

Security of Data on NVMe over Fabrics, The Armored Truck Way

Live Webcast

May 12, 2021

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Today's Presenters



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Ethernet, Fibre Channel, InfiniBand®

iSCSI, NVMe-oF™, NFS, SMB

Virtualized, HCI, Software-defined Storage

Storage Protocols (block, file, object)

Securing Data

Technologies We Cover

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Storage Security Overview

Nishant Lodha

What do customers expect from NVMe over Fabrics?



Cost



Operational simplicity



Performance



Applications



Scale



Security



Data Protection vs. Security

There is often confusion between storage/data security and protection.

Data security refers specifically to measures taken to protect the integrity of the data itself.

Data security primarily involves keeping private information out of the hands of anyone not authorized to see or modify it.

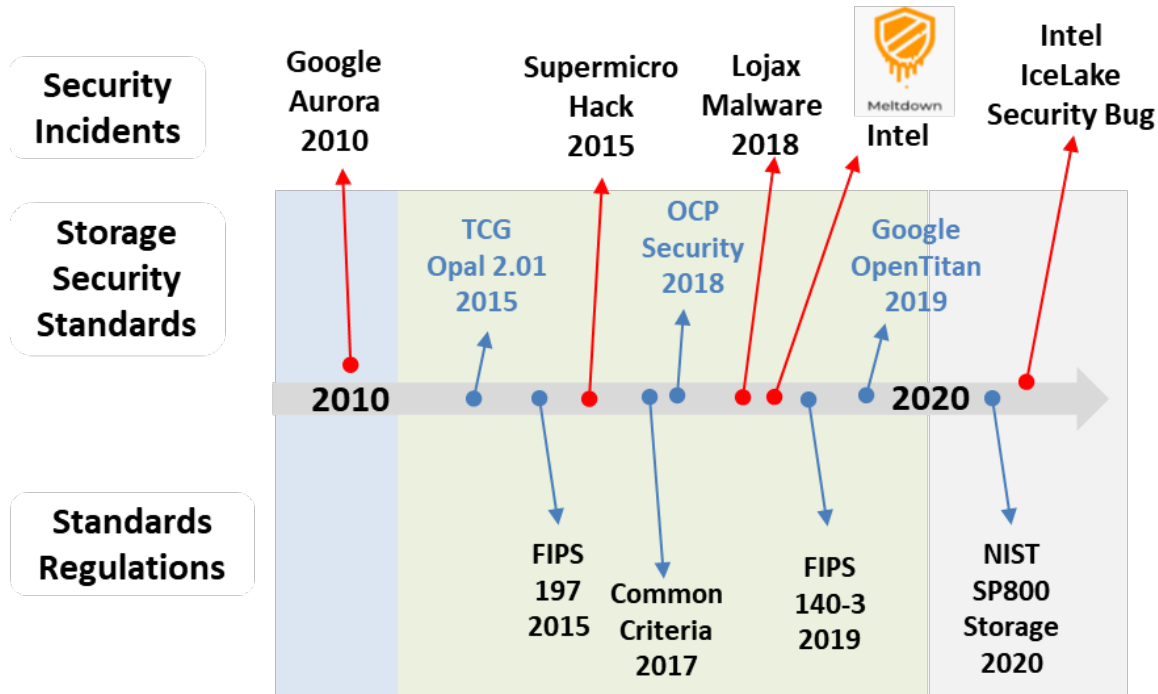
Unauthorized access and access control, auditing
Intentional or accidental loss/corruption of sensitive data

Data Security measures include encryption of data both at rest and in motion, physical and logical access controls that prevent unauthorized access etc.

Data Protection, refers to the mechanism of making copies of your data to restore in the event of a loss or corruption.

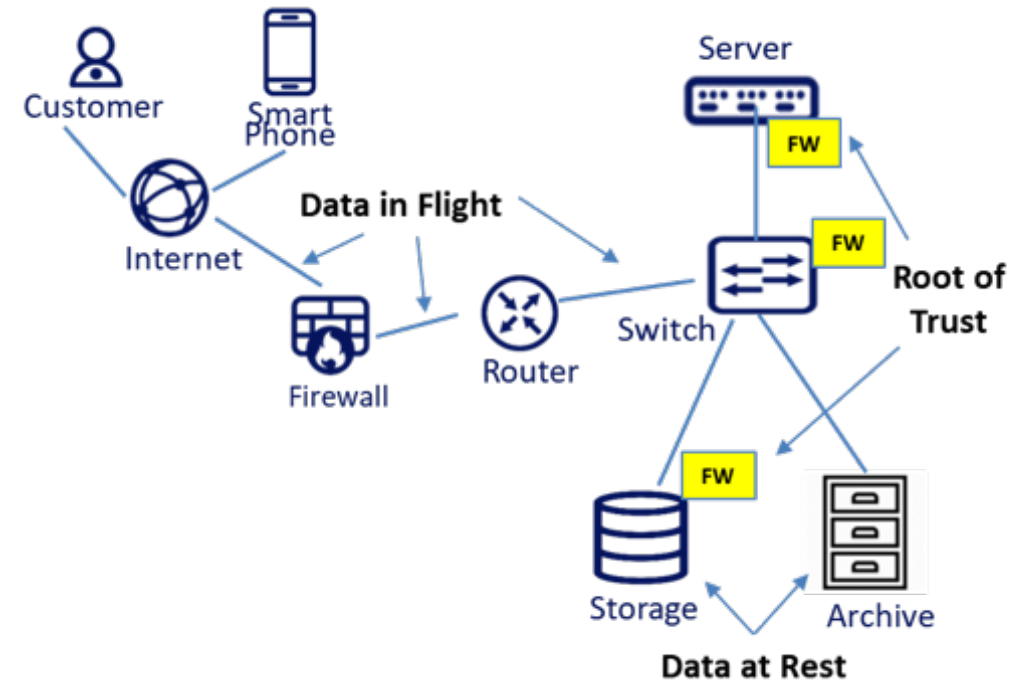
Datacenter Security and Standards

Incidents and Security Standards



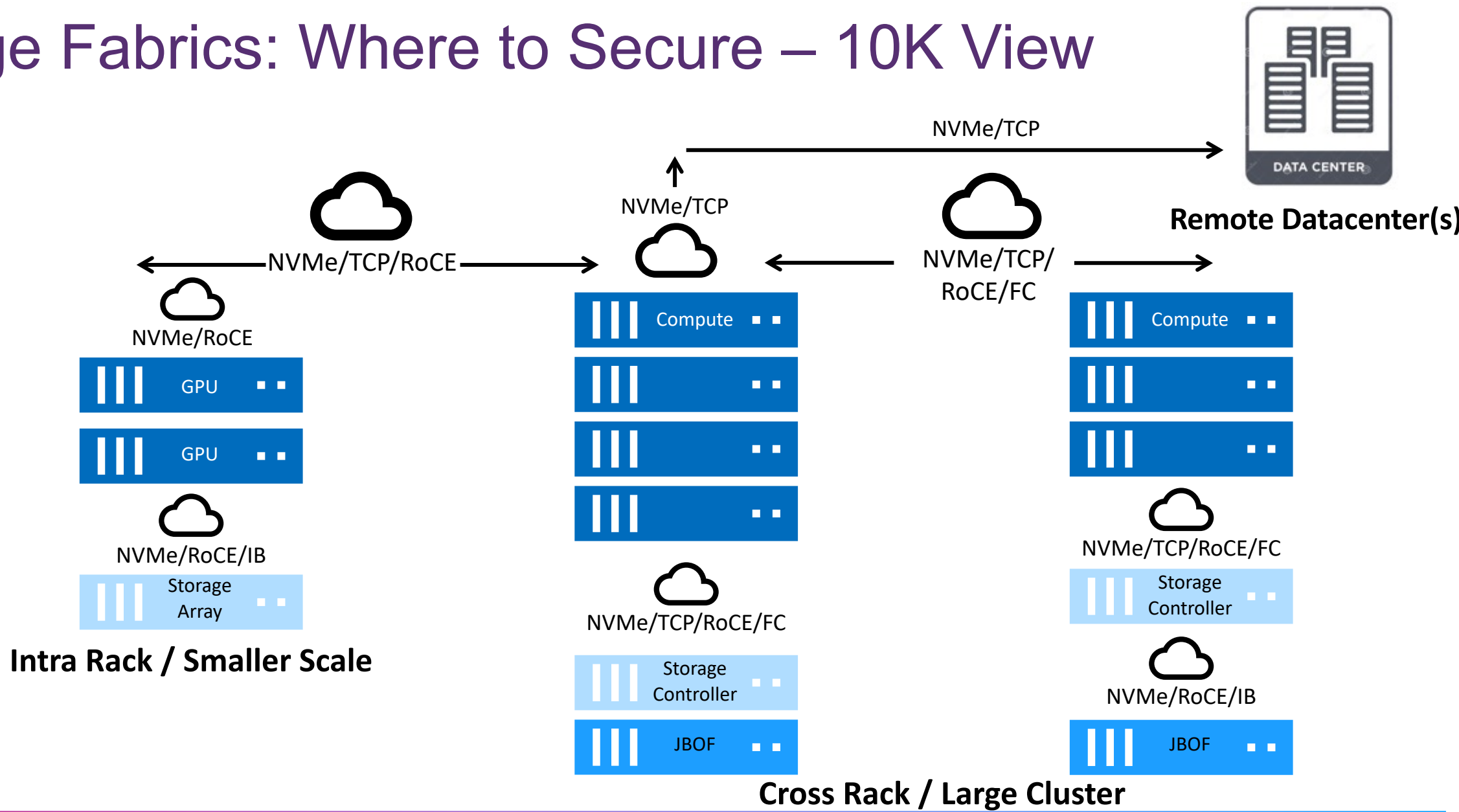
- Standards, Security threats growing in past 10 yrs.
- New Security Standards organizations emerged

