

### DAOS An Architecture for Extreme Scale Storage

Eric Barton High Performance Data Division Intel

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# **Emerging Trends**

FOLDERS VS METADATA

Increased computational power...

- Huge expansion of simulation data volume & metadata complexity
- Complex to manage and analyze
- ...achieved through parallelism
- 100,000s nodes with 10s millions cores
- More frequent hardware & software failures

Tiered storage architectures

- High performance fabric & solid state storage on-cluster
- Low performance, high capacity disk-based storage off-cluster



# **Disruptive Change**

### NVRAM + Integrated fabric

- Byte-granular storage access
- Sub-µS storage access latency
- µS network latency
- Conventional storage software
- Block granular access limits scaling
- High overhead
  - 10s µS lost to communications S/W
  - 100s µS lost to F/S & block I/O stack

### I/O stack requirements

- Minimal software overhead
  - OS bypass
    - Communications
    - Latency sensitive I/O
- Fail-out resilience
- Persistent Memory storage
  - Filesystem & application metadata / hot data

Integrated Fabric Persistent

**inte** 

Memory

- Block storage
  - SSD warm data / Disk lukewarm data

# **Storage Architecture**

Compute Node NVRAM

- Hot data
  - High valence & velocity
  - Brute-force, ad-hoc analysis
  - Extreme scale-out
- Full fabric bandwidth
  - O(1PB/s)→O(10PB/s)
- Extremely low fabric & NVRAM latency
  - Extreme fine grain
  - New programming models



- Semi-hot data / staging buffer
- Fractional fabric bandwidth
  - O(10TB/s)→O(100TB/s)
- Parallel Filesystem NVRAM/SSD/Disk
- Site-wide shared warm storage
  - SAN limited  $O(1TB/s) \rightarrow O(10TB/s)$
- Indexed data



# On-cluster (hot) storage requirements

Scalable capacity

O(10x) system DRAM

Scalable throughput

- Significant fraction of fabric bandwidth
- Significant fraction of fabric injection rate

Data integrity & consistency

- Tunable resilience/availability
- No silent failures
- Safe overwrite

### Minimal usage constraints

- Global namespaces
  - System namespace shared across jobs
    - Connected workflows
  - Object namespaces shared across processes

     Encapsulating entire simulation datasets
- Fine grained, random, massively concurrent
   Minimal interference
- Data movement by unrelated workflows

Security

Authorized/authenticated access



# **Global Namespaces**

#### Containers

- Shared System Namespace
  - "Where's my stuff"
- Private Namespaces
  - "My stuff"
  - Entire simulation datasets

### **Multiple Schemas**

POSIX\*

. . .

- Shared (system) & Private (legacy datasets)
- No discontinuity for application developers
- Scientific: HDF5\*, ADIOS\*, SciDB\*, ...
- Big Data: HDFS\*, Spark\*, Graph Analytics,





# Distributed Application Object Storage

### Exascale I/O stack

- Extreme scalability, ultra fine grain
- Integrity, availability, resilience
- Unified model over all storage tiers site-wide
   Multiple Top Level APIs
- Domain-specific APIs
- High-level data models
- **DAOS-CT**: Caching and Tiering
- Data migration over storage tiers
  - Guided by usage metadata
  - Driven by system resource manager

**DAOS-SR**: Sharding and Resilience

- Scaling throughput over storage nodes
- Fail-out resilience across storage nodes

DAOS-M: Persistent Memory Object Storage

- Ultra-low latency / fine grain I/O
- Fine-grain versioning & global consistency
- Location (latency & fault domain) aware



Tools

**Applications** 

(intel)

### **Transactions**

### Why

- Simplify application development
  - Safe update in-place
  - Guaranteed data model consistency
  - Concurrent producer/consumer workflows
- Support resilience schemas
  - Guaranteed consistency for redundant distributed data
- Support tiered storage
  - Preserve integrity on data migration

### How

Named Snapshots

Committed (immutable)

Lowest

Open

Epoch

Multi-version concurrency control

Consistent

Readers

- Snapshot consistency on read
- Maximize concurrency/asynchrony
- Process groups
  - Arbitrary numbers of collaborating processes

