THE REALITIES OF SOLID STATE STORAGE

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Abstract

Leveraging Flash Memory in Enterprise Storage

- This session is for Storage and Data Center Professionals seeking to understand how to best take advantage of flash memory in enterprise storage environments. The relative advantages of flash tiering, caching and all-flash approaches will be considered, across the dimensions of performance, cost, reliability and predictability.
Agenda

- The Storage I/O Crisis
- How to Evaluate Flash Solutions
- Approaches for Leveraging Flash
  - PCI cards
  - Flash caching in arrays
  - Flash tiering in arrays
  - All-flash arrays
- Analyzing Flash ROI
- Under the Hood: How Flash Works
Moore’s Law vs. Newton’s Law

Moore’s Law: 58% CAGR

HDD Areal Density: 40-100% CAGR

HDD Latency (Seek Time) -3% CAGR

Sources: Intel, IDEMA, IBM Research

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Deconstructing a Random Read

1. Spin disk ~2ms
2. Actuate arm ~4ms
3. Read Data ~0.04ms

Typical Resulting Duty Pattern:

Data ———> SEEK ———> Data ———> SEEK ———> Data ———> SEEK ———> Data

Disks spend the majority of their time seeking and rotating, not delivering data!
Randomization: The I/O Blender

Traditional Architecture
- Single-workload
- Serialized
- Cached
- Optimized

Virtualized / Consolidated Architecture
- Multi-workload
- No cross-VM optimization
- Highly randomized
Summary: The Storage I/O Crisis

Performance Demands

- Multi-core CPUs
- Data growth
- Consumerization of IT

Randomization: The “I/O Blender”

- Virtualization
- Data consolidation
- Cloud architectures

IOPS / TB

- 95% time seeking/rotating
- IOPS/TB dropping
Understanding Your Application’s I/O Fingerprint

- **I/O Load**
  - IOPS:
    - 1Ks-10Ks
    - 10Ks-100Ks

- **Block Size**
  - Small Block: <8K
  - Large Block: 10sK - Ms

- **Locality of Access**
  - Sequential, Predictable, Cacheable
  - Random, Unpredictable

- **Access Pattern**
  - Write-Centric
  - Read-Centric

- **Latency Sensitivity**
  - 10s of ms, Variability OK
  - <1-3 ms, Consistency Key
How to Evaluate Flash Alternatives

Cost: $/GB?

$/GB: Managed
+ Operations team

$/GB: Operational
+ Datacenter space
+ Power / Cooling

$/GB: Protected
+ Snapshot / replication software
+ Snapshot / replication copies
+ Backup software and media

$/GB: Usable
+ Cost of RAID parity
+ Over-provisioning waste
+ Management software

$/GB: Raw
+ Cost of raw disk / flash

Performance
• Read IOPS
• Write IOPS
• Latency
• Bandwidth
• $/IOP

Power & Size
• Rackspace / floor space
• Power consumption
• Cooling

Protection
• How to make HA?
• How to backup?
• How to manage drive loss?

Integration
• Fits in current architecture?
• Requires process change?

Operational Simplicity
• Makes operations more or less complex?
• Increases TBs/admin managed?
Flash Architecture Approaches

**Server-Attach PCI Flash**
- PCIe card in application server
- Host CPU typically used for flash management
- 100s of GB

**Array with Flash Cache**
- Flash as a read and/or write cache in array, or to cache FS metadata
- All data persisted to spinning disk
- Typically <1% - 5% of total capacity

**Array with Flash Tier**
- Sub-LUN/FS tiering, where a LUN or FS is spread across flash and disk
- Hot blocks/files moved to flash, cold left on disk
- Typically 1% - 10% of total capacity

**All-Flash LUN or Array**
- 100% flash LUN or array
- May or may not use DRAM for read/write caching
Host-Based PCI Flash

- Overview:
  - Host-based PCI flash cards, typically 100s of GB in size
  - Either non-RAIDed, or two cards mirrored in the server
  - Either non-HA, or application clustered across multiple servers
  - Leverages host CPU for flash management

- High-level Benefits:
  - Highest-performance flash architecture possible
  - Eliminates the cost/burden of shared storage
  - Ultra-small footprint
  - Potentially allows for the reduction of host DRAM

- High-level Challenges:
  - Very high-cost, creates islands of expensive flash in every server
  - Requires a re-architecture for most enterprises
  - Expensive and difficult to protect (HA, backup)
  - Requires mirroring for HA across multiple servers
  - Requires app-level or external replication for DR
Overview:
- Expansion of array’s DRAM cache with flash
- Implemented either via controller-connected PCI flash cards, or SSDs in drive bays
- Typically read-cache only, although some implementations of write cache as well
- Cache page size / caching scheme varies by array (typically 4-16K)

High-level Benefits:
- Expands arrays cache buffer from <<1% of total storage to 1-5%
- 10x latency improvement on IOs which hit the flash cache (10+ms → <1ms)
- Off-loading these IOs from disk reduces the load on disk and also makes the disk perform better
- Result: depends completely on cacheability of I/O stream, but typically 30-80% performance improvement

