NFSv4.1 — Using pNFS

SNIA WEBCAST

Presented by: Alex McDonald CTO Office, NetApp



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Alex McDonald Office of the CTO NetApp

Alex McDonald joined NetApp in 2005, after more than 30 years in a variety of roles with some of the best known names in the software industry .

With a background in software development, support, sales and a period as an independent consultant, Alex is now part of NetApp's Office of the CTO that supports industry activities and promotes technology & standards based solutions, and is co-chair of the SNIA Cloud Storage Initiative, and co-chair of the SNIA File Protocols Special Interest Group.







Ethernet Storage Forum Members





- File Protocol SIG drives adoption and understanding of SMB and NFS across vendors to constituents
 - Marketing, industry adoption, Open Source updates
- NetApp, EMC, Panasas and Sun founders
- White papers on migration from NFSv3 to NFSv4
 - An Overview of NFSv4; NFSv4, NFSv4.1, pNFS, and proposed NFSv4.2 features





Previous SNIA NFS Presentations

Education

BrightTalk SNIA Channel NFS Mini Series

NFSv4.1, pNFS & FedFS Protocol Development

- Part1 Four Reasons for NFSv4
 - Discusses the reasons behind the development of NFSv4 and beyond, and the need for a better-than-NFSv3 protocol

Part2 – Advances in NFS – NFSv4.1 and pNFS

- An overview and some details on NFSv4.1, pNFS (parallel NFS), and FedFS (the Federated filesystem); and a high level overview of proposed NFSv4.2 features
- Part3 Planning for a Smooth Migration
 - The key issues to consider when migrating from NFSv3 or implementing new applications with NFSv4.1; Unicode, security with Kerberos, statefulness, selecting the application and other aspects.
- Slides available from
 - <u>http://snia.org/forums/esf/knowledge/webcasts</u>





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The Four Reasons for NFSv4.1

	Functional	Business Benefit
Security	ACLs for authorization Kerberos for authentication	Compliance, improved access, storage efficiency, WAN use
High availability	Client and server lease management with fail over	High Availability, Operations simplicity, cost containment
Single namespace	Pseudo directory system, FedFS	Reduction in administration & management
Performance	management with fail oversPseudo directory system, FedFSRMultiple read, write, delete operations per RPC callBDelegate locks, read and writeL	Better network utilization for all NFS clients
	-	Leverage NFS client hardware for better I/O
	Parallelised I/O	
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Education

We'll cover

- Overview of pNFS terminology and operation
- How clients & servers co-operate to provide parallelism while supporting data consistency
- Some implementation considerations
- This is a high level overview
 - But more technical content for background
 - Use SNIA white papers and vendors (both client & server) to help you implement





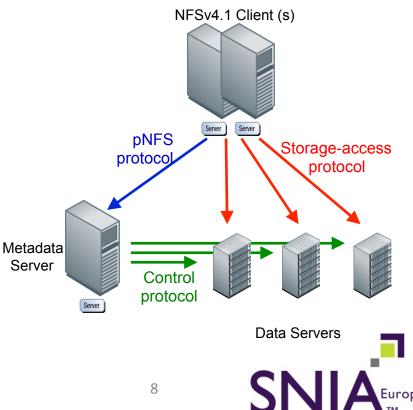
You've Done NFSv4.1; now for pNFS

Education

NFSv4.1 (pNFS) can aggregate bandwidth

- Modern approach; relieves issues associated with point-to-point connections
- pNFS Client
 - Client read/write a file
 - Server grants permission
 - File layout (stripe map) is given to the client
 - Client parallel R/W directly to data servers

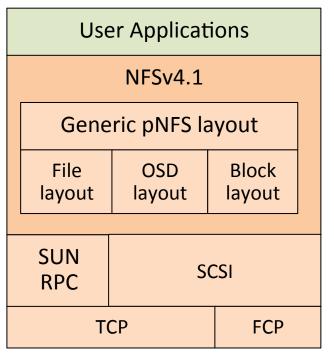
- Removes IO
 Bottlenecks
 - No single storage node is a bottleneck
 - Improves large file performance
- Improves
 Management
 - Data and clients are load balanced
 - Single Namespace



