#### Skylight – A Window on Shingled Disk Operation

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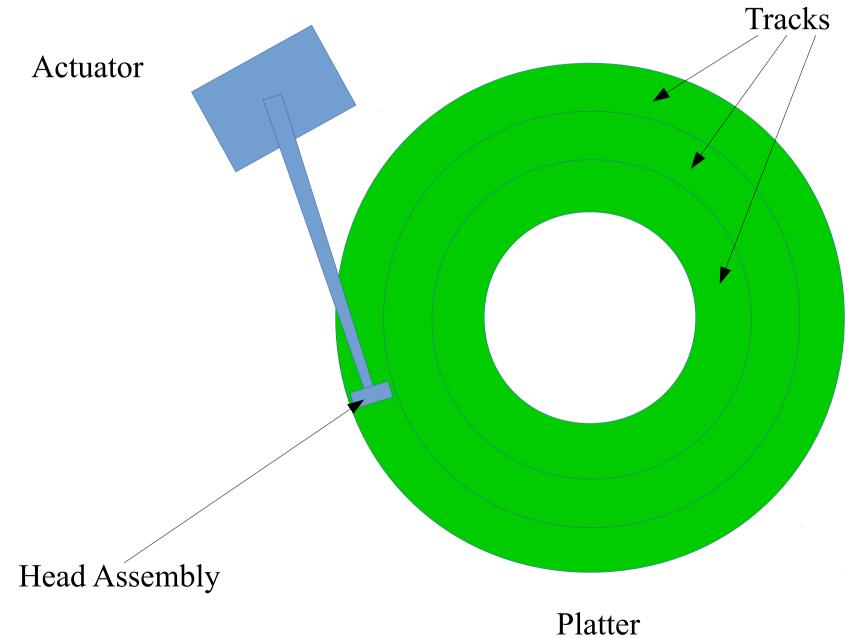




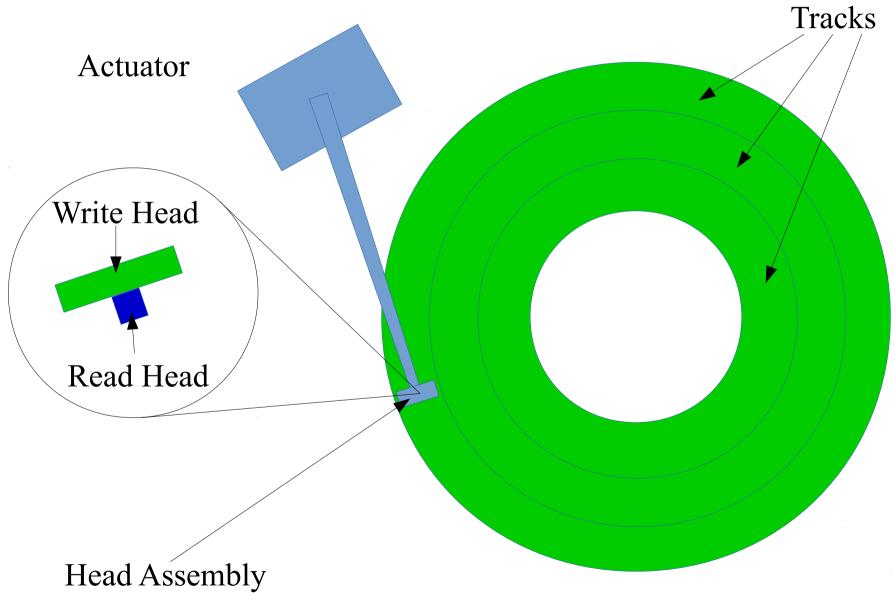
# What is Shingled Magnetic Recording (SMR)?

