

Pelican: A Building Block for Exascale Cold Data Storage

Austin Donnelly Microsoft Research

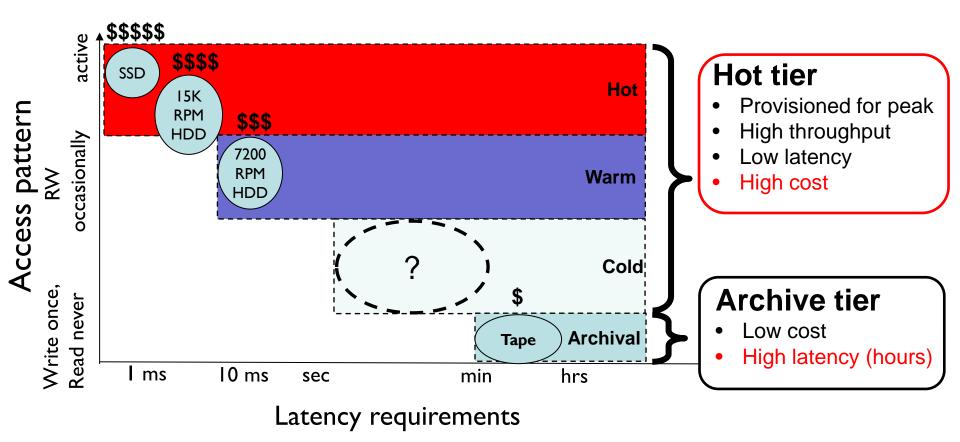
and: Shobana Balakrishnan, Richard Black, Adam Glass, Dave Harper, Sergey Legtchenko, Aaron Ogus, Eric Peterson, Ant Rowstron

Outline

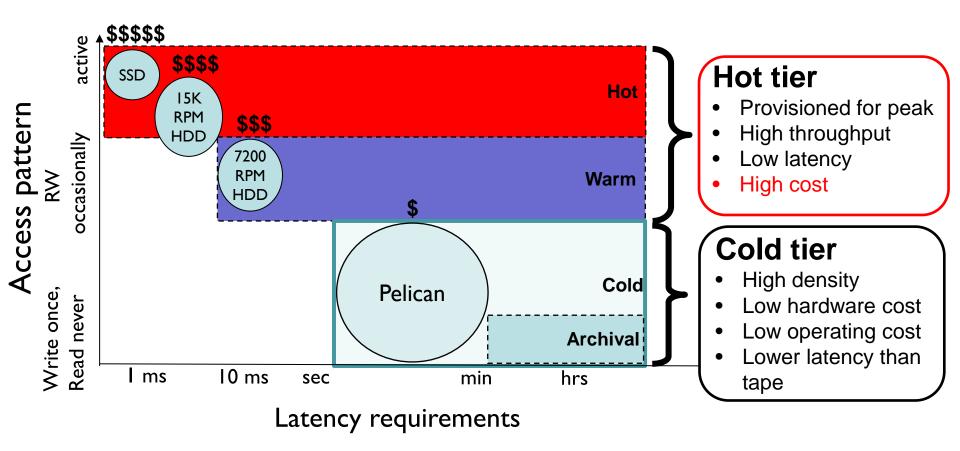
- Background
- □ Pelican co-design
- Research challenges
- Demo
- □ Performance results