Relationship of pNFS to NFSv4.1

 RFC 3530bis – Network File System (NFS) Version 4 Protocol

- NFSv4 (updated from RFC 3530 based on experience)
- RFC 5661 Network File System (NFS) Version 4 Minor Version 1 Protocol
 - Specifies Sessions, Directory Delegations, and parallel NFS (pNFS) for files
- RFC 5663 Parallel NFS (pNFS) Block/ Volume Layout
- RFC 5664 Object-Based Parallel NFS (pNFS) Operations
- PNFS is dependent on session support, which is only available in NFSv4.1





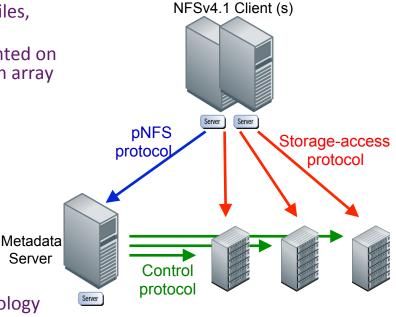


pNFS Terminology

Education

Metadata Server; the MDS

- Maintains information about location and layout of files, objects or block data on data servers
- Shown as a separate entity, but commonly implemented on one or across more than one data server as part of an array
- pNFS protocol
 - Extended protocol over NFSv4.1
 - Client to MDS communication
- Storage access protocol
 - Files; NFS operations
 - Objects: OSD SCSI objects protocol (OSD2)
 - Blocks; SCSI blocks (iSCSI, FCP)
- Control protocol
 - Not standardised; each vendor uses their own technology to do this
- Layout
 - Description of devices and sector maps for the data stored on the data servers
 - 3 types; files, block and object
- Callback
 - Asynchronous RPC calls used to control the behavior of the client during pNFS operations



Data Servers



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- Client requests layout from MDS
- Layout maps the file/object/ block to data server addresses and locations
- Client uses layout to perform direct I/O to the storage layer
- MDS or data server can recall the layout at any time using callbacks
- Client commits changes and releases the layout when complete
- pNFS is optional

education

standards

certification

oest practice

Client can fall back to NFSv4

- pNFS operations
 - LAYOUTCOMMIT Servers commit the layout and update the meta-data maps
 - LAYOUTRETURN Returns the layout or the new layout, if the data is modified
 - GETDEVICEINFO Client gets updated information on a data server in the storage cluster
 - GETDEVICELIST Clients requests the list of all data servers participating in the storage cluster
 - CB_LAYOUT Server recalls the data layout from a client if conflicts are detected





pNFS Pre-requisites

- NFSv4.1 and pNFS capable server
 - Contact your NAS vendor for availability
 - Commercial products available for all of files, blocks and object types
 - Open source Linux pNFS server in development
 - http://wiki.linux-nfs.org/wiki/index.php/ PNFS_Development
- pNFS capable client
 - Linux to date
 - See previous BrightTalks
 - Part3 Planning for a Smooth Migration





Linux Client and NFSv4.1

- Upstream (Linus) Linux NFSv4.1 client support
 - Basic client in Kernel 2.6.32
 - pNFS support (files layout type) in Kernel 2.6.39
 - Support for the 'objects' and 'blocks' layouts was merged in Kernel 3.0 and 3.1 respectively
- Full read and write support for all three layout types in the upstream kernel
 - Blocks, files and objects
 - O_DIRECT reads and writes supported





Linux Client and NFSv4.1

- pNFS client support in distributions
 - Fedora 15 was first for pNFS files
 - Kernel 2.6.40 (released August 2011)
- Red Hat Enterprise Linux (RHEL)
 - "Technical preview" support for NFSv4.1 and for the pNFS files layout type in version 6.2, 6.3
 - Full support in RHEL6.4, announced Feb 2013
- Ubuntu, SUSE & other distributions
 - Possible to upgrade to NFSv4.1
- No support in Solaris
 - Both server and client are NFSv4 only





pNFS Files Mount

Education

RHEL6.4 pNFS mount

- mount -o minorversion=1 server:/filesystem /mnt
- Check
 - (output edited)