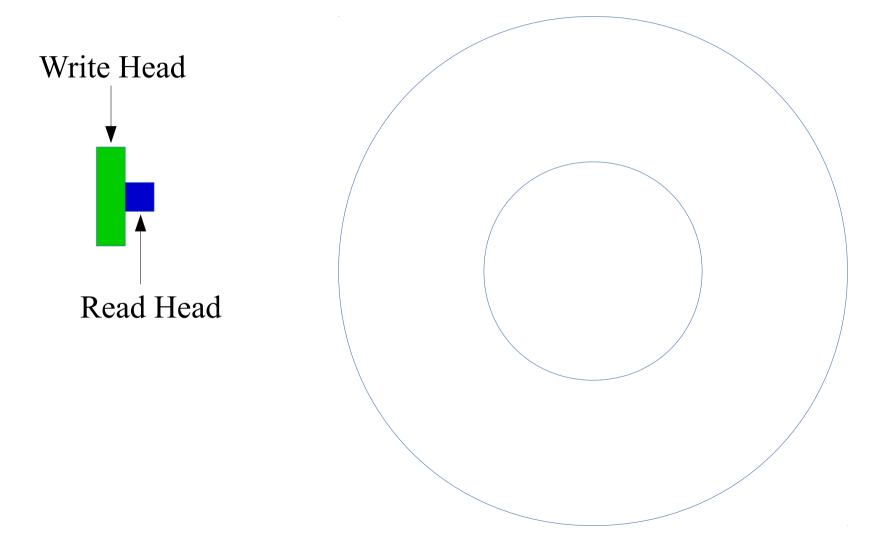
- A new way of recording tracks on the disk platter.
- Evolutionary uses existing infrastructure.
- Fits more tracks onto platter  $\rightarrow$  increases capacity.
- Disallows random writes  $\rightarrow$  increases complexity.

