# SNIA | DATA, NETWORKING, DNSF | & STORAGE

Ethernet in the Age of Al: Adapting to New Networking Challenges

Live Webinar November 19, 2024 9:00 am PT / 12:00 pm ET



# Today's Presenters



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- Edge implementation opportunities and factors
- Storage and networking security
- Acceleration and offloads
- Programming frameworks
- Sustainability

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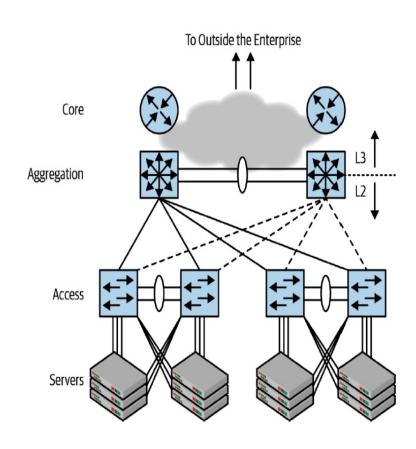
# Today's Agenda

- Overview of Data Center Networks
- LLM GPU Scale and Collective requirements
- Ethernet GPU Fabric Topology
- Ethernet GPU Fabric Requirements
  - Congestion Avoidance
  - Congestion Response



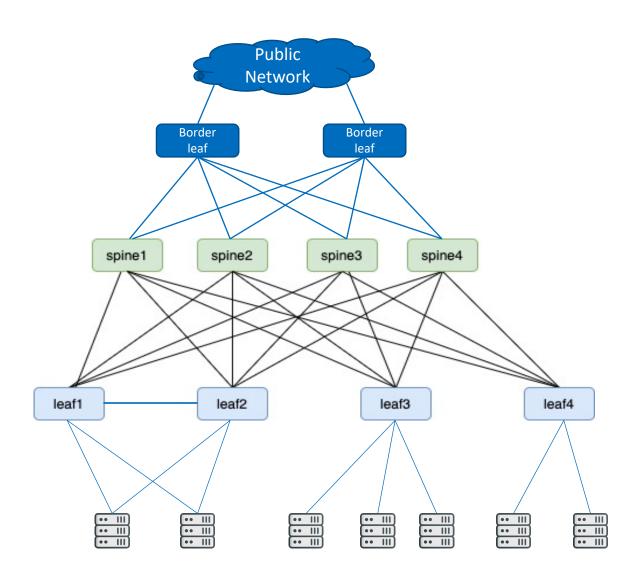
### **Data Center Networks**

- Access, Aggregation and core
- North-South Traffic
- Single Layer-2 domain below aggregation
- Single link failure makes the BW to half
- 2- aggregation switch
- Spanning Tree Protocol
- Vlan cannot span



#### **CLOS Network**

- East-West traffic
- High Bandwidth
- Load Balancing ECMP
- Vxlan
- Multitenancy
- High reliability and fast failure recovery
- Easily scalable



### GPU Scale – The LLM Connection

- Large Language Models (LLMs) are type of generative AI trained on vast natural language data using deep learning algorithms
- Most models stick to open data sets for training
- Tokenization translate the raw words into a sequence of integers (tokens).
- A typical data set can contain hundreds of billions to trillions of tokens
- Weights in LLM are learned variables that dictate how the model interprets and generate languages

## GPU Scale – The LLM Connection

#### Model Math

- For 175B parameters and 300 B tokens and 6 FLOPS per token per parameter ~ 3.15 x10^23
- A GPU with ~67 TeraFLOPS per sec it takes 4.7x10^9 seconds
- To finish the training in 1 month it would take ~1800 GPUs

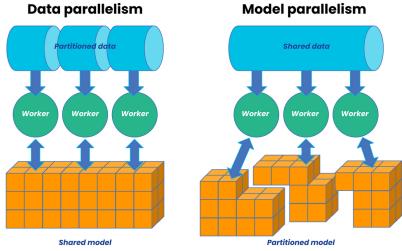


	GPT-3	LLaMA
Sequence length	2k	2k
Parameters	175B	65B
Tokens	300 B	1-1.2 Trillion
Training Time	1 month	21 days

