux PMFS at https://github.com/linux-pmfs
- Linux Pmem Examples: https://github.com/pmem/linux-examples



NVM Remote Access

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NVM Programming Remote Access



Use case:

- RDMA copy from local to remote persistent memory
 - > High availability memory mapped files
 - > Built on NVM.PM.FILE from version 1 programming model

Requirements:

- Assurance of remote durability
- Efficient byte range access (e.g., scatter-gather RDMA)
- Efficient addressing
- Efficient write security given fixed addressing in file context
- Resource recovery and hardware fencing after failure
- Security

NVM Memory Access Hardware Taxonomy



Various topologies examined

- Local persistent memory
 - > PM in the same servers as the accessing processor
- Disaggregated persistent memory
 - > PM is not contained within the server
 - Accessed at memory speed
- Network Persistent Memory
 - > PM accessed through a high speed network
- Virtual shared persistent memory
 - Emulating cache coherent shared memory across networked memory using software





High Durability

- Data will not be lost regardless of failures
- May be limited due to implementation of system

High Availability

- Data will remain accessible to hosts regardless of failures
- May be limited due to implementation of system

The distinction between high durability and high availability makes it clear that high availability requires networked access to persistent memory. The network plays an important "fault isolation" role for high availability.



Consistency Points

- All data items recovered must have correct values relative to each other from the application point of view
- Software uses hardware to insure that a failure or restart results in a consistent state

Crash Consistency

- Applications must be prepared to recover from any state of the writes that were in flight when a failure occurs
- Recovery from a crash consistent image is the same as a cold restart after a system crash

Crash consistency is a complex approach to recovery from an application standpoint. It forces considerable overhead to precisely communicate every sync action to networked persistent memory.

NVM Programming Model Approach to Recoverability



Atomicity and Atomic Granularity

Addresses crash consistency issues

Optimized NVM Flush

Addresses consistency point issues

Address Recovery Scenarios

- In-line recovery
- Backtracking recovery
- Local application restart
- Application failover

High Availability RDMA Example

- NVM Programming Model optimized flush
- RAID software for HA
 - local file system
 - remote file system
 - via network file system client and NIC

RDMA for HA

Examine durability, performance, and address space issues.

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Address security issues for NVM RDMA

- Data at rest
- Data in flight
- Authentication
- Authorization
- Threat models
- Transport security
- RDMA security model
 - In the context of NVM RDMA
 - Performance issues with existing models

Atomic Transactions

- Guidelines for atomic capabilities and transactions
- Use byte-addressable persistent memory (PM)
 - Like volatile memory (RAM)
 - Use processor load and store operations
- Retains contents across power loss
 - Like storage
- Provide atomic store operations
 - Assure data consistency
 - Known state in the event of store failure

Relationship with the SNIA NVM Programming Model

Application 1

- Links in a PM-aware library.
- Uses a compiler without PM support.

Application 2

- Uses a PM-aware compiler
- And run-time library.

- Provide for atomic store operations to PM
 - Including the event of system or power failure
- Provide operations for large address ranges
 - Or groups of ranges
- Provide atomic durability to PM
- Provide atomic transactions of arbitrary sizes
- No language extensions required
 - Does not preclude the use of C11 and associated memory model.
- No required hardware extensions
 - Atomicity and transactions rely on existing or near term processor capabilities.

Use Cases

Append to a file atomically

NVM Programming Model Atomics and Transactions

- Form a transaction model for PM
- Review academic and research efforts

Examine current methodologies

- ACID transactions
- Rio Vista
- Failure-Atomic msync
- Memory Mapped Transactions
- Lightweight Recoverable Virtual Memory
- C11 and C11++
- PM aware data structures

Open Source

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Open Source Contributions

Linux PMFS

https://github.com/linux-pmfs

Linux Pmem Examples

- http://pmem.io/nvml/libpmem/
- https://github.com/pmem/linux-examples
- Using NVM in a C Application
 - 4:50 PM 5:20 PM Today
- Processor Support for the NVM Programming Model
 - 4:20 4:50 PM Today
- Linux DAX Extensions
 - Support ext4 on NV-DIMMs
 - http://lwn.net/Articles/588218/

Questions

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