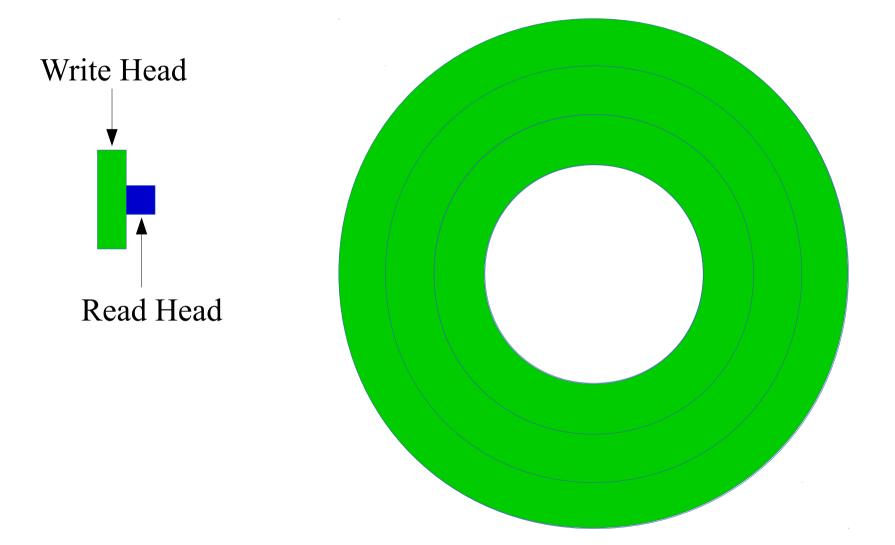
#### Disk Drive Internals

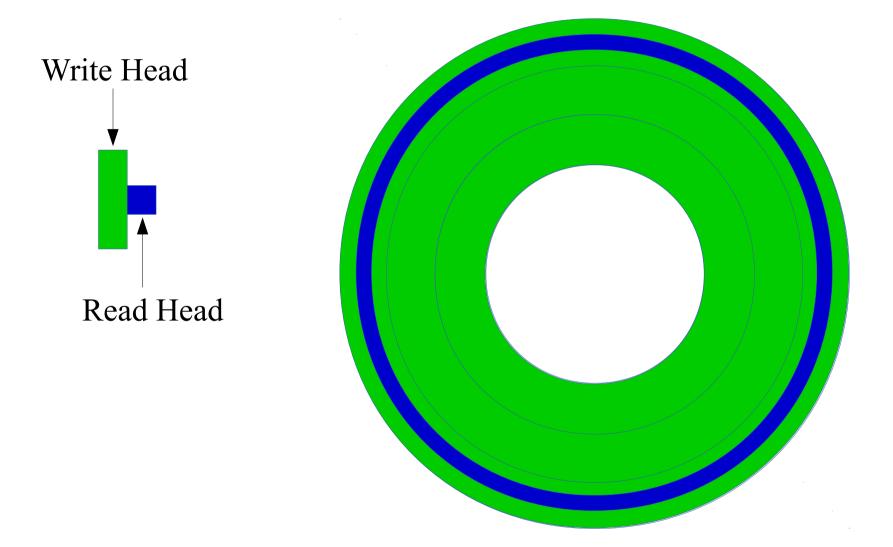


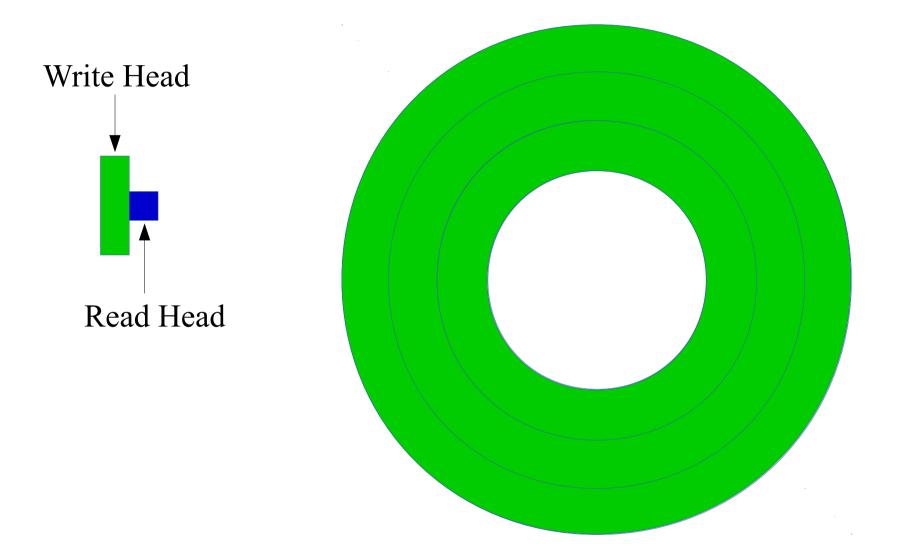
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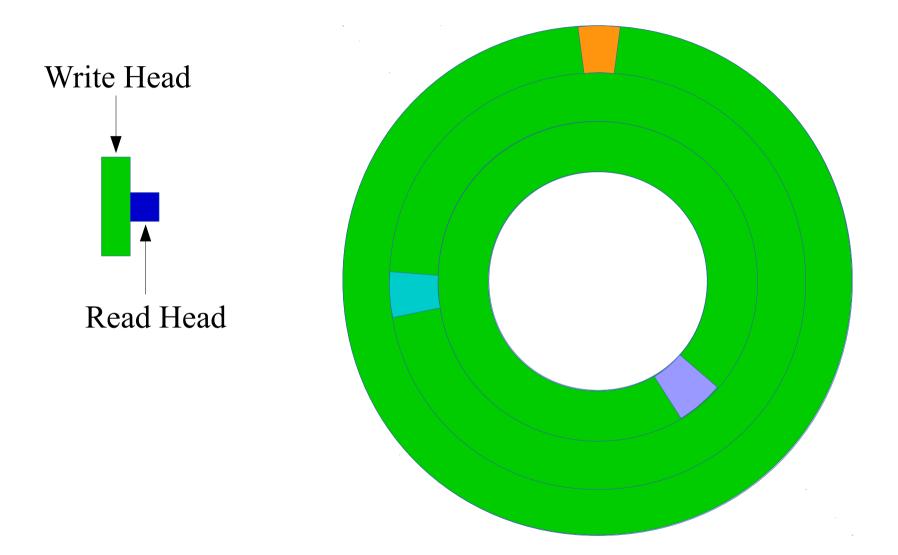