Datacenter Security Considerations



- **Data in Flight:** Network security (especially applicable for shared infrastructure)
- **Data at Rest:** Against theft of data or keys, and ransomware (esp. SSD media and key encryption with SED)
- **HW Root of Trust :** Dedicated security engine to ensure Secure Boot, Secure FW, and Key Management across all peripherals

Storage Fabrics: Where to Secure – 10K View



Drivers for Storage Security

Implementing a “ZERO Trust” framework requires advanced technologies

Healthcare
Finance
Defense
Government

Sensitive Verticals

Multi-tenancy
DR/Cloud Storage
Malicious Insiders

New Deployment
Use Cases

HIPAA
GDPR
ISO270001

Regulatory

Potential DC NVMe-oF Security Threats

**Sniffing
Storage Traffic**



**Storage
Masquerading**



Ransomware



Session Hijacking



Must secure NVMe payloads in flight and rest



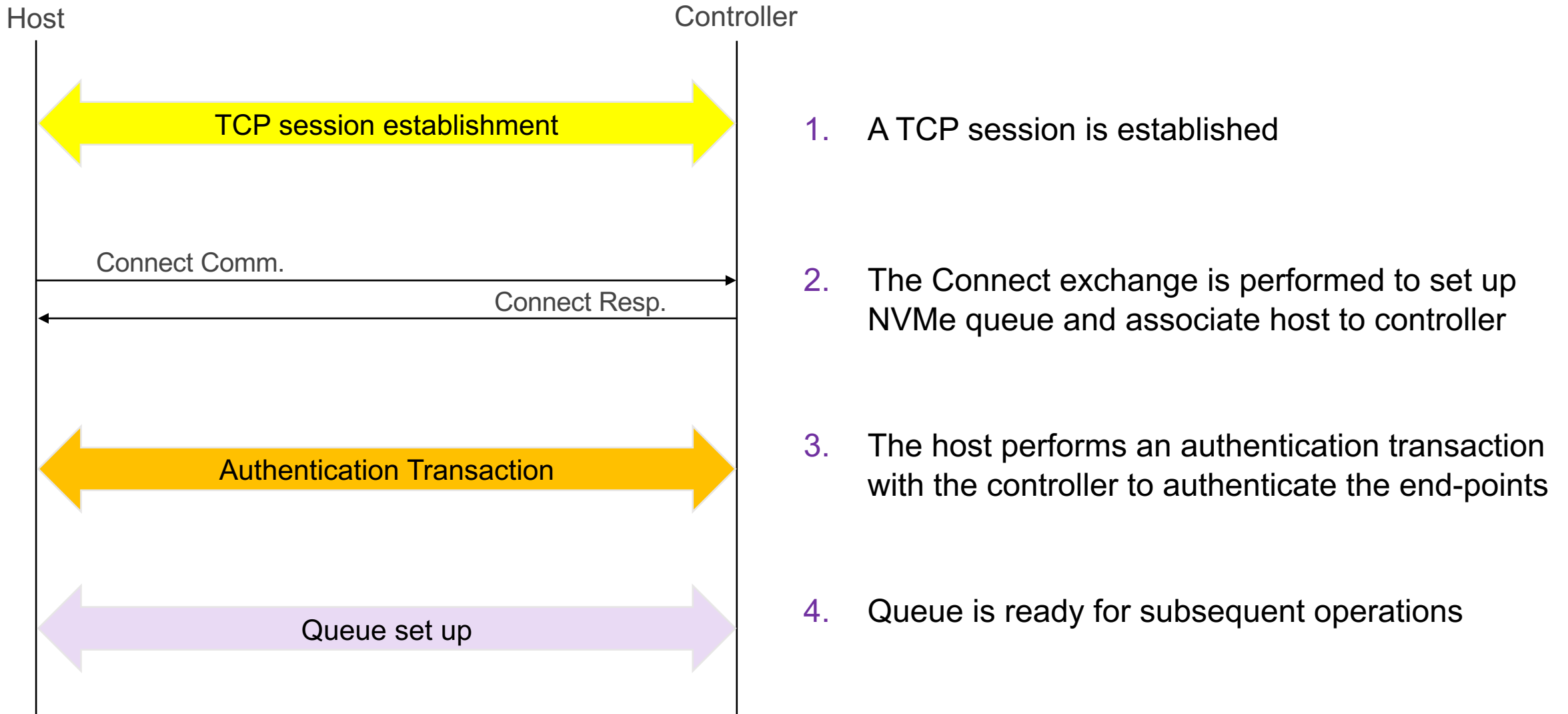
Securing Storage Area Networks Focus on NVMe-oF

Claudio DeSanti

SAN Protocols - Security Mechanism Comparison

	iSCSI	Fibre Channel	NVMe over Fabrics/IP
Storage Endpoint Authentication	CHAP (strong secret) SRP (weak secret, e.g., password) [not used in practice]	DH-CHAP (strong secret) FCPAP (weak secret, e.g., password) FCAP (certificates) FC-EAP (strong secret)	DH-HMAC-CHAP (strong secret)
Secure Channel (authenticated encryption & cryptographic integrity)	IPsec (e.g., in security gateway)	FC ESP_Header (IPsec-like)	TLS (pre-shared key) – for TCP only (not usable with RoCEv2 or iWARP) IPsec is also an option