Background: Cold Data in the Cloud



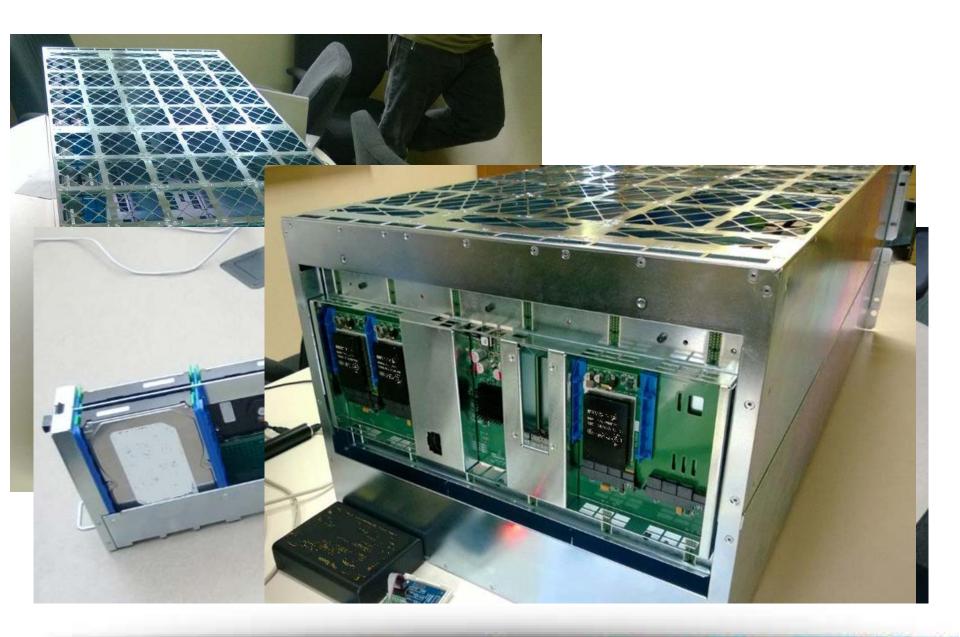
Background: Cold Data in the Cloud



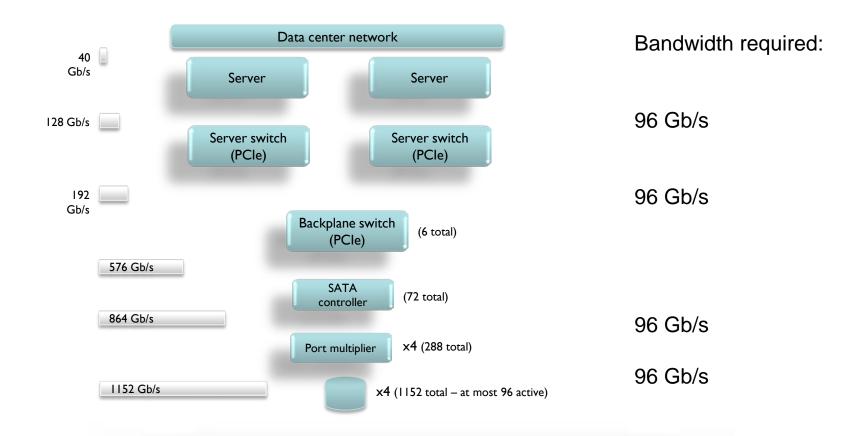
Pelican: Rack-scale Co-design

- Hardware & software co-designed:
 - Power, Cooling, Mechanical, HDDs & Software.
 - Trade latency for lower cost.
- Massive density, low per-drive overhead.
- □ 1152 3.5" HDDs per 52U.
- 2 servers, PCIe bus stretched rack-wide.
 - 4x 10G links out of rack.
- Only 8% of disks can spin.





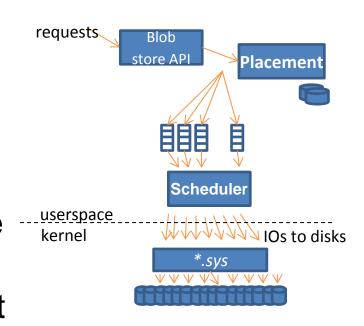
Interconnect Details



Research Challenges

- Not enough cooling, power, or bandwidth.
- How do we manage theses resource limits?
- Which disks to use for data? The data layout problem.
- How to schedule requests to get good performance.