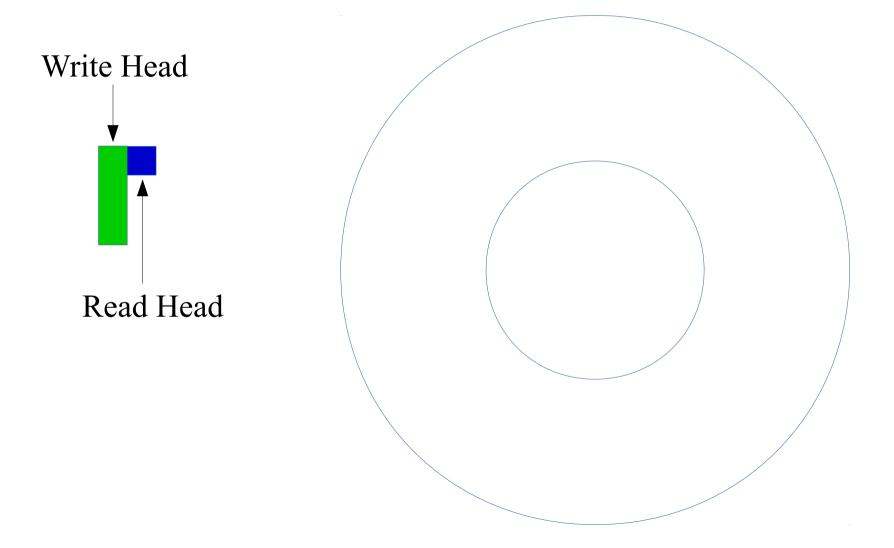


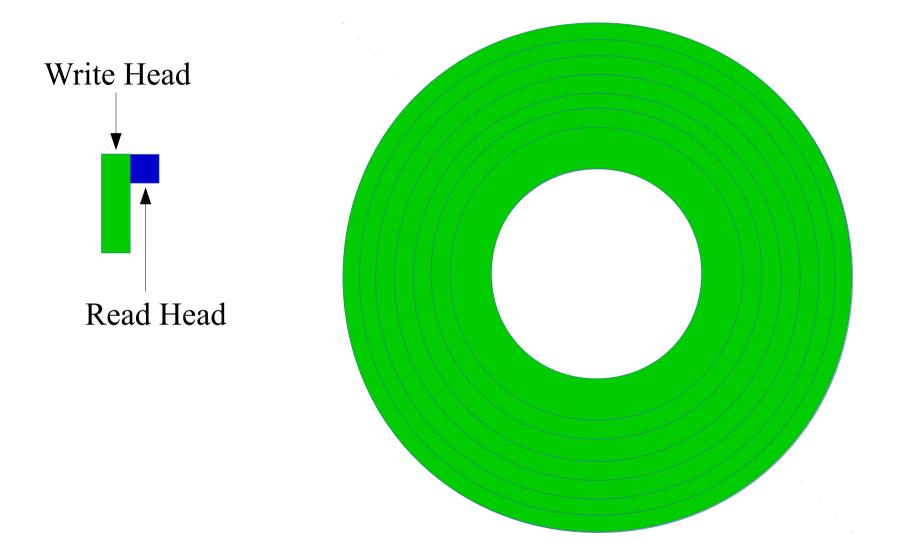


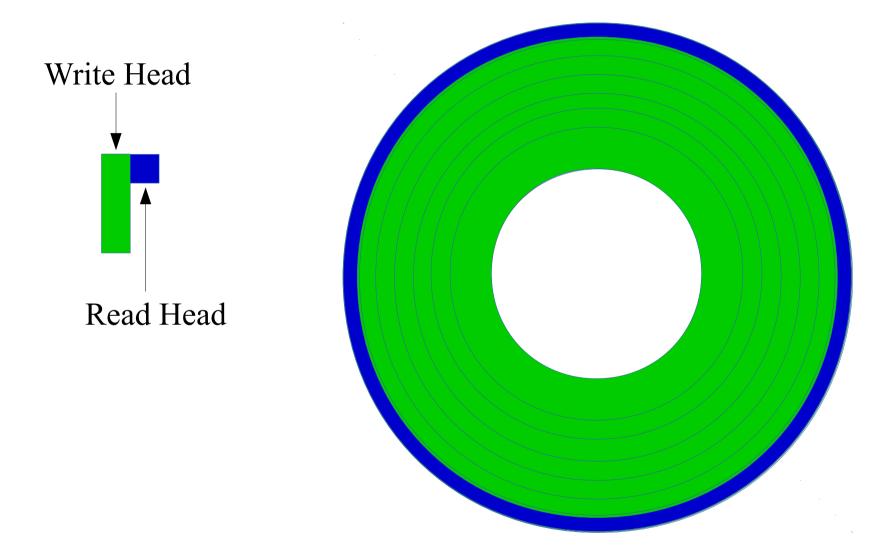


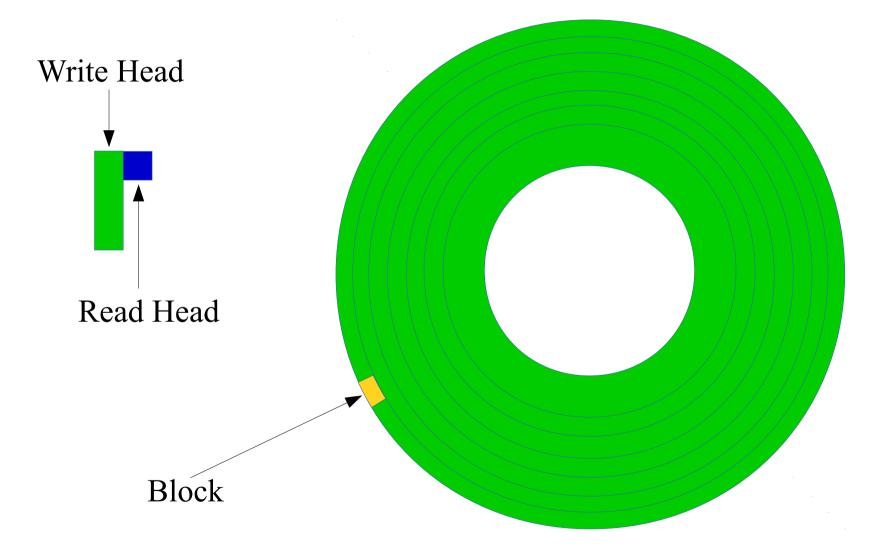


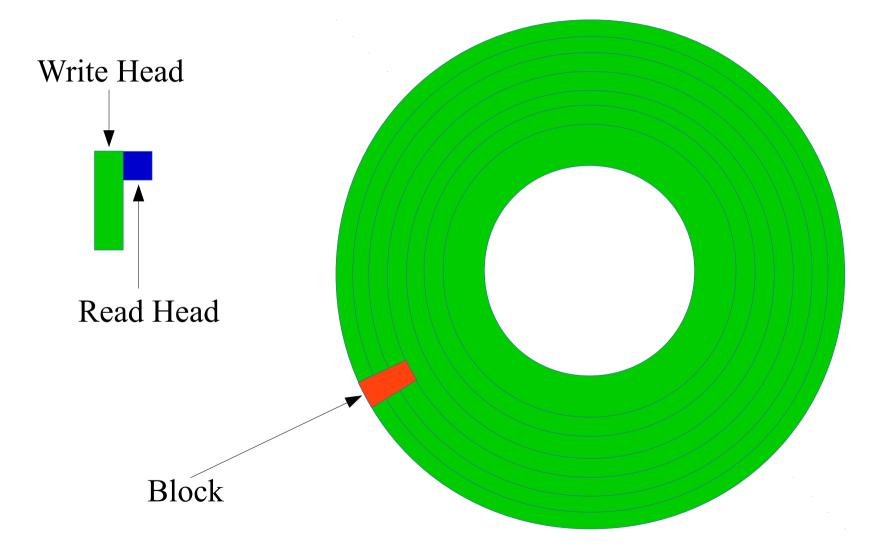


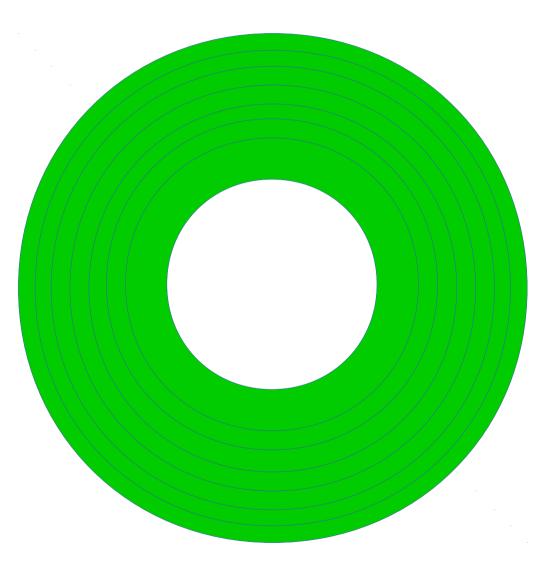