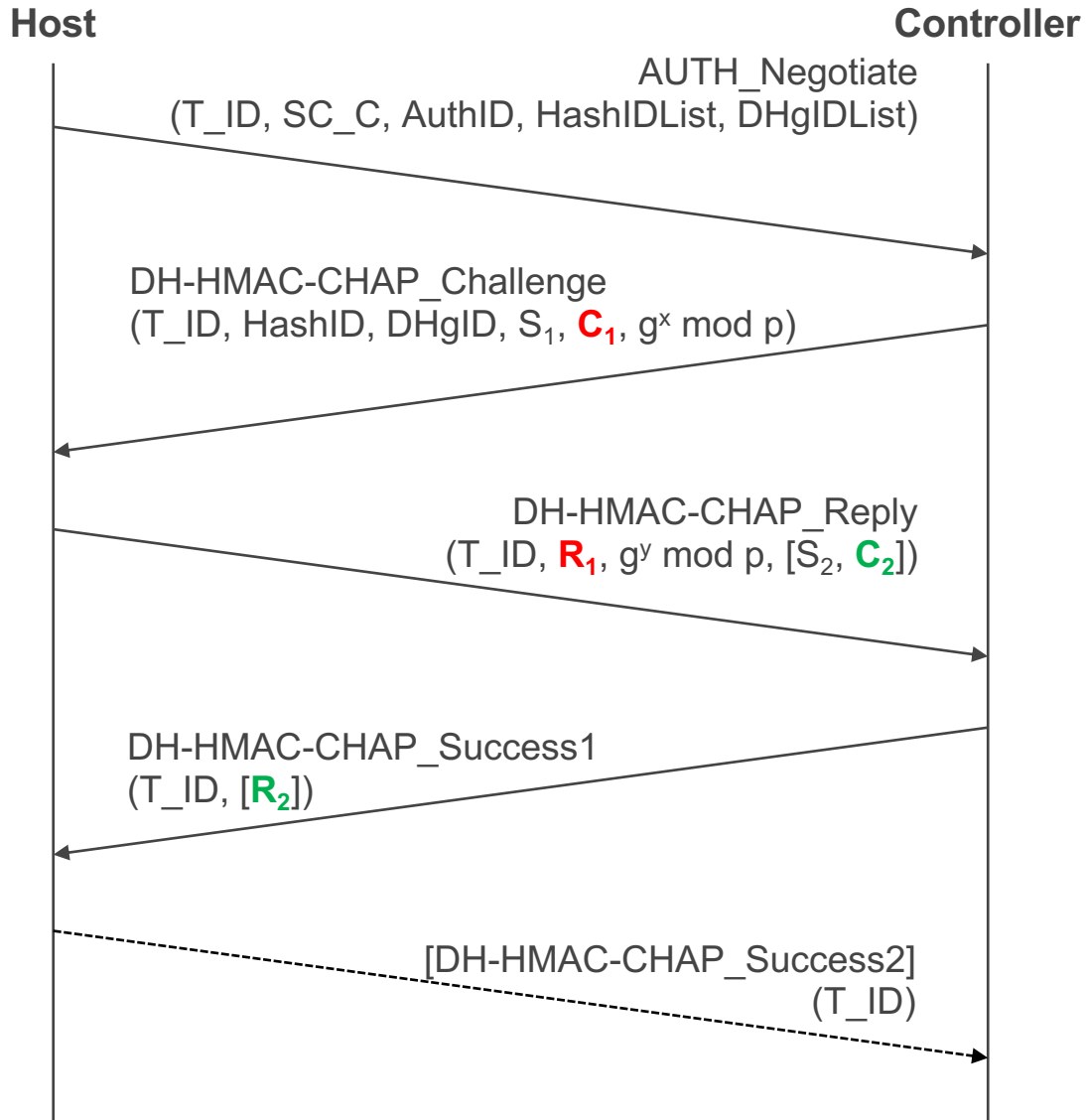
NVMe-oF Authentication Example



NVMe-oF Authentication: DH-HMAC-CHAP Protocol

- Defined in TP 8006
- Based on keys that need to be different for each NQN
- Challenge/response protocol: CHAP
 - the authenticator sends a challenge C;
 - the responder computes a response $R = \text{Hash}(C \parallel \text{key}_{\text{responder}} \parallel \text{other things})$;
 - the authenticator verifies the response (or delegates verification)
- **DH-HMAC-CHAP: Strengthened version of CHAP**
 - DH: Diffie-Hellman, adds (optional) key exchange to frustrate eavesdroppers
 - HMAC: Hashed MAC, uses secure hash twice to improve security
- Bidirectional authentication

DH-HMAC-CHAP Protocol: Bidirectional Authentication



Unidirectional challenge/response protocol

- Controller C sends a challenge C_1
- Host H computes a response $R_1 = \text{HMAC}(K_h, C_1 \parallel \text{other things})$
- Controller C verifies the response
- Unidirectional authentication (controller authenticates host)

Getting bidirectional authentication

- H sends a challenge C_2
- C computes a response $R_2 = \text{HMAC}(K_c, C_2 \parallel \text{other things})$
- H verifies the response
- Unidirectional authentication (host authenticates controller)

Verification:

- Controller computes R_1' and check if it matches the received R_1
- Host computes R_2' and check if it matches the received R_2

Authentication Responses

$$K_S = H((g^x \text{ mod } p)^y \text{ mod } p)$$

$$C_{a1} = (\text{DHgID} == 0) ? C_1 : \text{HMAC}(K_S, C_1)$$

$$R_1 = \text{HMAC}(K_h, C_{a1} \parallel S_1 \parallel T_ID \parallel SC_C \parallel \text{"HostHost"} \parallel \text{NQN}_h \parallel 00h \parallel \text{NQN}_c)$$

$$K_S = H((g^y \text{ mod } p)^x \text{ mod } p)$$

$$C_{a2} = (\text{DHgID} == 0) ? C_2 : \text{HMAC}(K_S, C_2)$$

$$R_2 = \text{HMAC}(K_c, C_{a2} \parallel S_2 \parallel T_ID \parallel SC_C \parallel \text{"Controller"} \parallel \text{NQN}_c \parallel 00h \parallel \text{NQN}_h)$$

SAN Protocols - Security Mechanism Comparison

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Secure Channel (authenticated encryption & cryptographic integrity)	IPsec (e.g., in security gateway)	FC ESP_Header (IPsec-like)	TLS (pre-shared key) – for TCP only (not usable with RoCEv2 or iWARP) IPsec is also an option

Secure Channel: TLS

- **TLS (Transport Layer Security): Widely used secure channel protocol**
 - Secure channel = authentication, confidentiality, cryptographic integrity (primary properties)
 - Typical (web) usage: Server uses certificate with TLS, client authenticates after TLS setup (e.g., TLS-protected HTTP)
- **TLS versions:**
 - TLS 1.0 and 1.1: Obsolete - should not be used
 - TLS 1.2: Baseline TLS version, widely implemented and used, getting replaced by TLS 1.3
 - TLS 1.3: New version, complete protocol redesign (TLS 2.0 in practice), usage rolling out
 - [Support available in some security libraries \(e.g. OpenSSL & LibreSSL\), expanding to more](#)
- **TLS 1.3 specification for NVMe-oF/TCP: completed in TP 8011**
 - TLS not specified for other NVMe-oF IP-based protocol (i.e., RDMA, e.g., RoCEv2)
 - Based on pre-shared keys (PSKs)
- **NVMe-oF/TCP TLS 1.2: discouraged by TP 8011**
 - Usage was specified by NVMe-oF 1.1 standard

Why TLS 1.3?