["Pelican: A building block for exascale cold data storage", OSDI 2014]



Resource use

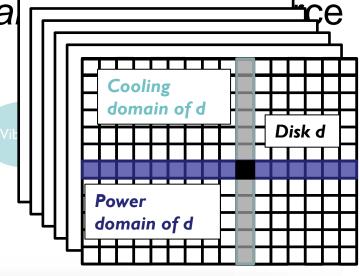
- □ Traditional systems:
 - Any disk can be active at any time.
- □ Pelican:

Rack: 3D array of disks

Disk is part of a doma

□ Domains, limits:

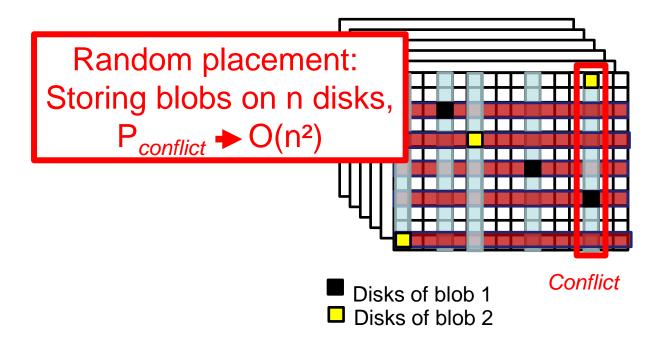
Power 2 of 16 Cooling 1 of 12



Data placement

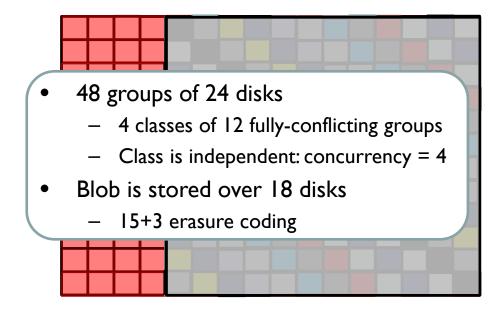
- Blob erasure-encoded on a set of concurrently active disks
- In traditional systems:
 - Any two sets can be active
 - No impact on concurrency
- □ In Pelican:
 - Sets can conflict in resource requirements
 - Conflicting sets cannot be concurrently active
- Challenge: form sets to minimize P_{conflict}

Data placement: random



Data placement: Pelican

- Intuition: concentrate conflicts over a few sets of disks.
- Store blob in one group
 - \square P_{conflict} \rightarrow O(n)
- Groups encapsulate constraints:
 - Unit of IO scheduling
 - No constraints managed at runtime.



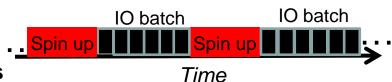
Schematic side-view of the rack

IO Scheduling: "spin up is the new seek"

Four independent schedulers

Each scheduler: 12 groups, only one can be active

- Naïve scheduler: FIFO
 - Avg. group activation time: 14.2 sec ·····
 - High probability of spinup after each request
 - Time is spent doing spinups!
- Pelican scheduler: Request batching
 - Limit on maximum re-ordering
 - Trade-off between throughput and fairness
 - Weighted fair-share between client and rebuild traffic



Time

0

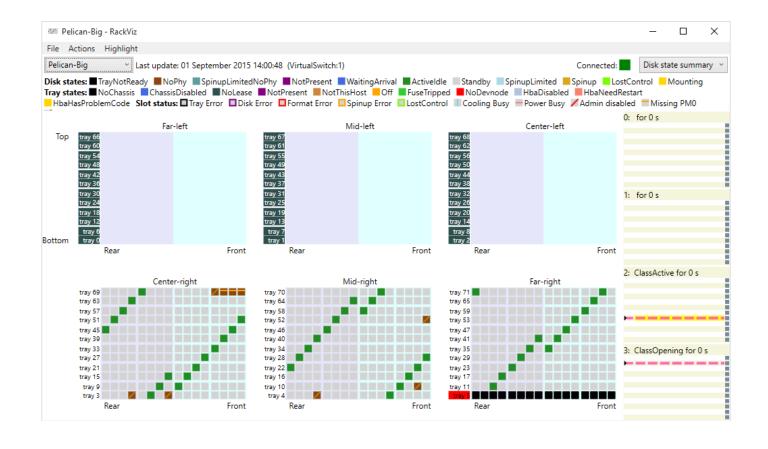


Outline

- Background
- □ Pelican co-design
- □ Research challenges:
 - Data placement: constraint-aware
 - Scheduler: batching to amortize spinups
- Demo
- Performance results