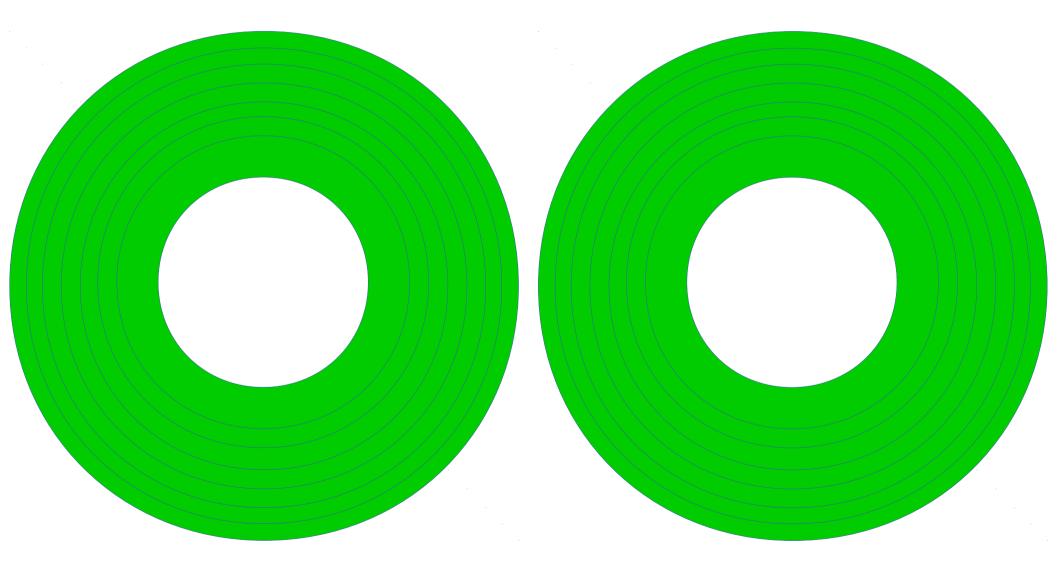


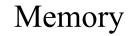


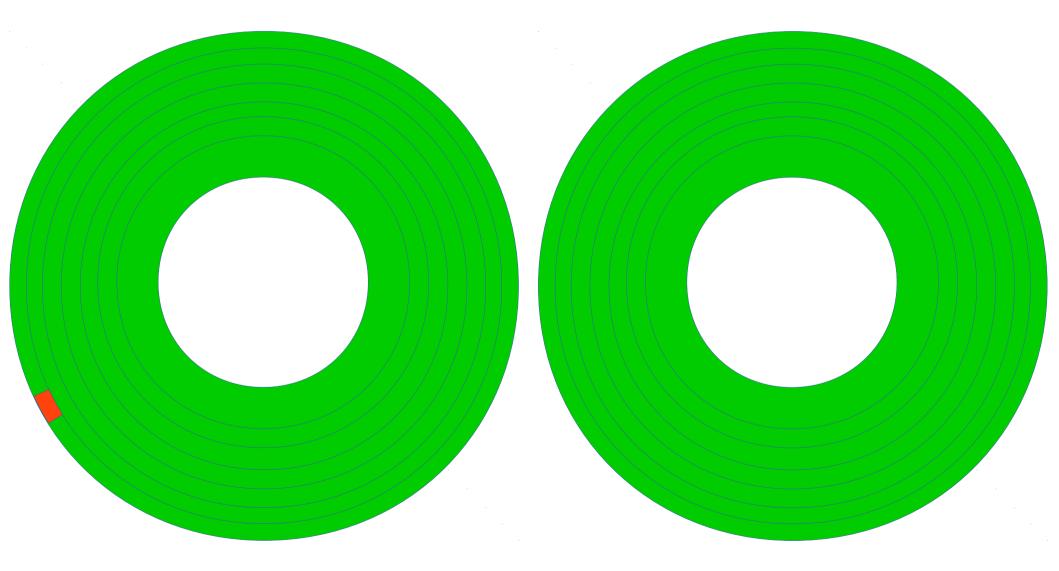


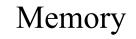


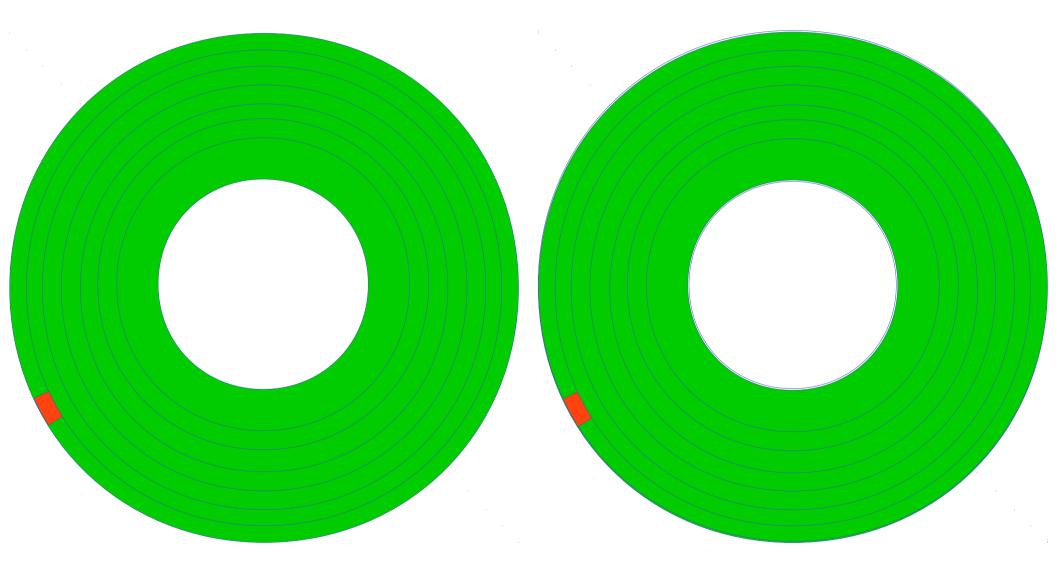
Memory

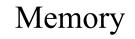


