10 Million I/Ops From a Single Thread

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Intel
Storage Performance Development Kit

- User-space block storage stack with similar features to an OS
  - Includes an NVMe driver
- Open Source, BSD 3-Clause License ([https://spdk.io](https://spdk.io))
- Very active community
  - 1200 commits from 56 committers in last 3 months
BIG NUMBERS

4KiB random reads at queue depth 128 to each device

10.39M iops

SPDK nvme perf
# System Configuration

<table>
<thead>
<tr>
<th>Configuration</th>
<th>Details</th>
</tr>
</thead>
<tbody>
<tr>
<td>CPU</td>
<td>Intel® Xeon® Platinum 8280L CPU @ 2.70GHz</td>
</tr>
<tr>
<td>Memory</td>
<td>12x 16GB 26667</td>
</tr>
<tr>
<td>Storage</td>
<td>21x Intel® SSD DC P4600 1.6TB</td>
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</tbody>
</table>
Setting the Stage

• NVMe drives are getting really, really fast
  • No, like really.

• PCIe Gen 4 and eventually Gen 5 are huge bandwidth boosts.

• We should expect storage servers doing tens of millions of I/Ops to be the *norm* going forward.
Overview

Talk based on this blog post:
https://spdk.io/news/2019/05/06/nvme/

Covering only 3 techniques from that post today for time reasons.

Not covering any active areas of research, but there are several!
Background: NVMe I/O Queues
Mechanics of Submitting an I/O

NVMe queues consist of two arrays in host memory (submission queue and completion queue) plus two doorbells (SQTAIL, CQHEAD) in the BAR.

Submission queue

Complete Queue
Mechanics of Submitting an I/O

To submit: Copy command into next slot. Write SQTAIL.

![Diagram of SQHEAD and SQTAIL]

To complete: Compare phase bit. If flipped, contains valid completion. When done, write CQHEAD.
Completions hold updates to SQHEAD.
SPDK NVMe Design Basics

• Assign NVMe queue pairs to threads.
  • No locks!
• Disable interrupts
  • Completions are handled when the application is ready to handle them. No context switch.
  • Code is compiled with –O2, LTO enabled, PGO disabled.
Keys To Performance

Don’t let the CPU stall!

- No cross-thread coordination (locks, etc.)
- Poll instead of interrupt
- Minimize MMIO
- Get the right things into the CPU cache at the right time
- Pack structures. Separate hot data from cold.
Strategy 1: Minimize MMIO
Tricks For Minimizing MMIO

while (true) {
    ... work ...
    spdk_nvme_ns_cmd_read(..., cb_fn);
    spdk_nvme_ns_cmd_read(..., cb_fn2);
    ... work ...

    spdk_nvme_qpair_process_completions(...); /* calls cb_fns if the read is done */
}

Program Flow: Submit several commands. Check for completions. Repeat.

Naïve implementation: For each command, 1 MMIO on submit, 1 MMIO on complete
Trick 1

- When checking for completions by reading the phase bit, don’t write the completion queue doorbell until all currently outstanding completions have been discovered.

- Everyone knows this trick.
Trick 2

- When a completion is posted, don’t write the completion queue doorbell unless the device actually needs more slots free.

- Devices often have large queues – say 1024. We can write the completion queue doorbell every ~512 commands.

- I’ve only seen SPDK do this. (Doesn’t help this benchmark because trick 1 is already finding 30 to 50 completions at a time)
Trick 3

- When a command is submitted, copy the command into the SQE slot but don’t ring the submission queue doorbell. Instead, ring the doorbell only when the user polls (which is happening very frequently).

- This is a tricky way to batch command submissions transparently to the user in a polling system.

2.89M I/Ops with this disabled. 10.39M with this enabled.
Strategy 2: Optimize Completion Handling
Completing I/O

```
for each cqe {
    if (cqe->phase != phase_flipped) {
        break;
    }

    struct tracker *tr = tracker_array[cqe->cid];

    struct request *req = tr->req;

    req->cb_fn(req->cb_arg);
}
```

- “Tracker” is the context for a command. Array of trackers pre-allocated.
- CQE points (via CID) at a tracker (1:1)
- Tracker points at a request (1 request to N trackers)
Eliminate Data Dependent Loads

```c
struct nvme_request {
    spdk_nvme_cmd_cb cb_fn; /* Callback function */
    void *cb_arg;
};

struct nvme_tracker {
    struct nvme_request *req;
    spdk_nvme_cmd_cb cb_fn; /* Copied callback function */
    void *cb_arg;

    /* Other stuff */
};
```

On submission, copy cb_fn and cb_arg into tracker
for each cqe {
    if (cqe->phase != phase_flipped) {
        break;
    }

    struct tracker *tr = tracker_array[cqe->cid];

    tr->cb_fn(tr->cb_arg); /* Can call cb_fn immediately! */
}
Strategy 3: Pre-fetch
for each cqe {
    if (cqe->phase != phase_flipped) {
        break;
    }

    next_cqe = cqe + 1;
    if (next_cqe->phase == phase_flipped) {
        __builtin_prefetch(tracker_array[next_cqe->cid]);
    }

    __builtin_prefetch(next_cqe + 1);

    struct tracker *tr = tracker_array[cqe->cid];

    struct request *req = tr->req;

    req->cb_fn(req->cb_arg);
}
Questions?

- https://spdk.io
- https://spdk.io/news/2019/05/06/nvme/
- https://spdk.io/doc/ssd_internals.html
Thank You