Demo



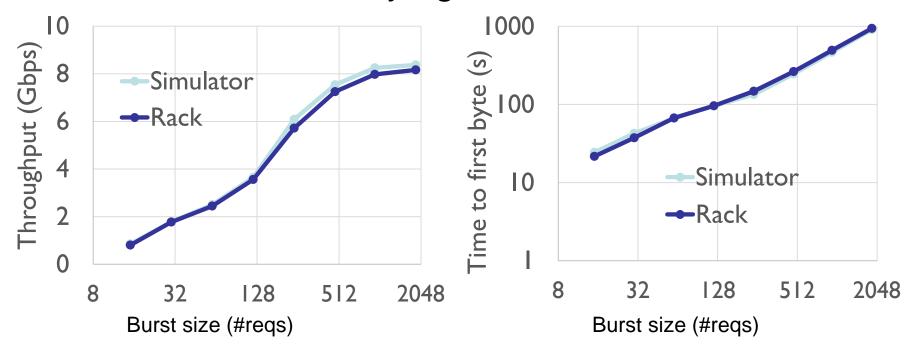
Performance

- Compare Pelican vs. all disks active (FP).
- Cross-validate simulator.
- Metrics:
 - Throughput
 - Latency (time to first byte)
 - Power consumption
- Open loop workload:
 - Poisson arrivals
 - Read requests, 1GB blobs