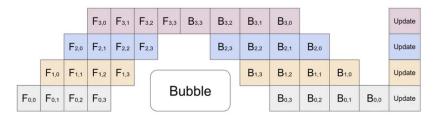


- Model with 175B parameters require greater 1TB of memory
- Storage required checkpoint training state typically around 4TB
- Typical High end GPU has 80 GB of High Bandwidth Memory
- One GPU cannot fit the model parameters or training sets.

### **Parallelism**



(Image source: anyscale.com)



Model + Data Parallelism (GPipe: Huang et al. 2018)

#### Data Parallelism

- Data Set is split into Mini-Batch(s) and each GPU with a copy of the Model make forward pass for predictions and Backward pass for Gradients
- End of one iteration Gradients are aggregated(averaged)
- The aggregated gradients are broadcast to all GPUs.
- Repeat the process till convergence

#### Model Parallelism

- Model is split into several partitions one per GPU.
- During Forward pass each GPU computes the output and pass on as input to the next GPU
- Backward Pass each GPU pass on the gradient to the previous GPU in sequence
- Both Forward/Backward pass create sequential dependency

#### Pipeline Parallelism

- Combines both Data and Model parallelism
- Dataset is further divided into Micro-batches and GPU works on a micro-batch and its model partition
- Instead of waiting for backward pass, it starts to process the next micro-batch
- Increases inter-GPU communication

### **Parallelism**

#### Tensor Parallelism

- Distributes compute intensive parts of the model/pipeline layer by splitting the computation across GPU
- Relies on nodes actively scattering and gathering interim results
- Reduces computation/storage requirement per GPUs for the model training/inference
- Increases inter-GPU computation significantly compared to pipeline parallelism
- Communication strategy would depend on the computation being split. For instance, a matrix multiplication split by row (column) would require All-Reduce (All-Gather) in the forward pass.

- Regardless of which ever parallelism is used it is evident that the inter GPU communication is significant.
- Any Congestion or Drops in the GPU fabric would result in poor performance
- Latency and Link utilization will also play significant role in performance numbers.

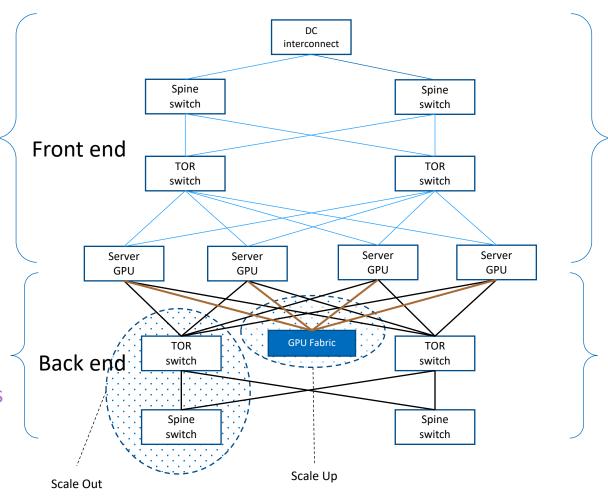
### Collectives

Collectives are set of operations involving communication among group of GPUs to perform a task.

