

PM Support in Linux and Windows

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Windows Support for Persistent Memory

Availability of Windows PM Support



Client Workstation:

August, 2016: Windows 10 Anniversary Update

April, 2017: Windows 10 Creators Update

October, 2017: Windows 10 Fall Creators Update

April, 2018: Windows 10 April 2018 Update

October, 2018: Windows 10 October 2018 Update

Server:

September, 2016: Windows Server 2016

November, 2019: Windows Server 2019

Supported Hardware:

- JEDEC NVDIMM-N
- HPE Scalable Persistent Memory

DAX Support



- Supported since Windows Server 2016
- Advantages:
 - Improved IO performance by eliminating OS overhead
- Disadvantages:
 - Functionality loss due to elimination of OS hook points, examples:
 - Software Encryption
 - Software Compression
- Only NTFS supports DAX

How DAX Works

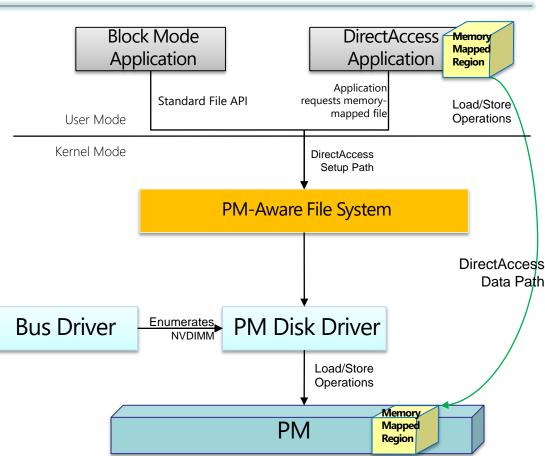


The Idea

- App has direct access to PM via existing memory-mapping semantics
- Updates directly modify PM
 - Storage Stack not involved

Characteristics

- True device performance (no software overhead)
- Byte-Addressable



IO on Windows DAX Volumes



Non-Cached IO

Is converted to cached IO by the file system

Cached IO

- Cache Manager maps directly to persistent memory
- Copies directly between user's buffer and persistent memory

Memory Mapped IO

- Memory mapped sections point directly to PM hardware
- Provides applications with zero-copy access to PM hardware

Windows PM Block Mode Volume Support



Is fully backwards compatible

- Maintains existing storage semantics
 - Sector atomicity supported by the PM disk driver
- Fully compatible with existing applications and filter drivers
- Supported by all Windows file systems

Windows Hyper-V PM Support



- Available in Windows Server 2019
- Windows & Linux guests in generation 2 VMs see virtual PMEM (vPMEM) devices
 - Memory mapped files in the guest have direct access to PM hardware on the host
 - Full Win32 and PMDK support
- New VHD file type: .VHDPMEM

What are Large and Huge Pages



- Modern CPUs manage memory using 4K pages
- An applications memory usage is managed via page tables controlled by the operating systems memory manager
- CPU's contain a mapping table cache called the TLB (translation lookaside buffer) that caches page table mappings
- For applications with a large memory footprint -- the CPU can spend a lot of time reading page table entries into the TLB
- A Large Page allows a contiguous 2mb region to be described with a single TLB entry
 - Applications typically see a significant performance improvement
- A Huge Page allows a contiguous 1gb region to be described with a single TLB entry

Windows PM support for Large & Huge Pages



- Available in Windows Server 2019
- DAX partitions are aligned to 2mb boundaries
- NTFS now supports cluster sizes up to 2mb (in powers of 2)
 - In Server 2016 the limit was 64K
- A memory mapped file on a DAX volume with a 2mb cluster size is guaranteed to be mapped using at least Large Pages
- Large and Huge page alignment is supported on cluster sizes <2mb</p>

Windows PM Management Support



- Windows Server 2019 introduces Powershell support for managing physical and logical persistent memory devices
 - Ability to enumerate, create and delete logical persistent memory devices
 - Ability to enumerate and initialize physical persistent memory devices
 - Example cmdlets:
 - Get-PmemDisk
 - > New-PmemDisk
 - > Remove-PmemDisk