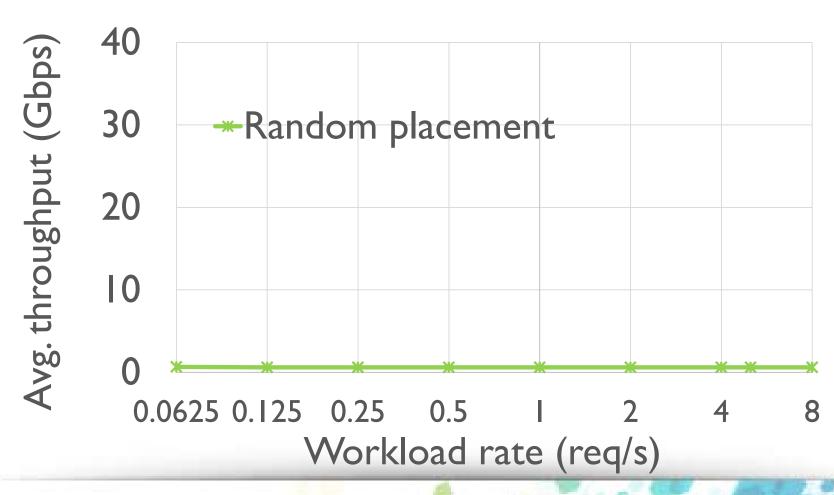
First step: simulator cross-validation

Burst workload, varying burst size



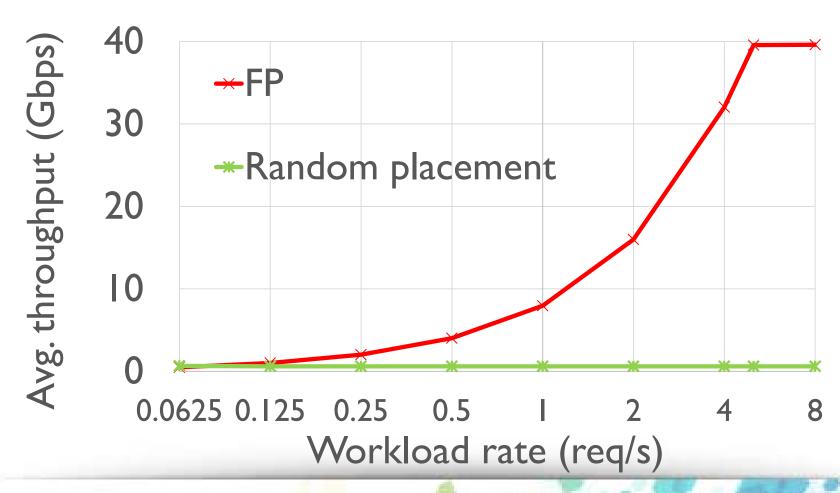


Rack throughput



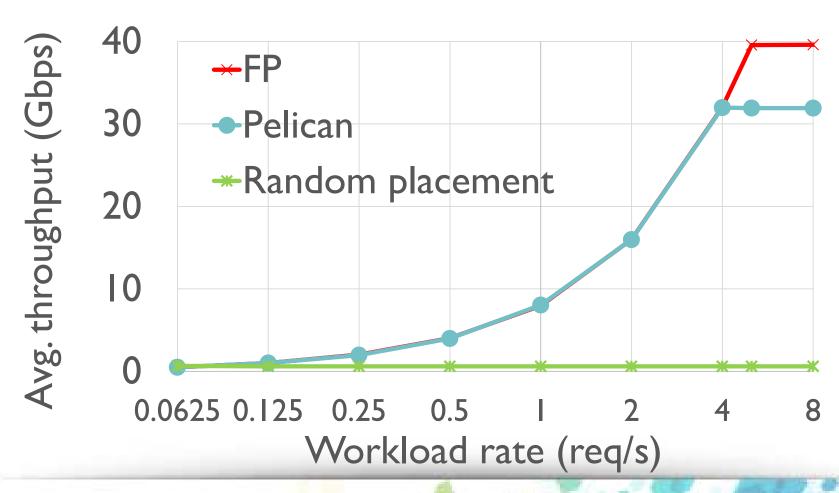


Rack throughput



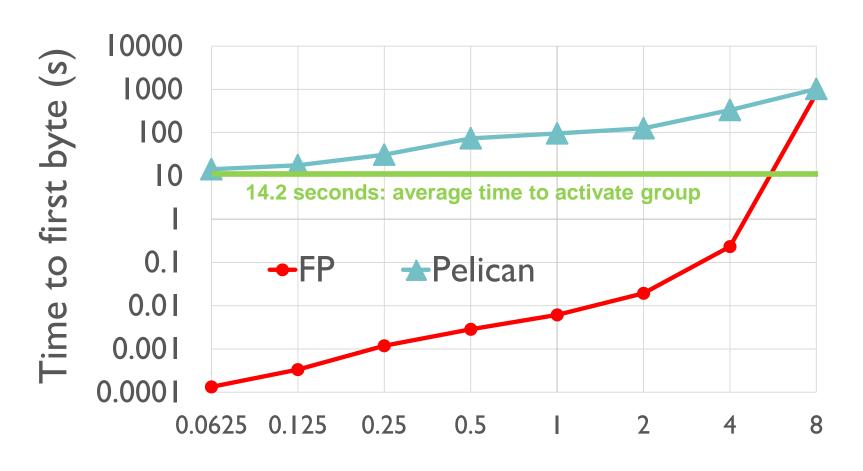