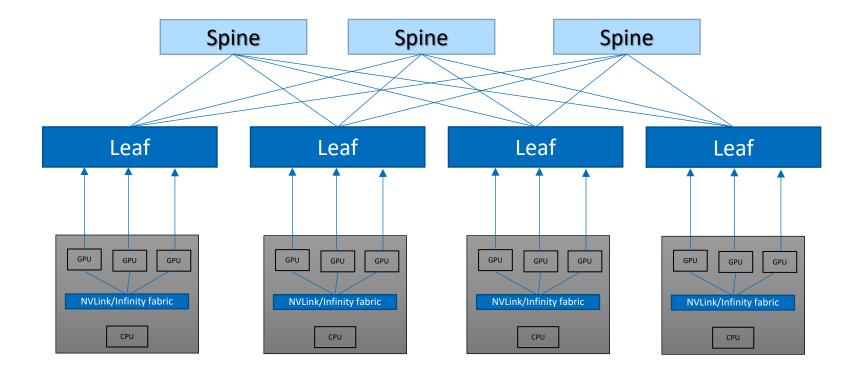
- Reduce
  - Aggregating data from all member and send the result to one member
- All Reduce
  - Aggregating data from all member and send the result to all member
- Scatter
  - Distribute different values from one member to all member.
- Reduce Scatter
  - Aggregate data from all member and scatter the results(unique subset of result) to all member
- Broadcast
  - Sending data from one member to all the member in the group
- All Gather
  - Gather all data and distribute it among all members
- AlltoAll
  - Scatter data from all members to all members

### GPU Scale – GPU Fabric Backend Network

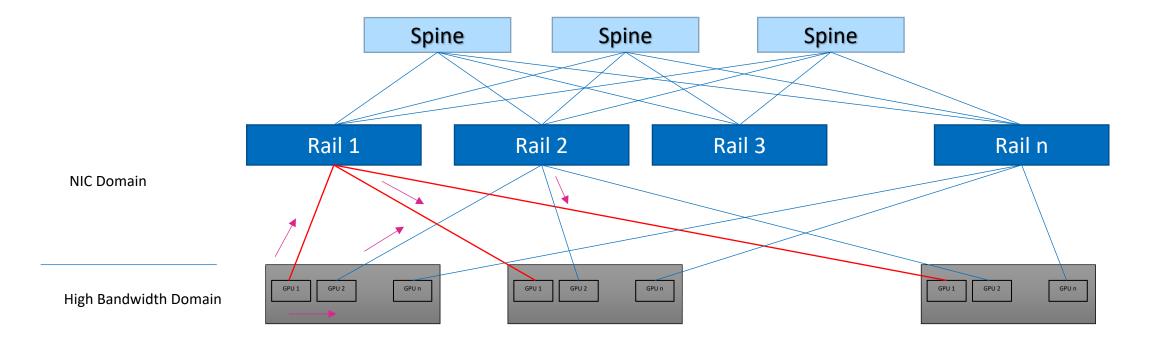
- With the GPU scale required for training there is need for an network to connect thousands of GPU
- There are three network in play here. The primary network is called front end network
- The GPU network is Back end network
  - Scale up: Full mesh NVLink/infinity fabric connections between the GPUs of the same server or same rack
  - Scale out: The network connects thousands of GPU in datacenter across racks
- For Scale Out network Ethernet becomes one of the primary choice
- Because Ethernet is modular, scalable, support high speed, cost effective and works with existing infrastructure
- Ethernet by nature is lossy technology. There are challenges in the area of congestion, latency and utilization etc.,



# Scale up and Scale out GPU Fabric

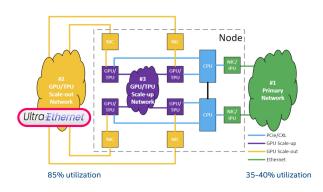


# Rail Optimized Topology



# Al Networking Characteristics and Challenges

#### AI/HPC Networks of Interest: Basic Characteristics

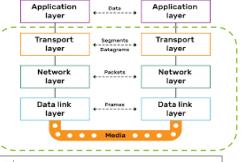


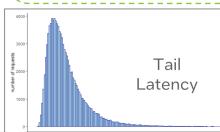
Ultra*Ithernet* 

- . Primary DC network
- Used by all 3 deployment models
- Main network for some HPC At Scale
- · Very large scale: up to 100K-1M Endpoints
- Distance: >150m; RTT ~100 uS +; BW/GPU ~10GB/S
- · Storage attached e.g., over RoCE RDMA
- Network semantics
- 2. GPU/TPU Scale-Out Network
- DL/Inference Cluster -10k nodes and 7
- Distance: <100m; RTT <10 uS +; BW ~100GB/S</li>
- · Main network for some HPC At Scale
- · Network semantics
- 3. GPU/TPU Scale-Up Network
- Within a node; small scale e.g., 256 XPU?
- Distance: ~1m; RTT ~1 uS+; BW ~1200 GB/S7
- · Direct connect and/or switched
- · Memory and Network semantics
- Optimal Link utilization
- Load Balancing Path aware, Adaptive, Lossless
- Loss Retransmission
- HPCC
- Low latency