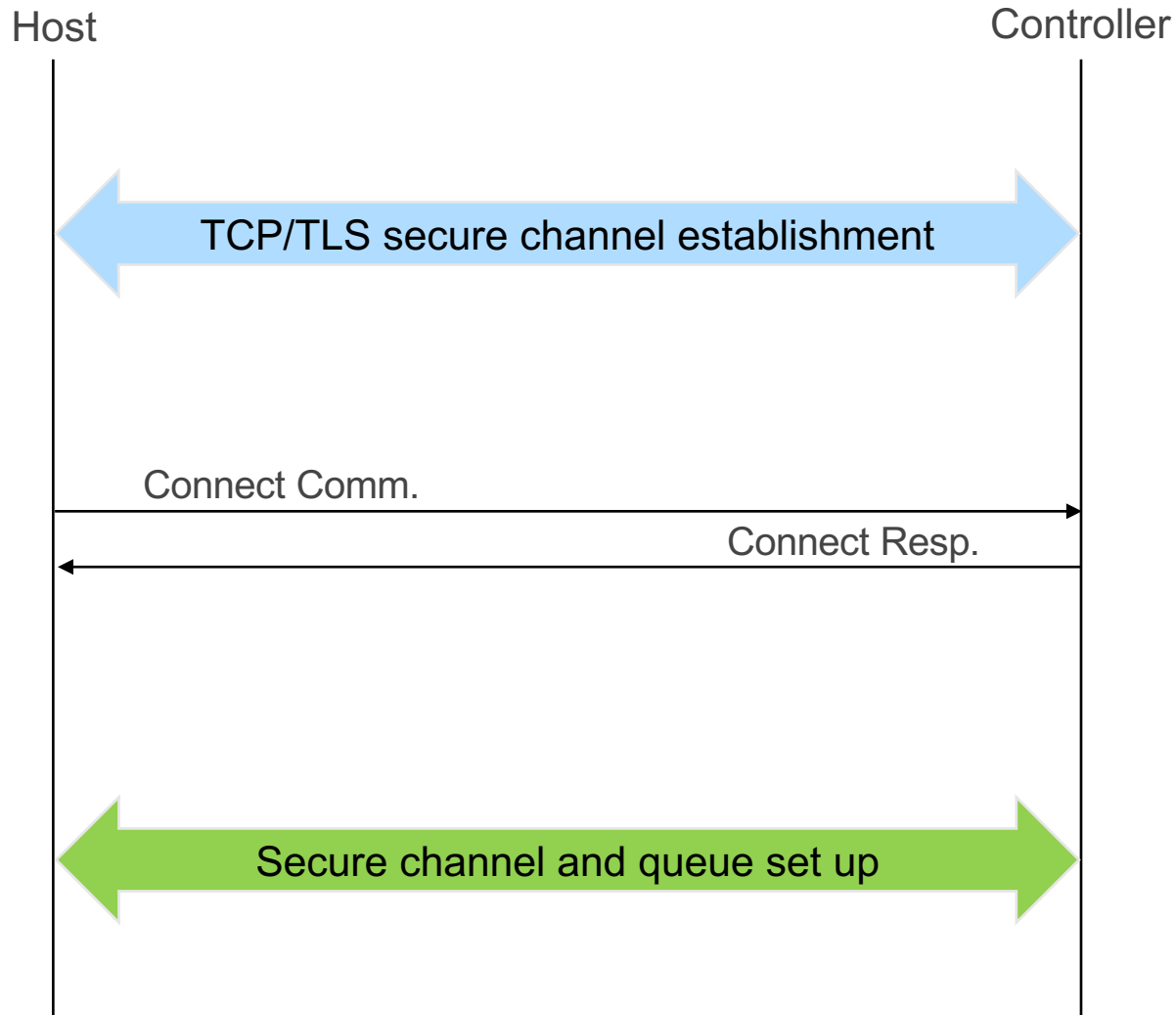
TLS 1.2

- IANA TLS registry has 300+ cipher suite code points
 - Uncertain security properties, difficult interoperability
- Encryption starts late in the handshake
 - Client cert and target site are sent in the clear
 - Poor privacy
- Many features with known security flaws

TLS 1.3

- 5 cipher suites, all with PFS and modern algorithms
 - Consistent security properties
- Encryption starts as early as possible, hiding content length
 - Minimal set of cleartext protocol bits on the wire
 - Less user information visible to the network
- All those features are omitted from 1.3

TLS 1.3

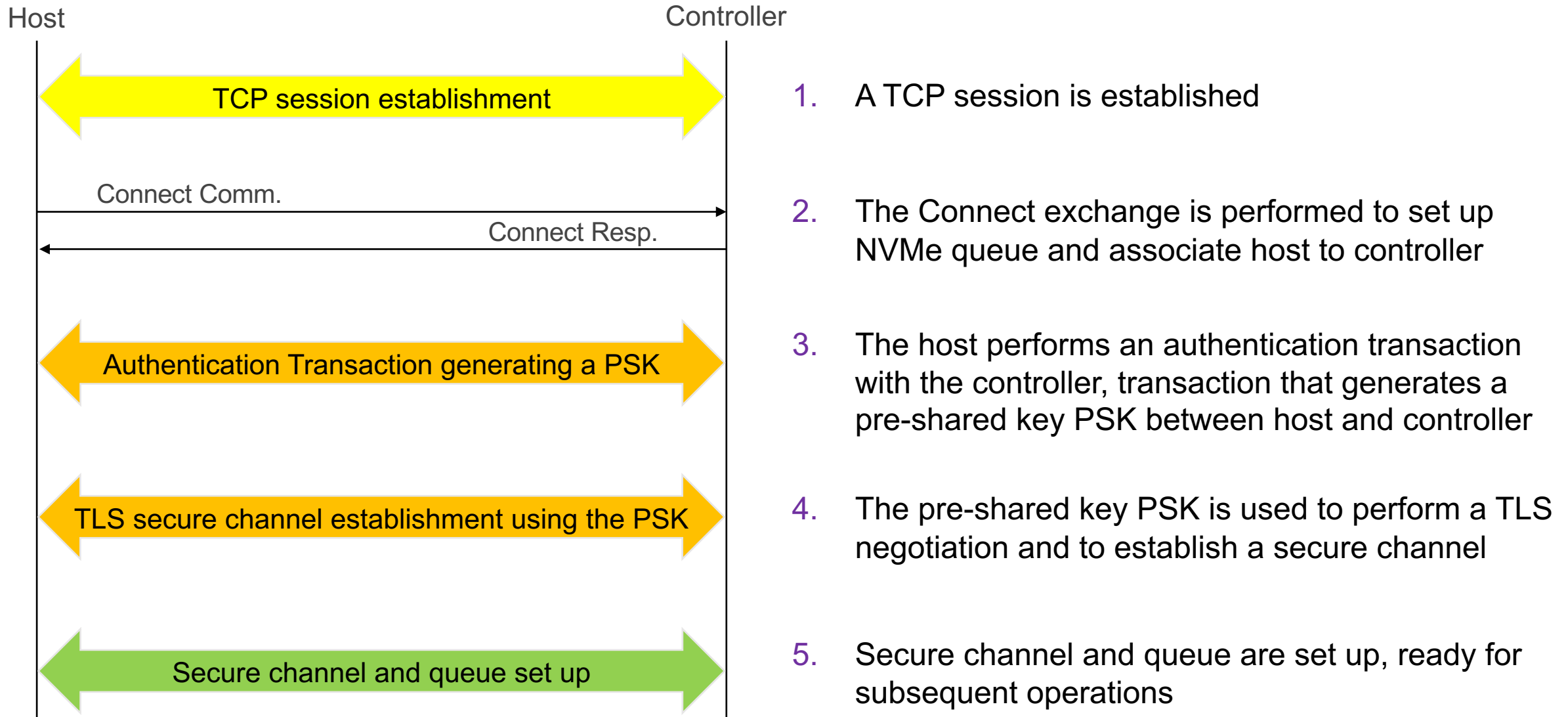


1. A TCP/TLS session negotiation is performed and a secure channel is established
2. The Connect exchange is performed to set up NVMe queue and associate host to controller
3. Secure channel and queue are set up, ready for subsequent operations

TLS Credentials

- TLS secure channel for NVMe-oF/TCP is based on pre-shared keys (PSKs)
 - In order to authenticate and establish a secure channel between themselves, two NVMe entities need to be configured with the same PSK
 - This can lead to a deployment option called ‘group PSK’: all NVMe entities share the same PSK
 - Big security concern (compromising a single node may allow an attacker to access all secure channels)
 - The proper way would be to have a PSK per each pair of entity that can communicate (n^2 problem)
- Authentication protocols to the rescue
 - Upon successful completion of an authentication exchange, the two involved NVMe entities generate an ephemeral shared session key (e.g., a ‘PSK’ computed on the fly)
 - The TLS negotiation can then be performed using a PSK derived from that shared key
 - No more need for ‘group PSK’
 - Implementation result: the TCP connection begins unsecured and then transitions to secured
 - Opportunistic TLS
 - Linear problem (not anymore n^2): need just one secret per entity

Authentication Followed by TLS



1. A TCP session is established
2. The Connect exchange is performed to set up NVMe queue and associate host to controller
3. The host performs an authentication transaction with the controller, transaction that generates a pre-shared key PSK between host and controller
4. The pre-shared key PSK is used to perform a TLS negotiation and to establish a secure channel
5. Secure channel and queue are set up, ready for subsequent operations