/proc/self/mountstats

```
device 172.16.92.172:/filesystem mounted on /mnt with fstype
nfs4 statvers=1.1
opts: ...,vers=4.1, ...
nfsv4: ...,sessions,pnfs=nfs_layout_nfsv41_files
...
```





pNFS Client Mount

47 27.086618 172.17.40.185	172.17.40.171	NFS	282 V4 Call (Reply In 48) EXCHANGE_ID		
48 27.086762 172.17.40.171	172.17.40.185	NES	266 V4 Reply (Call In 47) EXCHANGE_ID		
49 27.086883 172.17.40.185	172.17.40.171	NFS	330 v4 call (Reply In 51) CREATE_SESSION		
50 27.087003 172.17.40.171	172.17.40.185	NFS	146 V1 CB_NULL Call (Reply In 53)		
51 27.087032 172.17.40.171	172.17.40.185	NFS	194 V4 Replv (Call In 49) CREATE SESSION		
Ethernet II, Src: Netapp_20:7a:42	(00:a0:98:20:7a:42)	Dst: Int	elcor_2b:40:06 (00:1b:21:2b:40:06)		
			, Dst: 172.17.40.185 (172.17.40.185)		
			1007 (1007), Seq: 29, Ack: 261, Len: 200		
■ Remote Procedure Call, Type:Reply					
∃ Network File System, Ops(1): EXCH/	ANGE_ID				
[Program Version: 4]		1-	72.17.40.185 – IP address of the pNFS client		
Status: NF54_OK (0) 172.17		72.17.40.171 – IP address of the server			
⊞ Tag: <empty></empty>					
<pre> Operations (count: 1) </pre>		Client and Server handshake to determine respective			
Opcode: EXCHANGE_ID (42)		Capabilities. The Cluster replies with MDS and DS flags			
Status: NF54_OK (0)	-				
clientid: 0x638722000000004	1	set, indicating capability for both			
seqid: 0x00000001					
□ eir_flags.0x00060100					
0					
.0 = EXCHGID4_FLAG_UPD_CONFIRMED_REC_A: Not set					
$\dots \dots = EXCHGID4_FLAG_USE_PNFS_DS: Set$					
<pre> = EXCHGID4_FLAG_BIND_PRINC_STATEID: Set</pre>					
		HGID4_FLAG	S_SUPP_MOVED_REFER: Not set		
eia_state_protect: SP4_NONE	(0)				