#### **AI/HPC Common Requirements**







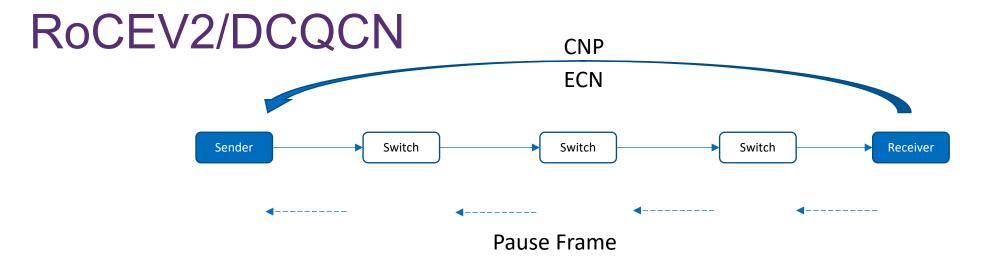
- Transport primitives for
  - Large Scale
  - Multi pathing
  - Relaxed ordering
  - Modernized Congestion Control
  - Optimized RDMA
  - Performance bandwidth, latency, tail latency, Packets/S
  - High network utilization
  - · Stability and Reliability



Key goals: high utilization & low tail latency!



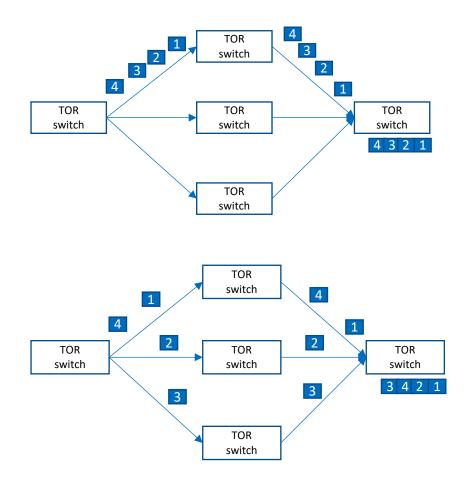




- RoCEV2 RDMA over Converged Ethernet (L3)
- Uses DCQCN as the congestion control mechanism which combines PFC and ECN for congestion management in the RDMA networks
- Enable Lossless ethernet traffic
- PFC is priority based not per flow based. Non congested flows can be affected by a congested flow on the port with same priority
- PFC storm
- ECN is not per flow based.
- Start/stop nature of traffic increases latency