Rack throughput



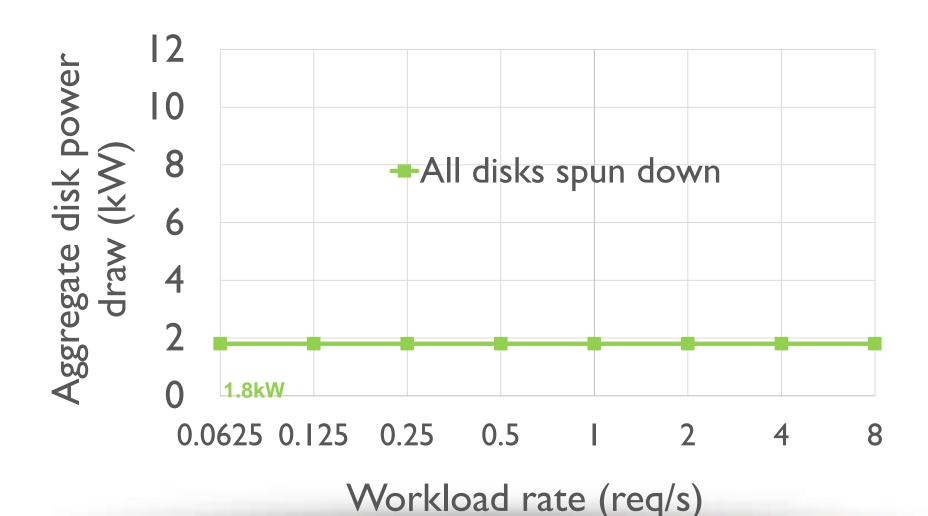


Time to first byte



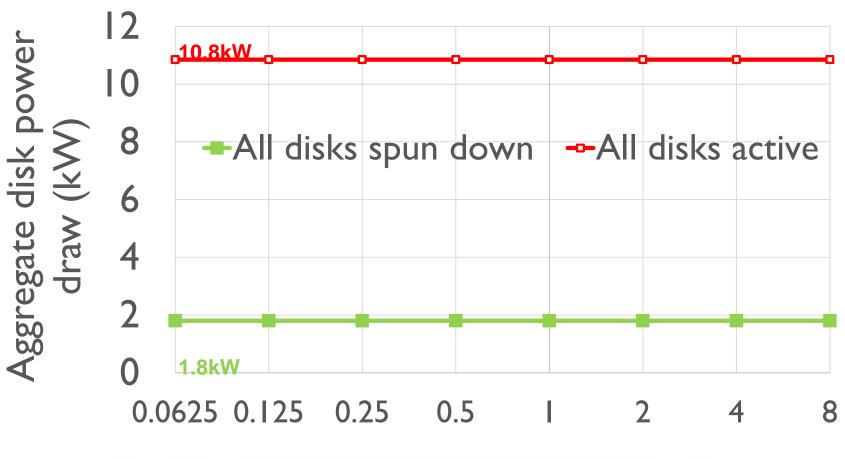


Power consumption



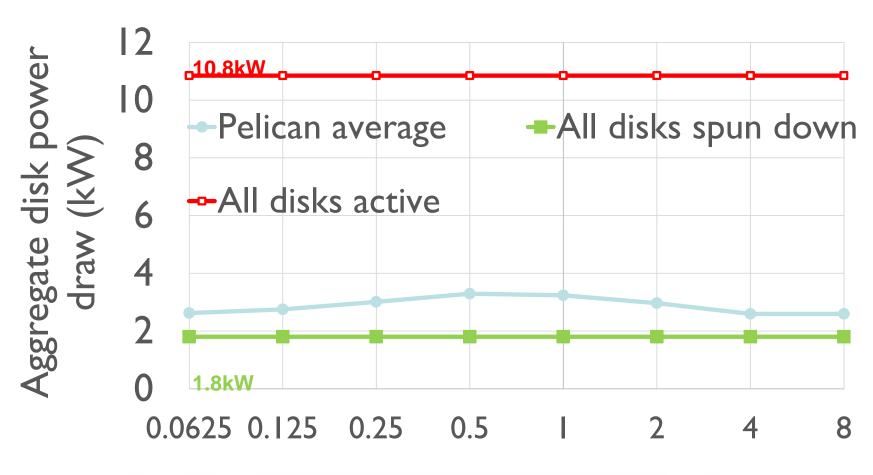


Power consumption



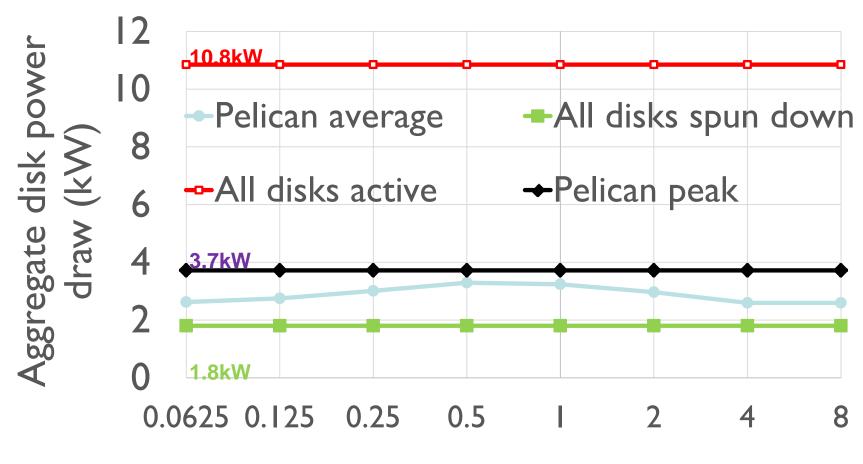


Power consumption