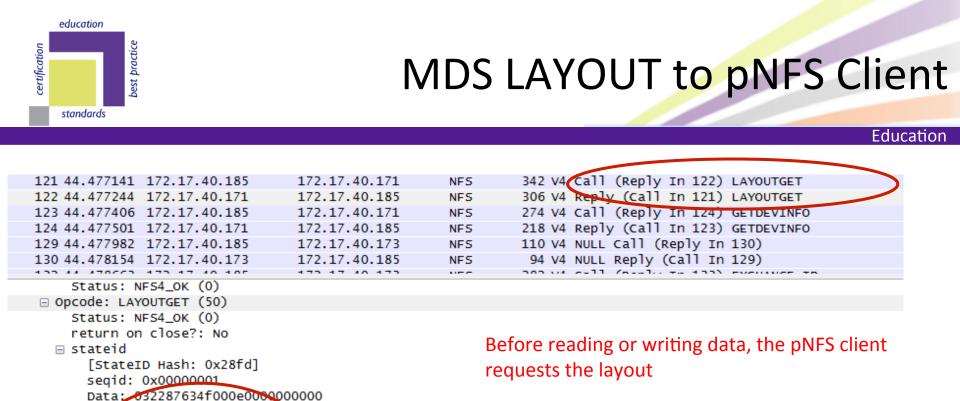




pNFS Client to MDS

117 44.370851 172.	.17.40.185	172.17.40.171	NES	418 v4 call (Reply In 118) OPEN DH:0x7f69f7d7/testfile5
118 44.470682 172.	.17.40.171	172.17.40.185	NFS	566 V4 Reply (Call In 117) OPEN StateID:0xa36e
119 44.470856 172.	.17.40.185	172.17.40.171	NFS	338 v4 call (Reply In 120) SETATTR FH:0x4c99adea
120 44.471391 172.	.17.40.171	172.17.40.185	NFS	318 V4 Reply (Call In 119) SETATTR
121 44.477141 172.	.17.40.185	172.17.40.171	NFS	342 V4 Call (Reply In 122) LAYOUTGET
122 44.477244 172.	.17.40.171	172.17.40.185	NFS	306 V4 Reply (Call In 121) LAYOUTGET
123 44.477406 172.	.17.40.185	172.17.40.171	NFS	274 V4 Call (Reply In 124) GETDEVINFO
124 44.477501 172.			NFS	218 V4 Reply (Call In 123) GETDEVINFO
129 44.477982 172.			NFS	110 V4 NULL Call (Reply In 130)
130 44.478154 172.	.17.40.173	172.17.40.185	NFS	94 V4 NULL Reply (Call In 129)
device index: □ r_netid: tcp length: 3 contents: t fill bytes: □ r_addr: 172.1	: 0 5lot id: 15 INFO (47) _OK (0) LAYOUT4_NFSV4_1_ : 0 ССР	_FILE5 (1)		The OPEN and SETATTR are sent to the MDS
length: 17 contents: 172.17.40.173.8.1 fill bytes: opaque data [Main Opcode: GETDEVINFO (47)]				





The map of data servers and file handles is returned



□ Layout Segment (count: 1)

IO mode: IOMODE_RW (2)

nfl util: 0x00010000

⊮ File Handles (count: 1) [Main Opcode: LAYOUTGET (50)]

length: 18446744073709551615

Nirst stripe to use index:

offset: 0

offset: 0



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pNFS Client DEVICEINFO from MDS

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117 44.370851 172.17.40.185	172.17.40.171	NES	418 v4 Call (Reply In 118) OPEN DH:0x7f69f7d7/testfile5
118 44.470682 172.17.40.171	172.17.40.185	NES	566 V4 Reply (Call In 117) OPEN StateID:0xa36e
119 44.470856 172.17.40.185	172.17.40.171	NFS	338 V4 Call (Reply In 120) SETATTR FH:0x4c99adea
120 44.471391 172.17.40.171	172.17.40.185	NFS	318 V4 Reply (Call In 119) SETATTR
121 44.477141 172.17.40.185	172.17.40.171	NFS	342 V4 Call (Reply In 122) LAYOUTGET
122 44.477244 172.17.40.171	172.17.40.185	NFS	306 V4 Reply (Call In 121) EAVOUIGET
123 44.477406 172.17.40.185	172.17.40.171	NFS	274 v4 Call (Reply In 124) GETDEVINFO
124 44.477501 172.17.40.171	172.17.40.185	NFS	218 V4 Reply (Call In 123) GETDEVINFO
129 44.477982 172.17.40.185	172.17.40.173	NFS	110 V4 NULL Call (Reply In 130)
130 44.478154 172.17.40.173	172.17.40.185	NFS	94 V4 NULL Reply (Call In 129)
<pre>seqid: 0x00000017 slot ID: 0 high slot id: 0 target high slot id: 15 status: 0 Opcode: GETDEVINFO (47) Status: NF54_0K (0) layout type: LAYOUT4_NF5V4 device index: 0 Ir_netid: tcp length: 3 contents: tcp fill bytes: opaque data r_addr: 172.17.40.173.8.1 length: 17</pre>	_1_FILE5 (1)		Meta-data node provides the pNFS client with the IP information for the DS. In this example – 172.17.40.173
contents: 172.17.40.173.4 fill bytes: opaque data [Main Opcode: GETDEVINFO (47))		ormation is cached for life of the layout or til recalled (for example, when the data is moved)



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pNFS Client Uses Direct Data Path