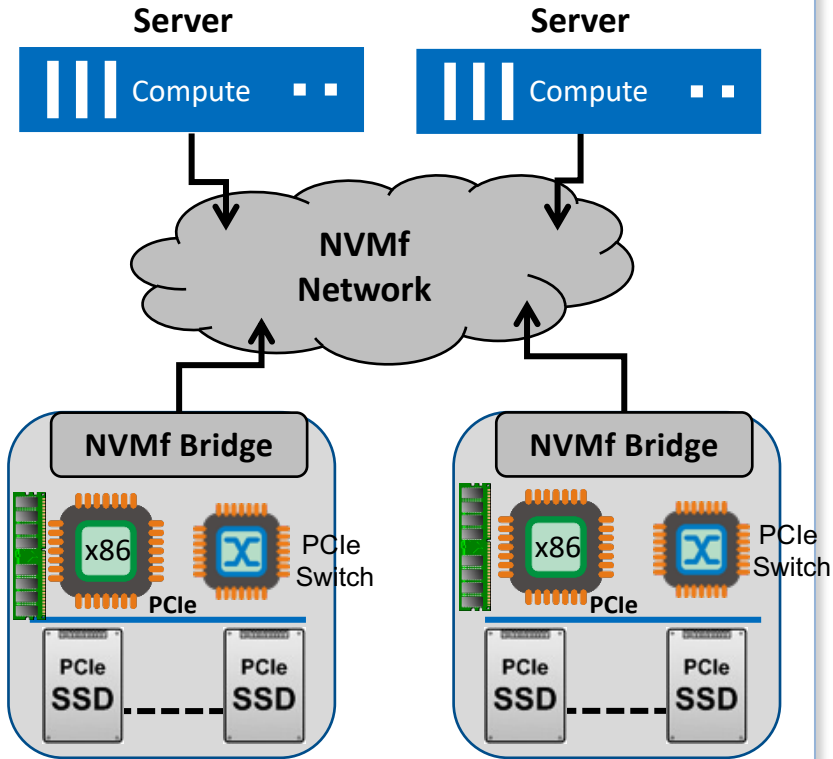


NVMe-oF Security in Action A Use Case in E-SSDs

Hrishikesh Sathawane and Eric Hibbard

NVMe Over Fabrics Architectures for Disaggregated Storage

JBOF with x86

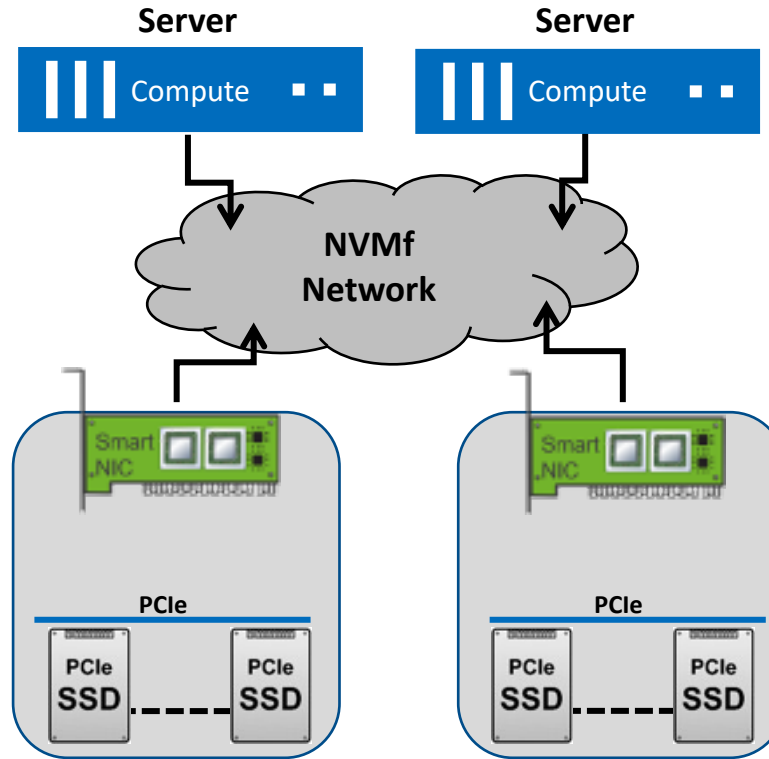


Pros
 Current production
 Established Ecosystem
 PCIe SSDs in production

Cons
 BW bottleneck
 Added PCIe latency
 High power

Production

JBOF with SmartNIC

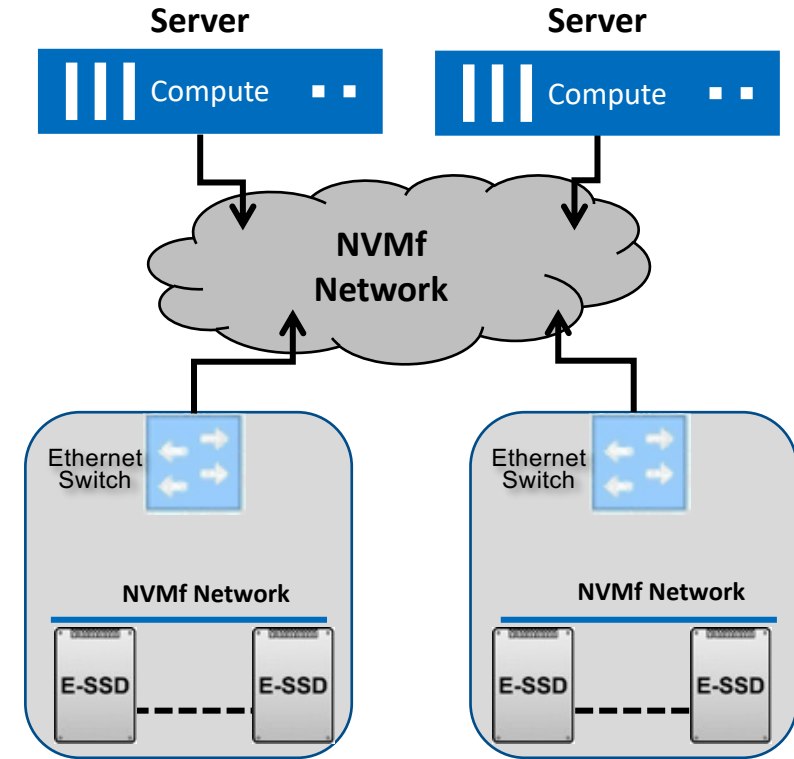


Pros
 Finding some use cases
 Emerging Ecosystem
 PCIe SSDs in production
 Lower Power