Power consumption: 3x lower peak

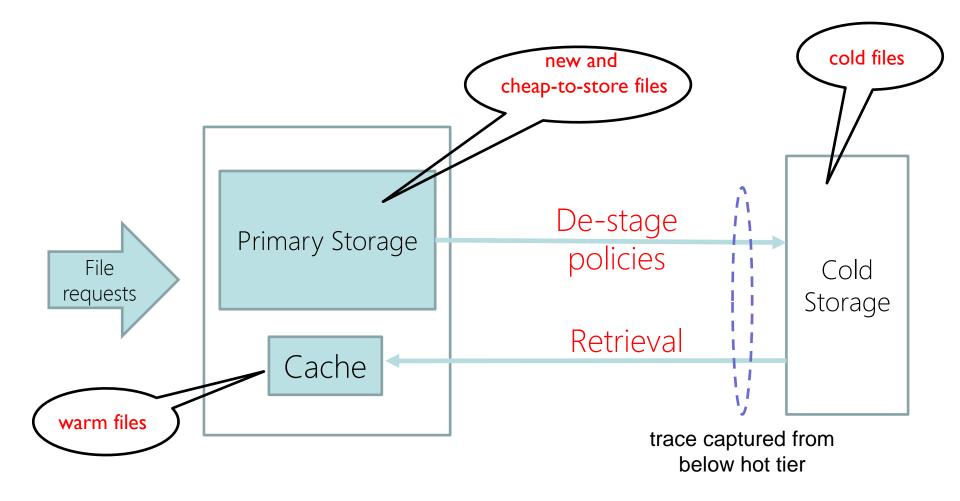




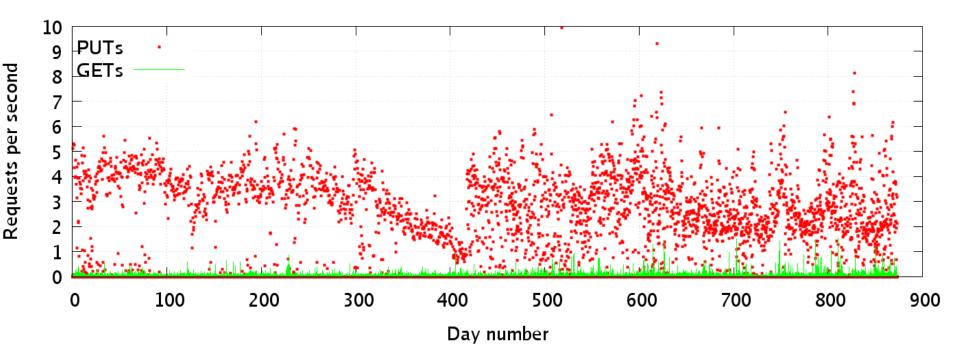
Trace replay

- European Centre for Medium-range Weather Forecasts [FAST 2015]
 - ECFS trace is every request for 2.4 years.
 - Run through a tiering simulator

