

Object Drives: Simplifying the Storage Stack

Mark Carlson/Toshiba

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Abstract



Object Drives: Simplifying the Storage Stack

A number of scale out storage solutions, as part of open source and other projects, are architected to scale out by incrementally adding and removing storage nodes. Example projects include:

- Hadoop's HDFS
- Ceph
- Swift (OpenStack object storage)

The typical storage node architecture includes inexpensive enclosures with IP networking, CPU, Memory and Direct Attached Storage (DAS) block devices. While inexpensive to deploy, these solutions become harder to manage over time. Power and space requirements of Data Centers are difficult to meet with this type of solution. Object Drives further partition these object systems allowing storage to scale up and down by single drive increments.

This talk will discuss the current state and future prospects for object drives. Use cases and requirements will be examined and best practices will be described.

Learning Objectives

- · Definition of object drives
- · Learn the value that they provide
- · Discover where they are they best deployed



- Key/Value semantics (Object store) among others
- Hosted software in some cases
- Interface changed from SCSI based to IP based (TCP/IP, HTTP)
- Channel (FC/SAS/SATA) interconnect moves to Ethernet network

This work is ongoing in the SNIA Object Drive TWG. Please join us at:

https://members.snia.org/apps/org/workgroup/objecttwg/



Key Value example is Kinetic¹

- Metadata is not part of the interface
- Metadata understood by higher levels would be stored as values
- Not expected to be used directly by end user applications (similar to existing Block I/F in that regard)

Higher level example is CDMI

- Includes rich metadata (both user and system)
- Expected to be used by end user applications (and is)

Object Drives, because of their abstraction allow many more degrees of innovation freedom in the implementations behind this abstraction

- Hosts needn't manage the mapping of objects to a block storage device.
- An object drive may manage how it places objects on the media without reference to host specified locations
- 1, Kinetic is a Key Value storage protocol from the Linux Foundation



A number of scale out storage solutions expand by adding identical storage nodes incrementally

Typically use an Ethernet interface and may be connected directly to the Internet

Open source examples include:

- Scale out file systems
 - > Hadoop's HDFS
 - > Lustre
- Ceph
- Swift (OpenStack object storage)

Commercial examples also exist



• System vendors and integrators

- Enables simplification of the software stack
- Hyperscale Data Centers
- Using commodity hardware and open source software
- Enterprise IT
- Following the Hyperscale folks



- For these solutions, typically a commodity server is used as a storage node with direct attached block storage, CPU, memory, networking
- These generalized solutions for specific use cases utilize multiple commodity servers and therefore consume more power and are more complex to manage
- Although less expensive to acquire than previous solutions, they still require higher long term ownership costs



- CPU power is moved to the drive where it can be optimized to the task (I/O) – does not need to be general purpose
- Similar to what happened with (dumb) cell phone processors
- For Smart phones, as desired applications increase their needs, cell phone resources have grown
- CPU/Memory/Network/Storage is one component, managed as such
 - Volume shipments drives the cost of this down
- Enable the resources to be matched/tuned with each other and sized appropriately



Management Complexity

- 10x or more of things to manage
- . Still no management software included with the drives
- SNIA is developing a standard that will allow management of drives at scale

• End to end security

- Data is secured on the drive
- End user applications don't do this securing of the data
 - No secure multi-tenancy in reality
- Higher level software will typically be trusted for authentication and access control
 - Drive has basic message security support (private key)
- End to end integrity
- Ensure that data is not altered during I/O
- e.g. T10 DIF



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Example: Ceph Architecture

Ceph with Object Drives





Traditional Block Storage Device



SATA/SAS/PCIe

ATA/SCSI /NVMe

Disk and/or NVM

Traditional Drive

Interconnect via SATA/SAS/PCIe

- Limited routability
- High development costs
- Typically single host

ATA/SCSI/NVMe Protocols

- Block addressable
- Not designed for lossy network connectivity
- Not typically used in multi-client concurrent access use cases

Kinetic Key Value Drive





Interconnect via Dual Ethernet

- Full routability
- Lower development costs
- Intended for multi-client access
- Path failover for availability

Kinetic Protocol

- Higher-level Key Value interface
- Designed for lossy network connectivity (TCP/IP)
- Typically used in multi-client concurrent access use cases

Pre-Configured In-Storage Compute Drive





Interconnect via Ethernet

Full routability

- Custom Applications installed at factory or provisioning time
 - No pre-determined client-facing interface
 - Example is Ceph, or Kinetic API
 - Custom application can run in a container in an embedded OS environment
 - Custom application accesses storage via standard OS interfaces

Provisionable In-Storage Compute Drive





- Interconnect via Ethernet
 - Full routability
- Custom Applications installable at any time

Interposers





- Allows existing drives to be used as Ethernetconnected Object Drives
- Allows collections of drives to be virtualized as Ethernet-connected Object Drives



• Key Value Protocol (Object Drive)

- Minimal incremental CPU/Memory requirements
- · Simple mapping to underlying storage
- In-Storage Compute (Object Drive)*
- Enough CPU/Memory for Object Node Software to be embedded
 on the drive
- General purpose download or factory installed
- May have additional requirements such as solid state media and more/higher bandwidth networking connections
- In both cases, the interface abstracts the recording technology

*Jim Gray memorial



• Eliminates existing parts of the usual storage stack

- Block drivers, logical volume manager, file system
- And the associated bugs and maintenance costs
- Existing applications need to be re-written, or adapted
- Mainly used by green field developed applications
- Firmware is upgraded as an entire image
- Hyperscale customers are already doing this
 - Using open source software and creating their own "apps" Facebook, Google, etc.
- Key Value organization of data is growing in popularity
 - Examples: Cassandra, NoSQL



- Same advantages as Key Value protocol plus
 - No need for a separate server to run Object Node service (other services still need a server but scale separately)
 - scaling is smoother only adding drives
 - Additional features of the Object Node software can be deployed independently
 - Fewer hardware types that need to be maintained (for selected use cases)
 - Failure domains are more fine grained, thus overall data availability is enhanced



As data on the drive becomes colder, CPU/Memory becomes less utilized

- Possible to host software that then uses this spare resource and works against the cold data
- Extracting metadata

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- Performing preservation tasks
- Other data services: advanced data protection, archiving, retention, deduplication, etc.
- Data Analysis off-load



- Ethernet speeds standardized at 1Gbps, 10 Gbps
- Object Drives are high volume, low margin devices
 - 10 Gbps too expensive for the drive form factor for the foreseeable future
- Desire to have standardized speeds of 2.5 Gbps, 5 Gbps
 - With auto-negotiation among them
- SFF 8601 effort to specify auto-negotiation using existing silicon implementations (switch firmware upgrade – done in a few months)
- 802.3cb will standardize single lane 2.5 and 5 Gbps speeds in silicon for backplanes (multi-year effort)



- The Object Drive TWG is producing a specification for scalable management of Object Drives
- Based on the RedFish management specification from DMTF
 - https://www.dmtf.org/sites/default/files/standards/documents/DSP0266_1.0.1.pdf
- Uses Odata (OASIS) for RESTful interface
- Minimizes and simplifies the management of resources
- New Chassis Type: Object Drive
 - Most other parameters are already part of the base spec
- SNIA standard specifies common features and references the other standards



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Authorship History

Mark Carlson and the Object Drive TWG September 2015

Additional Contributors

David Slik Robert Qiuxin Paul Suhler Dave Landsman Bill Martin

Please send any questions or comments regarding this SNIA Tutorial to <u>tracktutorials@snia.org</u>