Cons
 Added PCIe latency
 Needs SmartNIC on both sides

Engage-Short Term

E-BOF



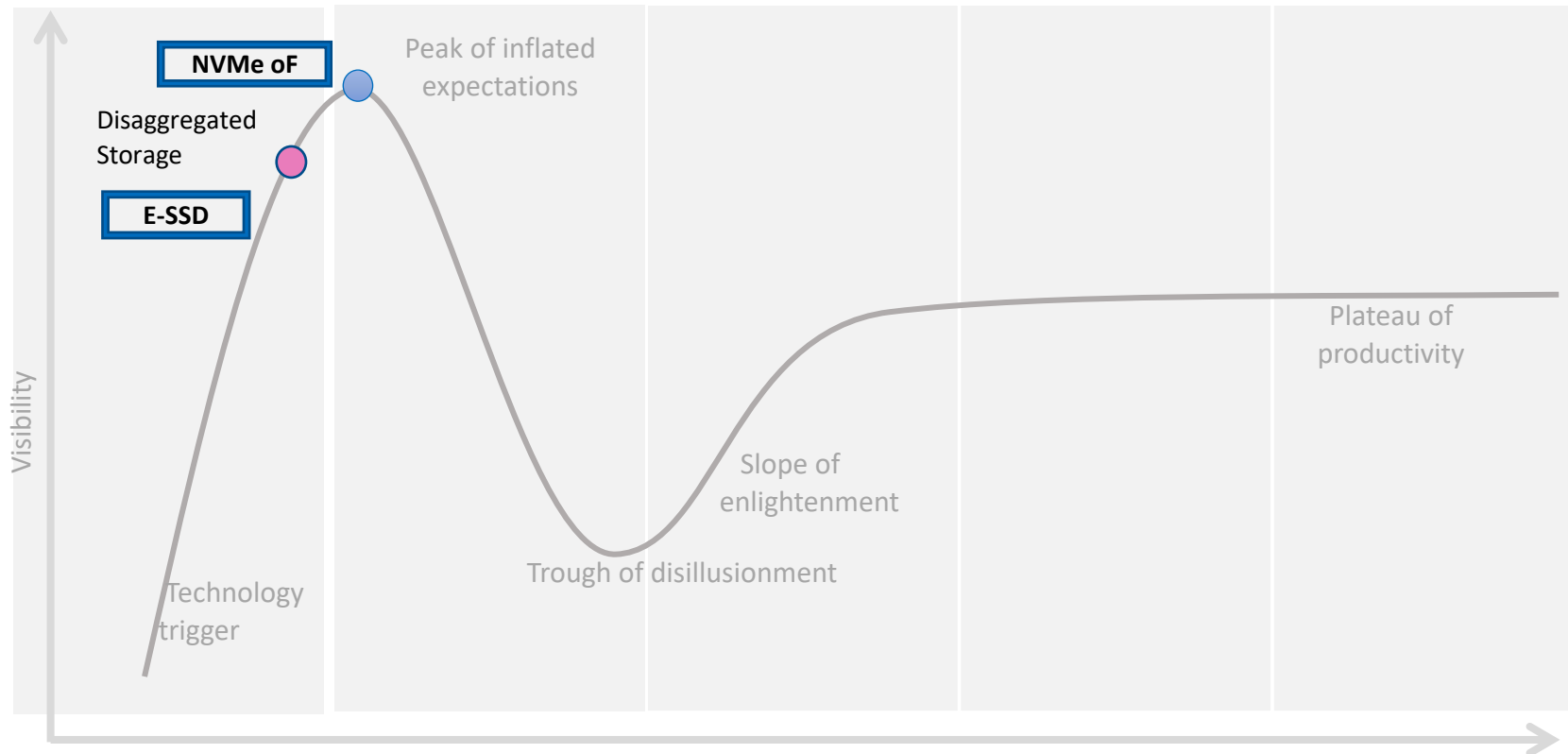
Pros
 BOM cost savings
 Lower latency/power
 True disaggregation possible

Cons
 No Ecosystem yet
 E-SSDs are in PoC

Engage-Long Term

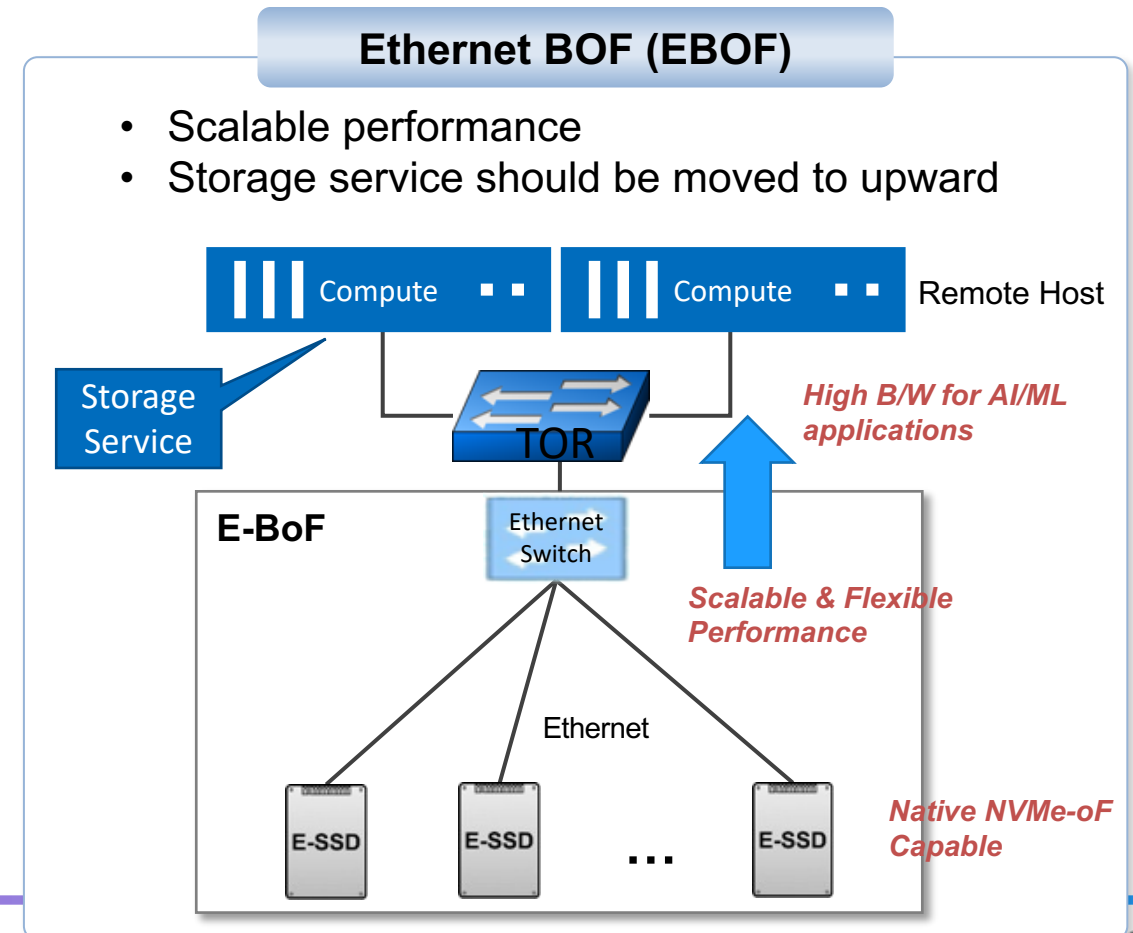
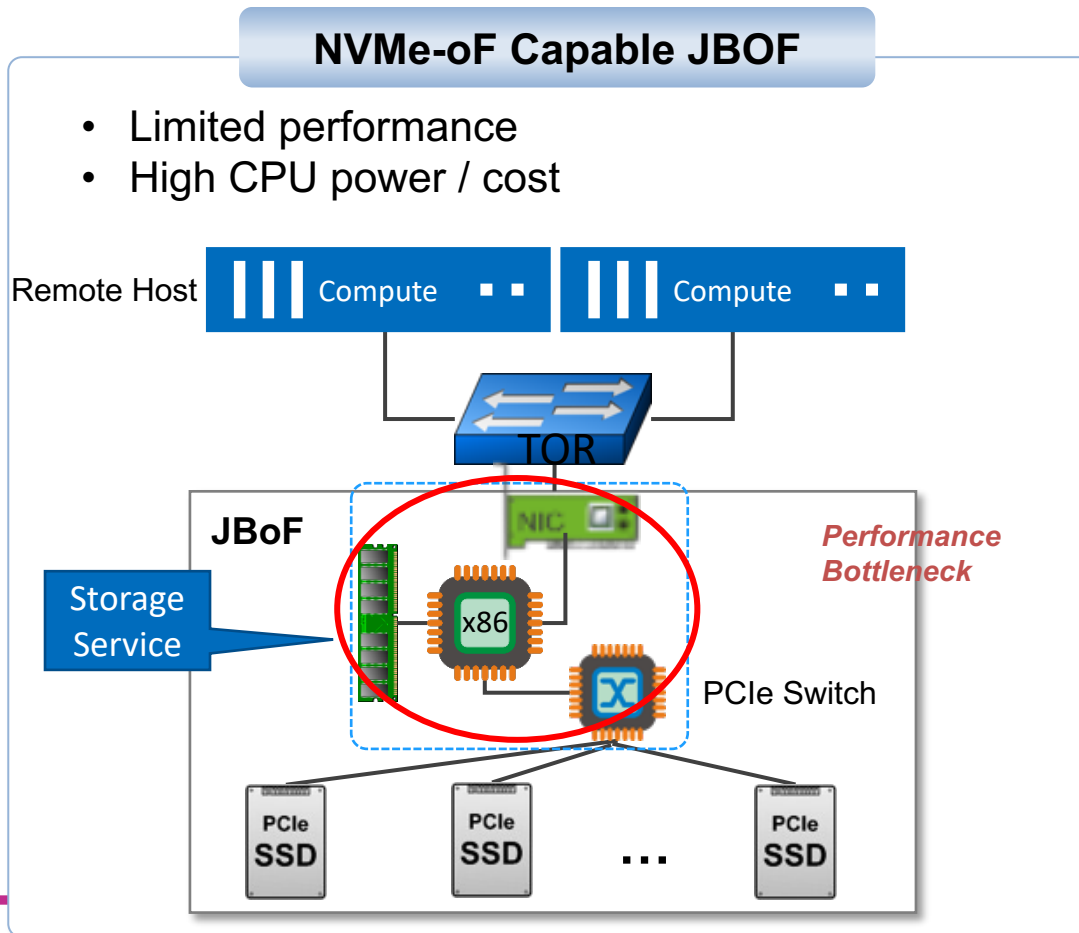
E-SSD and the Hype Cycle

E-SSD is in the early phase of the hype cycle.