Highest

Committed

Epoch

Writers

&

Inconsistent Readers

- Arbitrary numbers of storage targets
- Leader commit/snap/migrate



# **DAOS-M Object Storage**

Multiple Independent Object Address Spaces

- Versioning Object PGAS
- Container = {container shards} + metadata
  - Container Shard = {objects}
    - Object = KV store or byte array
    - Sparsity exposed
  - Metadata = {shard list, commit state}
    - Minimal
    - Resilient (Replicated state machine)

#### Maximum concurrency

- Byte-granular MVCC
- Deferred integration of mutable data
- Writes eagerly accepted in arbitrary order
- Reads sample requested version snapshot

#### **Distributed Transactions**

- Prepare: Send updates tagged with version 't'
- Commit: Mark version 't' committed in container MD
  - Version 't' now immutable and globally consistent
- Abort: Discard version 't' updates everywhere

#### Low latency

- End-to-end OS bypass
- Persistent Memory server
- Userspace fabric drivers





# **DAOS-M latency sensitive server operations**

### Byte array objects

- Write (log data)
  - Allocate extent buffer in NVRAM
  - Copy immediate / RDMA READ remote
  - Insert into persistent extent.version index
- Read (data integration)
  - extent.version index traversal => gather descriptor
  - RDMA WRITE remote

Key-Value objects

Insert/remove/retrieve value into key.version table



# **Sharding & Resilience**

Multiple mixed schemas

Performance schemas

Scale IOPs & bandwidth

**Resilience schemas** 

- Data integrity
  - Checksums + data stored separately
- N-way replication
  - High performance for shared write objects
- Erasure codes
  - High efficiency for non-shared objects
- Asynchronous refactor, scrub & repair
  - Exploit immutability of committed data

#### Leverage DAOS-M

- Distributed consistency
- Sparsity

#### **Scaling Requirements**

- Onerous!
  - Aliasing of access & distribution patterns
  - Bulk synchronous workload == Amdahl's law vector
- Extreme object size dynamic range
  - "Megaliths" v. "grains of sand"
- Declustering
  - Rebalance on node addition
  - Distributed rebuild on node failure



### **Sharding & Resilience**

Algorithmic (O(0)) layout metadata

- Consistent hash randomizes placement
- Replicas placed adjacently
  - Hash must guarantee minimum separation of nodes in same fault domain
- Multiple hash rings for declustering

Explicit (O(n)) layout metadata

- Layout responsive to usage
  - Preserve locality / feed "hungriest mouth"
- O(0) structures used to store layout





# **Caching & Tiering**

#### Metadata

- Residence maps
  - Whole object maintained directly
  - Sub-object leverages lower layers
- Access patterns
  - Historical
  - Explicit notification by upper levels
  - Data "colouring"

#### Data migration

- Resharding between tiers
  - Maintain distributed object semantics
  - Maximize performance on subsequent access
  - Select appropriate resiliency schemas



"Near" storage

"Far" storage

#### **Explicit control**

- Persist & prestage APIs / JCL
- System resource manager driven migration
  - Rebalance & minimize interference

#### **Transparent caching**

- Write-back & demand cache
- Prefetch guided by usage metadata
- Residence maps



# Top level I/O APIs

### **POSIX Containers**

- POSIX namespace over DAOS-HA objects
  - Dynamically sharded directories & files
- Private POSIX namespaces
  - Library for parallel applications and middleware targeting POSIX
- System POSIX namespace
  - Parallel server exporting shared namespace

DAOS for application programmers

- Simplified APIs
- Distributed Persistent Memory



### High level HPC object databases

- Complex application datatypes & metadata
- HDF5 + derivatives / ADIOS / SciDB etc...

### **Big Data**

- HDFS compatibility layer
  - Hadoop ecosystem
- Spark / Graph Analytics etc...



### **NVRAM Storage Revolution**

Cost-effective storage & fabric integration

- Challenge: Extreme scale-out
  - Amdahl's law
  - Fault Tolerance
- Reward: Storage @ full fabric bandwidth
  - O(1000) increase in data velocity

Byte-granular data access @ uS latency

- Challenge: Deliver benefit to applications
  - Software overhead of conventional storage & communications stacks
- Reward: Ultra fine-grain access
  - Remove constraints on applications
  - Enable new programming models