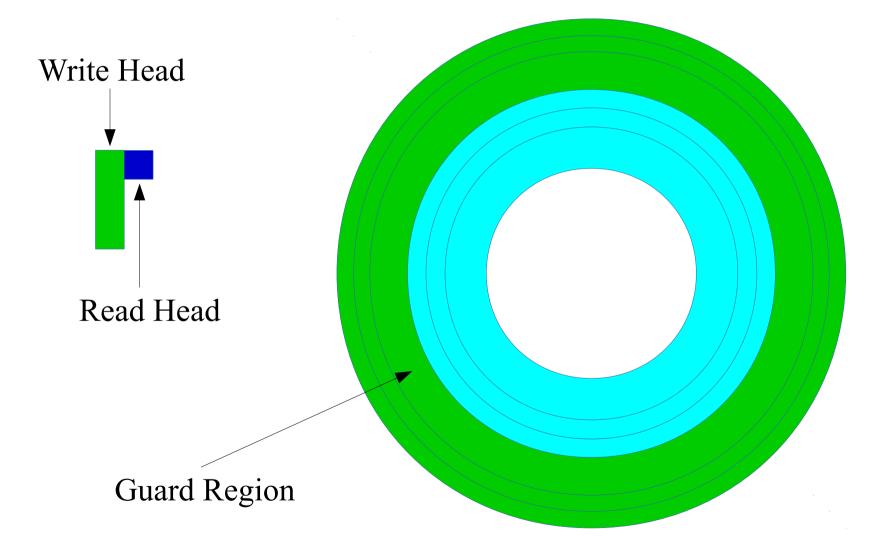






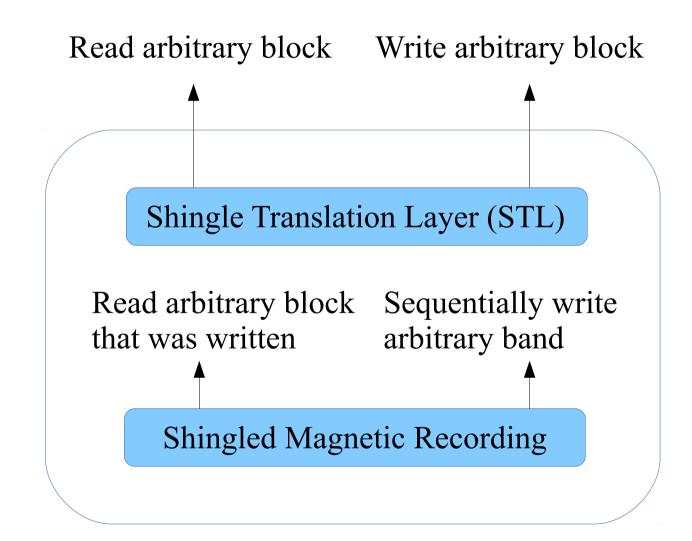


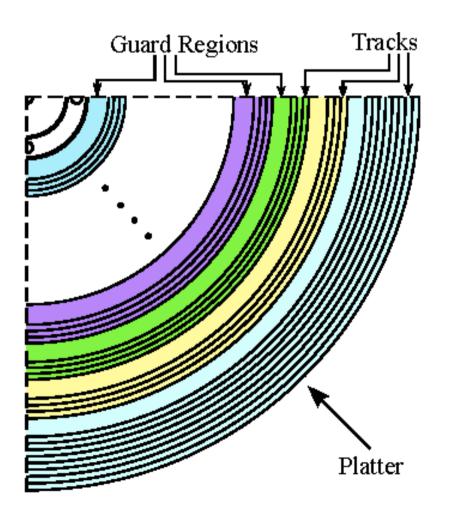




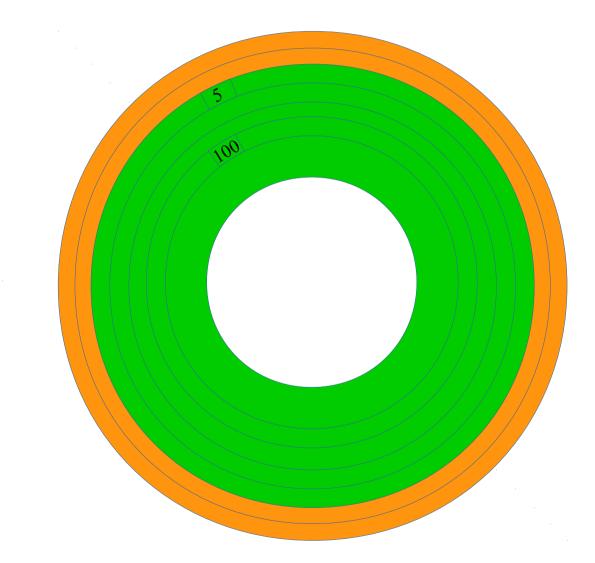
# SMR Drive Implementations

- Host-Managed
  - Reports band to host.
  - Bands must be written sequentially.