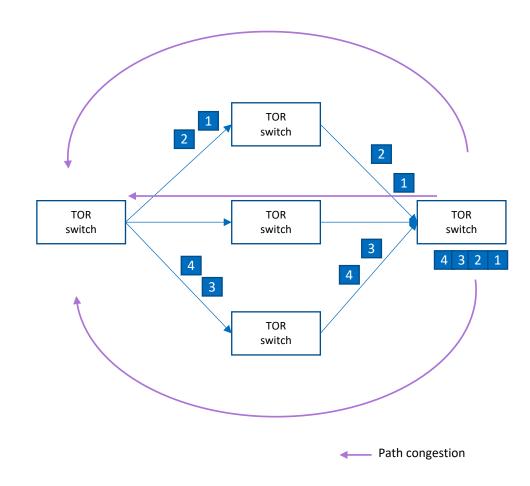
# Link Utilization/Packet Spraying

- In traditional ECMP hashing the packet fields are hashed to pick a path, the path for the flow is constant and ensure inorder delivery.
- Under utilizes the paths available to the destination.
- An elephant flow can create congestion on the path, even other paths are available
- Packet spraying is not new to the switches. Sprays the packets from the same flows to available paths
- Receiver will get out of order delivery based on the path length and congestion status.
- Receiver should be capable of arrange the packets in the right order.
- Utilizes the paths in a balanced way better than before
- Spraying can happen from the NICs or the leaf switches
- Need sophisticated hardware for out of order reassembly
- Need careful configuration as it might affect other schemes?



# Adaptive Load Balancing/ Path aware CC

- Real time telemetry about the congested paths
- Dynamically avoid congested paths
- Traditional load balance uses only that nodes queue depths whereas in this case the entire path congestion status is used for load balancing
- During transient congestion the same flow can be load balanced to different path. Out of order delivery to be handled.
- Careful tuning is required as it might cause unwanted out of order packets

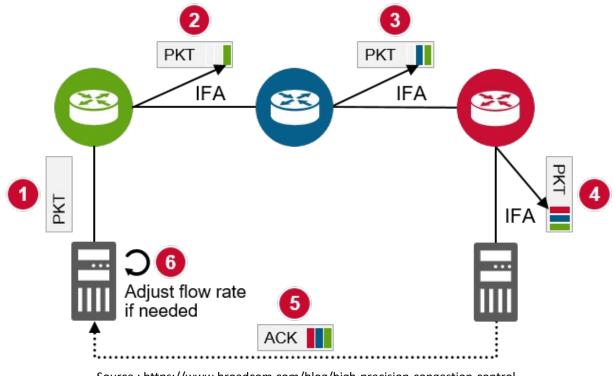


### INT/CSIG - HPCC

- In band Telemetry (INT)/Congestion signal(CSIG) carry fine-grained network signals for congestion control and traffic management to the end host
- In INT every switch along the path add meta data about the congestion on the node to the packet
- CSIG provides a simple, low-overhead, and extensible packet header mechanism to obtain fixed-length summaries from bottleneck devices along a packet path
- The end host can reflect this congestion information to the source real time to adjust the rate and the window size per flow
- Can help to ramp up higher speeds at faster rate
- IFA/CSIG capable hardware is needed for the entire path. Cost would be an issue

#### Reference:

https://www.ietf.org/archive/id/draft-ravi-ippm-csig-01.html



Source: https://www.broadcom.com/blog/high-precision-congestion-control

# **Congestion Control**

#### Congestion Control (CC)

...is required, (in addition to optimal load balancing on multiple paths)

- How is UET CC different from TCP?
  - High bandwidth, short RTT (10us)
  - · Get to wire rate very quickly
    - 1MB takes 10 usec at 800gbps = 1 RTT
    - · No time to wait for TCP slow start
  - · And back off quickly when congestion is noticed

**UEC Congest Control is Designed for short RTT, high BW** 

40

%30 20

#### **UET Congestion Control**

Two flavors - that can work together

- Sender-based (default)
  - fast ramp, fast slowdown
  - uses delay as a measure of congestion
- Receiver-based (optional)
  - · receiver credit manages receiver buffer
  - simple rate-control for oversubscribed networks

sender control



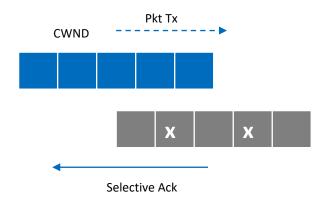
receiver control



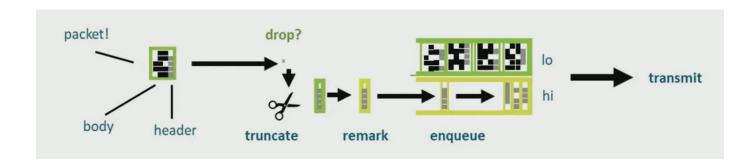
Both are designed for AI and HPC workloads and multi-pathing