- Get-PmemPhysicalDevice
- > Initialize-PmemPhysicalDevice
- Get-PmemUnusedRegion

Windows PMDK Support



- Available since Windows Server 2016
- Defines a set of application API's for efficient use of PM hardware
 - Abstracts out OS specific dependencies
 - Underlying implementation uses memory mapped files
 - Makes its own atomicity guarantees
 - Works in both PM and non-PM environments
 - Simpler application development model
- Open source library available for Windows and Linux via GitHub
 - https://github.com/pmem/pmdk/

Examples uses of PM on Windows



SQL Server 2016

- Tail-of-the-Log: Transactions committed to DAX volume
- 2X performance gain by:
 - Halving the transaction latency
 - De-serializing the SQL logging threads

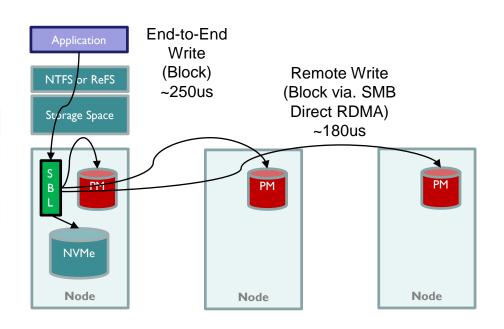
Write Caching for HCI with Storage Spaces Direct & Windows Server 2019



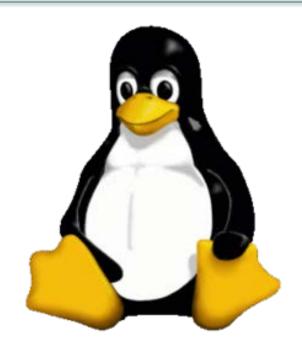
- Ultra-low-latency write caching (SBL) for all-NVMe deployments
- Industry record performance demonstrated at Ignite '18

| Benchmark | Performance |
|---------------------|---------------------|
| 4K 100% Rand Read | 13.8 Million IOPs |
| 4K 90/10% Rand R/W | 9.45 million IOPs |
| 2MB Sequential Read | 549 GB/s Throughput |

12-nodes with Triple-Mirrored Scoped Spaces with ReFS
12 x Intel® S2600WFT, 384 GiB memory, 2 x 28-core "CascadeLake", 1.5 TB
Intel® Optane™ DC persistent memory as cache, 32 TB NVMe (4 x 8TB Intel®
DC P4510) as capacity, 2 x Mellanox ConnectX-4 25 Gbps







Linux Support for Persistent Memory

Track development via linux-nvdimm mailing list!

Managing PM in Linux: ndctl



11 23 contributors

O Watch = 28 ★ Star 45 ¥ Fork 34

2 years ago

4 months ago

6 months ago

2 years ago

4 years ago

Latest commit cb2d678 on Oct 5, 2018

- How do you determine if you have PM in your system?
- How do you manage the PM in your system?
- How do you determine the health of PM in your system?
- Ties into physical layer specifications like NFIT and HMAT [1].