Tiering model



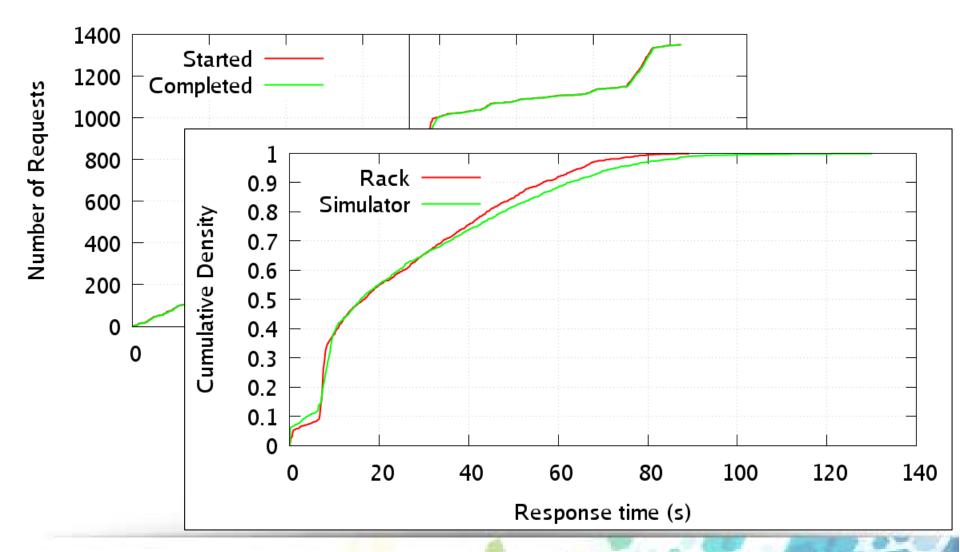
Requests per second, over 2.4 years



- We replay two 2-hour segments:
 - □ G1: highest response time
 - □ G2: deepest queues

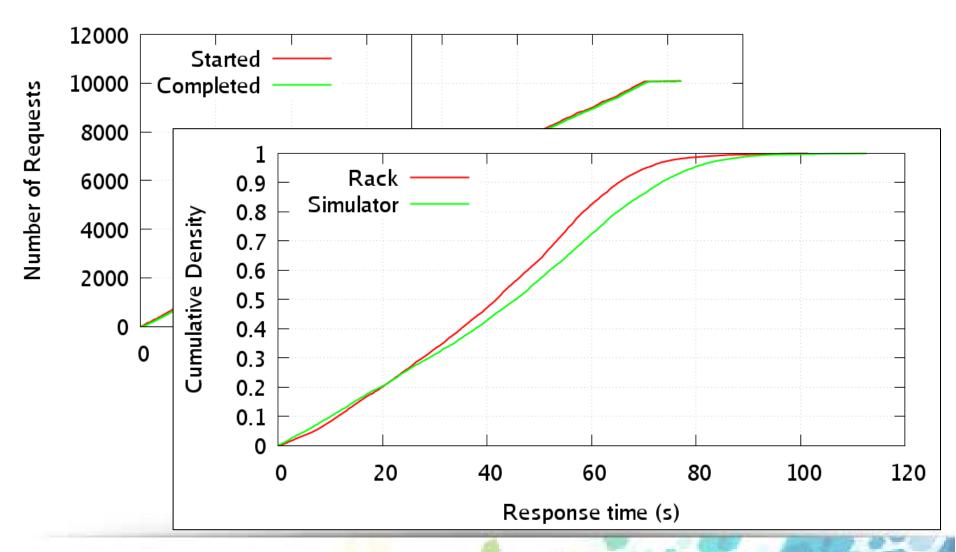


G1: Highest response time





G2: Deepest queues





War stories

- Booting a system with 1152 disks
 - BIOS changes needed
- Object store vs. File system
- □ Data model for system:
 - Serial numbers on all FRUs
 - Disks, Volumes, Media



Thank you!

Questions?