### Loss Retransmission

- Enhanced network performance through inorder message delivery and selective acknowledgment (SACK) retransmission.
- Unlike RoCEv2's Go-back-N mechanism, which resends all packets from the point of failure, SACK allows the receiver to identify and retransmit only lost or corrupted packets.
- This targeted approach optimizes bandwidth utilization, reduces latency in packet loss recovery and minimizes redundant data transmission
- Faster job completion times, lower tail latencies, and more efficient bandwidth use



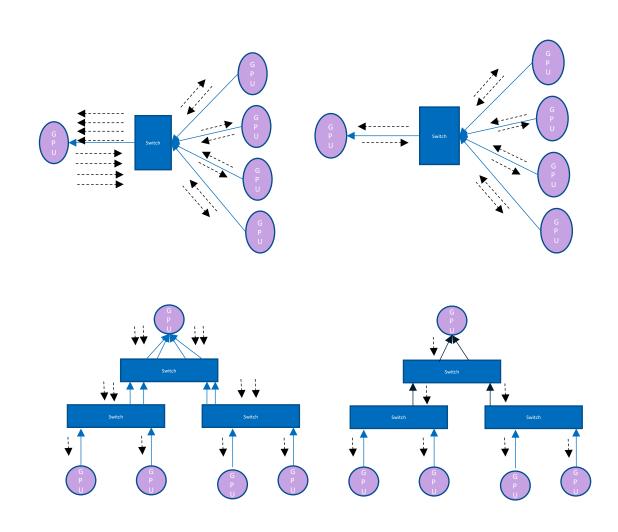
# **Packet Trimming**



- Fast loss detection
- Trim the packet to 64 bytes instead of dropping
- Mark the DSCP to trimmed for the receiver to identify
- Send via the high priority queue
- This would be useful for the receiver to identify the packet loss faster and send SACK to the sender for retransmission

# INC – In Network Collectives

- It is a method to offload collective operations to the switches instead of GPUs
- Offloading collectives to network devices reduces traffic bottlenecks and enhance the performance
- Implemented in form of aggregate tree where each node in the tree aggregate the downstream flows and forward the data to the upstream node
- By design avoids in cast congestion during reduce collective operations.



Requirement	UEC Transport	Legacy RDMA	UEC Advantage
Multi-Pathing	Packet spraying	Flow-level multi-pathing	Higher network utilization
Flexible Ordering	Out-of-order packet delivery with in-order message delivery	N/A	Matches application requirements, lower tail latency
Al and HPC Congestion Control	Workload-optimized, configuration free, lower latency, programmable	DCQCN: configuration required, brittle, signaling requires additional round trip	Incast reduction, faster response, future-proofing
In Network Collective	Built-In	NONE	Faster Collective operation, lower latency
Simplified RDMA	Streamlined API, native workload interaction, minimal endpoint state	Based on IBTA Verbs	App-level performance, lower cost implementation
Security	Scalable, 1st class citizen	Not addressed, external to spec	High scale, modern security
Large Scale with Stability and Reliability	Targeting 1M endpoints	Typically, a few thousand simultaneous end points	Current and future-proof scale

# Summary

- The size of number of tokens in data set and model parameters in the LLM training require high number of GPUs
- GPU Fabric to scale up to some extent beyond that Scale out network is needed
- Ethernet is one of the primary options for Scale out network because of its proven advantages
- CLOS and Rail-optimized Scale out topology
- The traffic pattern generated because of the AI workloads demand new methods to
  - utilize bandwidth better
  - reduce latency and
  - congestion control
- UEC is working on a UE Transport specification to address the AI workload demands

### References

- Large Language Models The HW connection
- GPU Fabrics for Gen Al Workload
- Rail Optimized Topology Meta Paper
- SwitchAgg: A Further Step Towards In-Network Computation

#### **UEC**

- UEC Introduction
- Networking for Al and HPC, and Ultra Ethernet

# Q&A

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