ndctl, list: fix the verbosity level formatting in the man page

ndctl list; always output array without --human

ndctl: remove unreferenced ccan license file

ndctl: Revert *ndctl: Create ndctl udev rules for dirty shutdown

○ 78 releases

1 9 branches

ndetl: clamp dimm formats

Import initial infrastructure

pmem / ndctl

@ 840 commits

iiii contrib

IIII dayeti

iiii m4

stellarhopper ndctl: release v63 ==

Using PM in Linux: A block device



- PM can be consumed by Linux applications using a block interface.
- Useful when a file-system is not required.
- Not super-efficient as block device changes on sector, not byte granularity.
- All your favourite block-device tools can be used.

```
O
                     batesste@tyrone: ~/vpns — ssh vm
generic@pmsummit-1:~$ sudo fdisk -l /dev/pmem0
Disk /dev/pmem0: 15.9 MiB, 16646144 bytes, 32512 sectors
Units: sectors of 1 * 512 = 512 bytes
Sector size (logical/physical): 512 bytes / 4096 bytes
I/O size (minimum/optimal): 4096 bytes / 4096 bytes
[generic@pmsummit-1:~$
[generic@pmsummit-1:~$ lsblk
NAME
                   SIZE RO TYPE MOUNTPOINT
vda
       252:0
                    12G
                         0 disk
`-vda1 252:1
                    12G
                         0 part /
      259:0
                0 15.9M 0 disk
pmem0
[aeneric@pmsummit-1:~$
generic@pmsummit-1:~$
```

Using PM in Linux: A filesystem



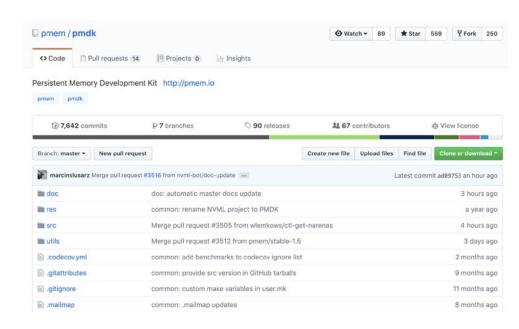
- Like any standard block device we can put a filesystem over it.
- In this case we can put ANY filesystem we like on top.
- But if we use a DAX aware filesystem we get added benefits (PM optimizations).
- Now we can have files and directories and all that good stuff!
- In Linux both EXT4 and XFS have DAX support.

```
quest@pmsummit-1:~$ mkfs.xfs -f /dev/pmem0
mkfs.xfs: cannot open /dev/pmem0: Permission denied
[guest@pmsummit-1:~$ sudo mkfs.xfs -f /dev/pmem0
[[sudo] password for quest:
meta-data=/dev/pmem0
                                 isize=512
                                               agcount=4, agsize=56320 blks
                                 sectsz=4096
                                              attr=2, projid32bit=1
                                 crc=1
                                               finobt=1, sparse=0, rmapbt=0, reflink=0
                                 bsize=4096
                                              blocks=225280, imaxpct=25
data
                                 sunit=0
                                               swidth=0 blks
         =version 2
naming
                                 bsize=4096
                                               ascii-ci=0 ftype=1
                                 bsize=4096
                                              blocks=1605, version=2
         =internal log
                                              sunit=1 blks, lazy-count=1
                                 sectsz=4096
realtime =none
                                 extsz=4096
                                              blocks=0, rtextents=0
[guest@pmsummit-1:~$ sudo mount -o dax /dev/pmem0 /mnt/
[quest@pmsummit-1:~$ dmesq | tail
    16.212472] pmem0: detected capacity change from 0 to 130023424
  184.699585] pmem0: detected capacity change from 0 to 134217728
   203.539578] pmem0: detected capacity change from 0 to 939524096
   210.083474] pmem0: detected capacity change from 0 to 922746880
   232.851656] random: crng init done
   232.851694] random: 7 urandom warning(s) missed due to ratelimiting
[10604.527337] SGI XFS with ACLs, security attributes, realtime, no debug enabled
[10604.588508] XFS (pmem0): DAX enabled. Warning: EXPERIMENTAL, use at your own risk
[10604.588563] XFS (pmem0): Mounting V5 Filesystem
[10604.627537] XFS (pmem0): Ending clean mount
guest@pmsummit-1:~$ ■
```

Using PM in Linux: mmap()



- Now we have files that live in a PM-aware filesystem on a PMaware block device on physical PM.
- mmap() has been around for a while ;-). Maps the file into the Virtual Address space of the running process.
- Now the world becomes our oyster (see PMDK for more!)



git clone https://github.com/pmem/pmdk.git



Thank You