Education

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123 44.477406 172.17.40.185	172.17.40.171	NFS	274 V4 Call (Reply In 124) GETDEVINFO		
124 44.477501 172.17.40.171	1/2.1/.40.185	NFS	218 V4 Reply (Call In 123) GETDEVINFO		
129 44.477982 172.17.40.185	172.17.40.173	NFS	110 V4 NULL Call (Reply In 130)		
130 44.478154 172.17.40.173	172.17.40.185	NFS	94 V4 NULL Reply (Call In 129)		
132 44.478663 172.17.40.185	172.17.40.173	NFS	282 V4 Call (Reply In 133) EXCHANGE_ID		
133 44.478784 172 17.40.173	172.17.40.185	NFS	266 V4 Reply (Call In 132) EXCHANGE_ID		
134 44.478918 172.17 40.185	172.17.40.173	NFS	330 v4 call CREATE_SESSION		
163 60.480000 172.17.40.185	172.17.40 173	NFS	330 v4 call (Reply In 206) CREATE_SESSION		
169 64.476795 172.17.40.185	172.17.40.171	NFS	242 V4 Call (Reply In 170) SEQUENCE		
170 64.476916 172.17.40.171	172.17.40.185	NFS	150 v4 Reply (Call In 169) SEQUENCE		
191 76.480717 172.17.40.185	172.17.40.173	NFS	330 v4 call CREATE SESSION		
Network File System, Ops(2): SEQUE	NCE GETDEVINFO				
[Program Version: 4]					
[V4 Procedure: COMPOUND (1)]					
Status: NF54_OK (0)					
<pre>Operations (count: 2)</pre>					
□ Opcode: GETDEVINFO (47)					
Status: NF54_OK (0)					
layout type: LAYOUT4_NFSV4_1_FILES (1)					
device index: 0	/ _ /				
□ r_netid: tcp					
length: 3					
contents: tcp			ow the pNFS client is reaching out to the		
			mate velumes on a direct noth using ID		
□/r_addr: 172.17.40.173.8.1			mote volume on a direct path using IP		
			drage 172 17 40 172		
		30			
-	1	au	ldress 172.17.40.173.		
sontents: 172.17.40.173.8.	1	au	Iuress 172.17.40.173.		
-		au	Iuress 172.17.40.173.		



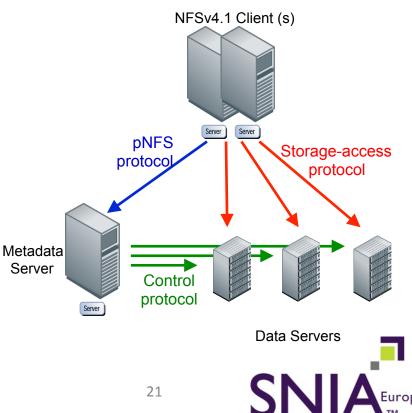


In Summary: The Benefits of pNFS

NFSv4.1 (pNFS) can aggregate bandwidth

- Modern approach; relieves issues associated with point-to-point connections
- pNFS Client
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 Management
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Other NFS Performance Capabilities

- Trunking (NFSv4.1 & pNFS)
 - A single data server connection limits data throughput based on protocol
 - Trunking "bundles" connections into a single pipe
 - Open multiple sessions via different physical Ethernet connections to the same file handle/data server resource
 - Expands throughput and can reduce latency
 - No implementations as yet
- Compound operations (NFSv4 and above)
 - Example: LOOKUP, OPEN, READ, CLOSE as a single RPC call
 - Benefits WAN operations
- Caching & delegation (NFSv4 and NFSv4.1)
 - Allows client and server to agree on data that will be processed by the client
 - Reduces IO and provides data locality





Summary/Call to Action

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- Start using NFSv4.1 today
 - NFSv4.2 nearing approval
 - pNFS offers performance support for modern NAS devices
- Planning is key
 - Application, issues & actions to ensure smooth implementations
 - PNFS
 - First open standard for parallel I/O across the network
 - Ask vendors to include NFSv4.1 and pNFS support for client/ servers
 - pNFS has wide industry support
 - Commercial implementations and open source







Question & Answer



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