Ethernet SSD Introduction and Value Proposition

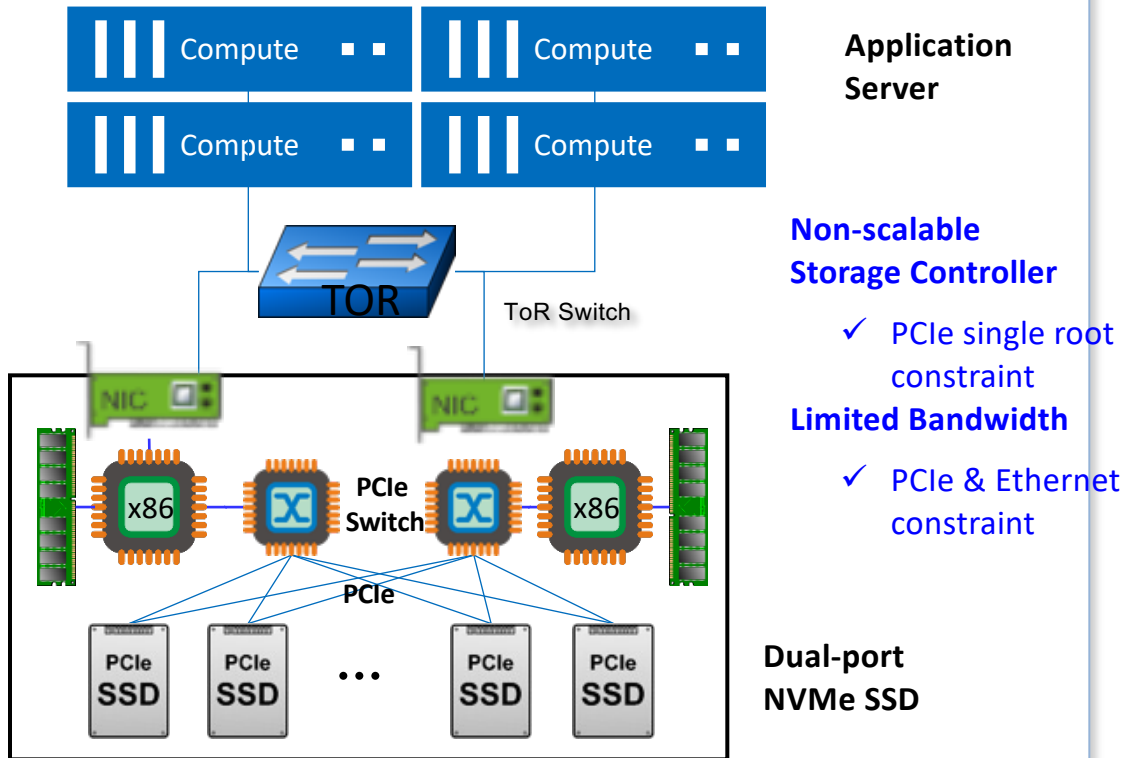
- Ethernet SSD can be one of the solution for scalable performance of JBOF
 - Also targeting for TCO saving through replacement from high CPU & BOM to Ethernet switch
 - Samsung is also continuing to study the architecture, ecosystem and benefits of Ethernet SSD



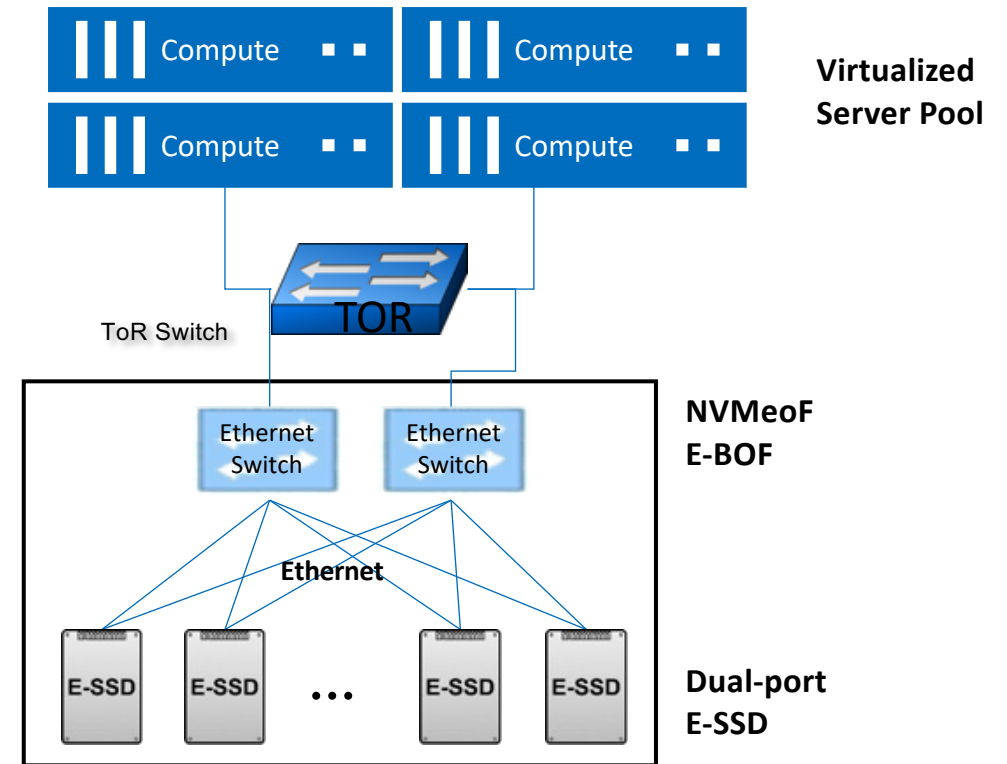
Advantage of NVMe-oF SSD

- NVMeoF JBOF can solve performance, scalability, and flexibility
 - Scalable and flexible data center solution through Ethernet-only infrastructure