High-level Challenges:
- Improvement depends completely on cacheability of the I/O stream, results vary
- Flash as a cache heavily exercised: requires SLC flash
- Very expensive: $50-150/GB list prices typical from major vendors
- Cache requires time to “warm” to see the performance benefit, often doesn’t persist across re-boots or HA events
- Minimal benefit for truly random I/O streams: random I/O cant be cached, it is random
Array with Flash Tier

**Overview:**
- Creation of multiple tiers of storage (flash, enterprise FC/SAS HDD, SATA HDD) with the ability to spread a LUN / FS across them
- Typically implemented via 3.5” or 2.5” SSDs in the drive bays/shelves
- Generally 5-10% Flash, 90-95% disk in typical implementations
- Back-end storage is virtualized, and blocks/files are migrated up to flash or down to disk, based upon access patterns
- Chunk size for migration varies by vendor: typically MBs to GBs large
- Frequency varies by vendor: typically daily during low I/O windows

**High-level Benefits:**
- Allows one to leverage flash for IOPs, HDDs for capacity
- Less cost than an all-flash solution, blends the economics of flash and HDD
- Suitable for MLC and SLC implementations of flash
- Potentially allows for significantly less disk / footprint / power if disk has been over-built for performance

**High-level Challenges:**
- Success depends on the predictability/randomness of the I/O stream
- Only works well if the “hot blocks” are consistent
- Current implementation block sizes are large; a single hot block could pin a whole MB or GB chunk in flash
- Not a “real-time” adaptable solution, chunks only promoted/demoted once daily
- Difficult to setup and manage – is it worth the complexity of another tier to manage?
All-Flash LUN or Array

Overview:
- Creation of a LUN on 100% flash
- Typically implemented via 3.5” or 2.5” SSDs in the drive bays/shelves
- Suitable for MLC or SLC flash implementations

High-level Benefits:
- 10x+ potential performance improvement vs. 1-2x for cached/tiered solutions
- Much more consistent latency: no “cache-miss” penalty, all IOs at the speed of flash (typically <1ms)
- Dramatically smaller size, allows for reduction of storage footprint by 4-5x
- Flash fast enough to eliminate the need for significant DRAM caching in the array
- RAID re-build times greatly improved due to the speed of flash (hours → 10s of minutes)

High-level Challenges:
- Cost: flash varies from $20/GB - $150/GB depending on the vendor and type (MLC vs. SLC)
- Connectivity: most flash arrays have limited connectivity options compared to their more mature disk array alternatives
- HA & DR: varies by vendor, but the HA and DR models of these arrays are less mature than existing disk array alternatives
Understanding Flash ROI

**Cost: $/GB?**

- $/GB: Managed
  + Operations team

- $/GB: Operational
  + Datacenter space
  + Power / Cooling

- $/GB: Protected
  + Snapshot / replication software
  + Snapshot / replication copies
  + Backup software and media

- $/GB: Usable
  + Cost of RAID parity
  + Over-provisioning waste
  + Management software

- $/GB: Raw
  + Cost of raw disk / flash

**Flash Impacts...**

- Reduction in management cost (performance troubleshooting)
- 2-5x reduction in space
- 2-5x reduction in power
- 2-3x reduction in disk over-provisioning for performance
- 5-10x the raw cost...

**Source:** averaged / anonymized customer examples

Leveraging Flash Memory in Enterprise Storage
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Summary

Datacenters are facing a storage I/O crisis

Understanding your application's "I/O fingerprint" is key to choosing the best flash strategy for your environment

Flash should be evaluated on several dimensions, it's not all about performance

Look at the full picture (beyond $/GB raw!) to build your flash ROI
Flash Fundamentals

- Flash Cells and How They Work
- Flash Devices and Dynamics
How Flash Cells Work

Electrons Tunnel and Get Trapped in Gate

Floating Gate

Source

Drain
What is Flash Wear?

Electrons tunnel through oxide to charge the floating gate.

Electron traps interfere with charge.
Impact of Geometry Shrink
What is SLC/MLC/TLC?

Smaller and smaller windows to determine signal’s value

Single bit  Multi bit  TLC

Incoming Signals

0 00
0 01
0 10
0 11
1 00
1 01
1 10
1 11
Flash Cells are Arranged Like a Classroom

Cells Share Routing Resources to Reduce Space (and Cost/Power)

They Read/Write/Erase in Different Sizes
  - Erases take longest
Read Write Disturb

- Erase Block
- Read Data Here
- Disturb Data Here
- Write Data Here
- Disturb Data Here
- Write Page
Write Amplification

“You need to move, I am going to erase this block.”

“You need to move, I am going to erase this block.”

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Exercise

- **Reads**
  - Many, fast, predictable
  - Disturbs data at rest

- **Writes**
  - Cannot read and write concurrently
  - Writes disturb data at rest
  - Writes erode flash cells

- **Erases**
  - Long compared to reads and writes
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