Conventional SSD System

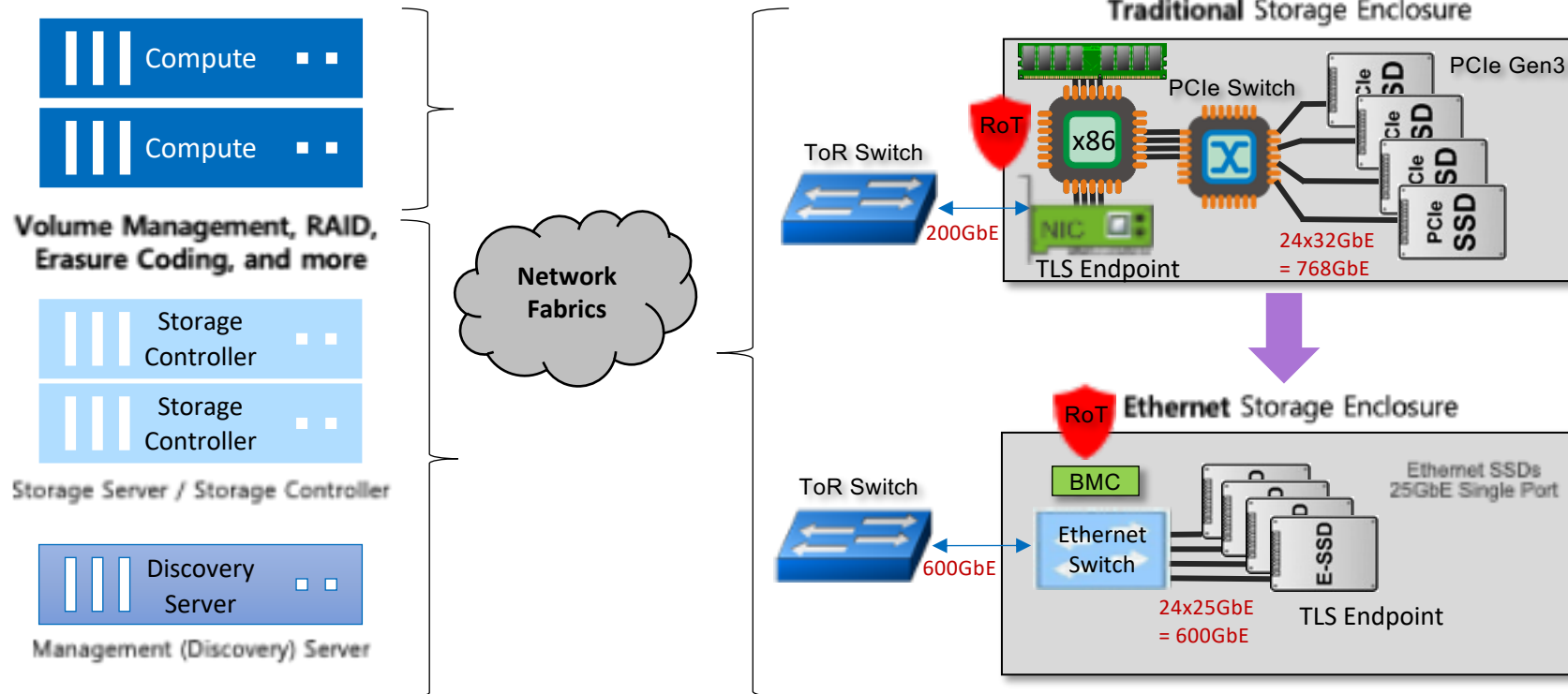


NVMeoF SSD System



Disaggregated Architecture

- **NIC card's essential features are offloaded to E-SSD.**
- **Storage controller can communicate with multiple E-BOFs as it is done with JBOFs.**
 - Connection and discovery are added to maintain connection in NVMe-oF specification.
 - NVMe I/O command processing is equal as PCIe based NVMe SSD for E-SSD.
 - Additional target configuration may not be needed, since configured information are saved in E-SSD.
- **SW modification**
 - Schema of JBOF for Redfish can be changed.



NIC + CPU

- IP address configuration
- Manage Connections
- Network Management

PCIe SSD

- NVMe I/O Processing

BMC (Expected)

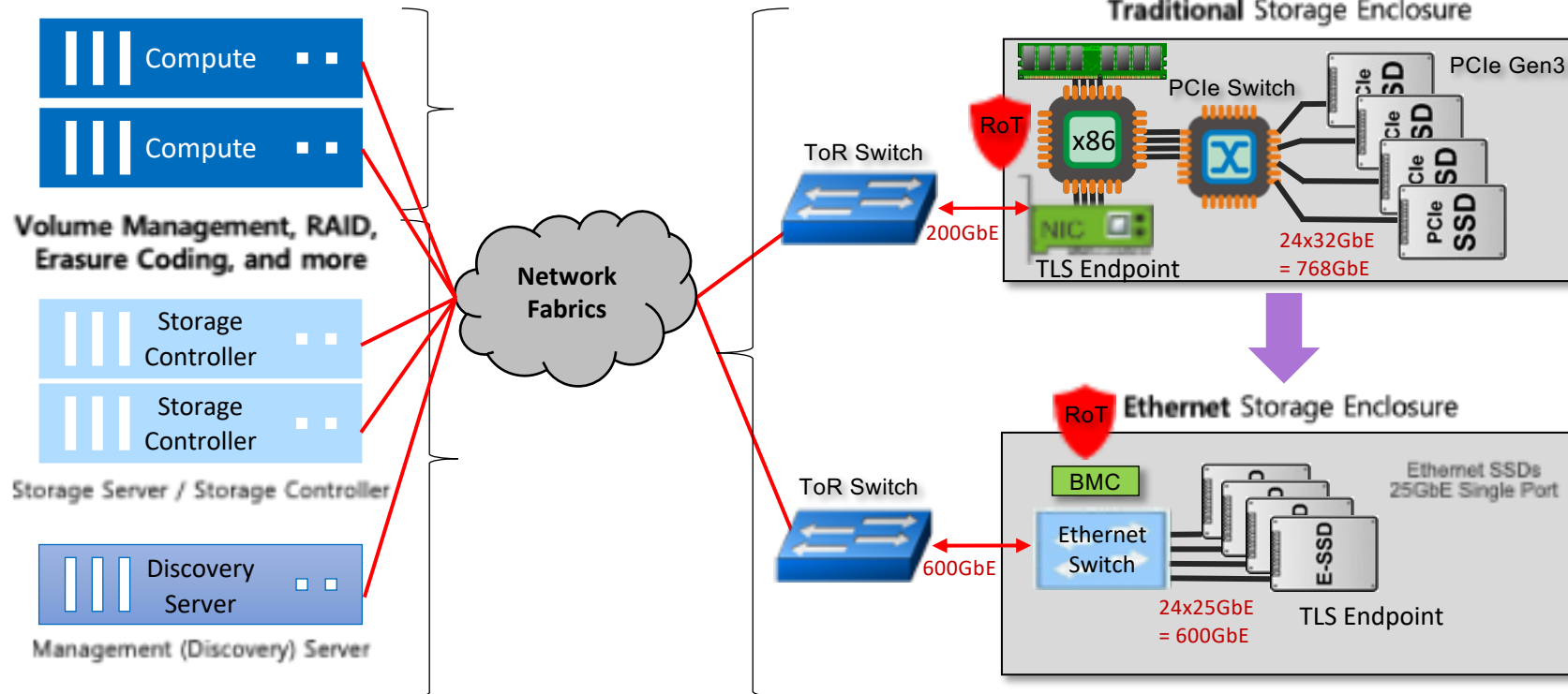
- E-SSD IP address configuration

E-SSD (offloaded)

- IP address configuration (manual)
- Manage Connections
- Network Management
- NVMe I/O Processing

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PCIe SSD

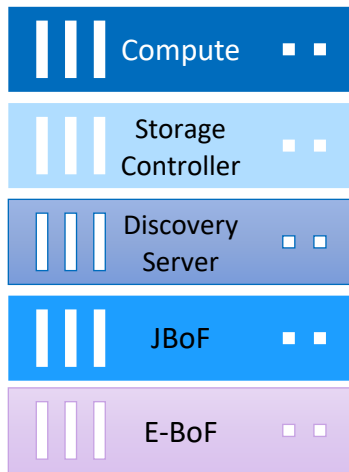
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BMC (Expected)

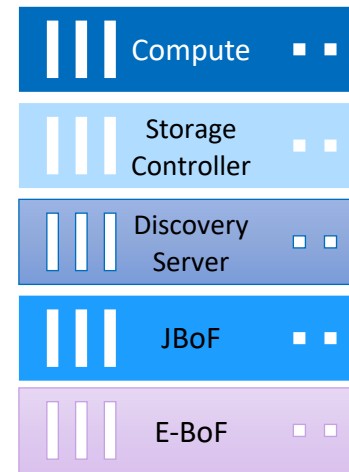
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E-SSD (offloaded)

- IP address configuration (manual)
- Manage Connections
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- NVMe I/O Processing



Questions



Other Resources

- Webcast: NVMe over Fabrics: Looking Beyond Performance Hero Numbers
 - <http://bit.ly/NVMeoFHero>
- Multiple resources: SNIA Geek Out on NVMe-oF
 - <https://bit.ly/GeekNVMeoF>
- Blog: NVMe over Fabrics for Absolute Beginners
 - <https://bit.ly/3aKf3JS>

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Thank You