  - Random writes or reads before writes will fail.
- Host-Aware
  - Reports band to host.
  - Also handles random writes backward compatible.
- Drive-Managed
  - Hides SMR details.
  - Drop-in replacement for existing drives.

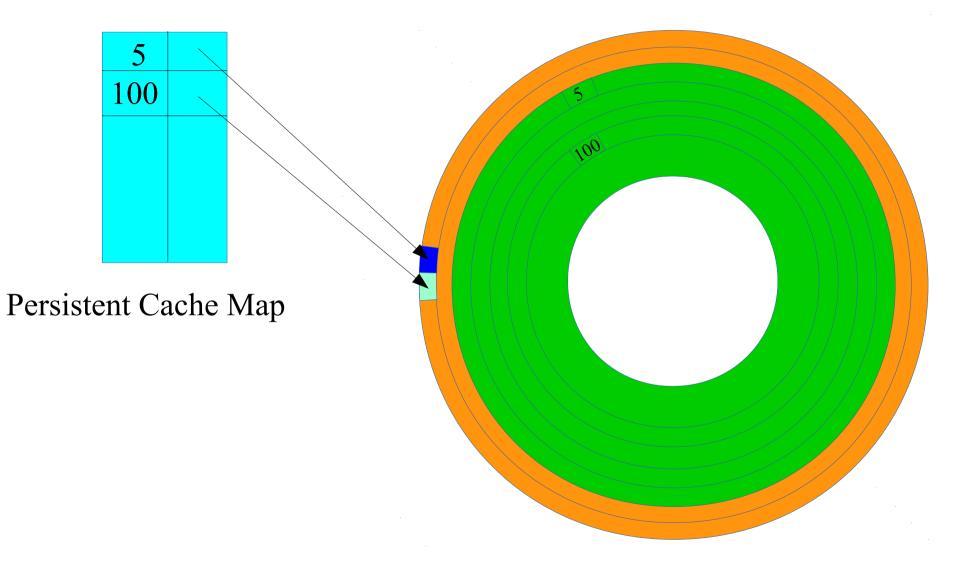




- Small region of disk, called persistent cache, used for staging random writes.
- Other non-volatile memory like flash can also be used for persistent cache.
- Disk is mapped at band granularity; persistent cache uses extent mapping.

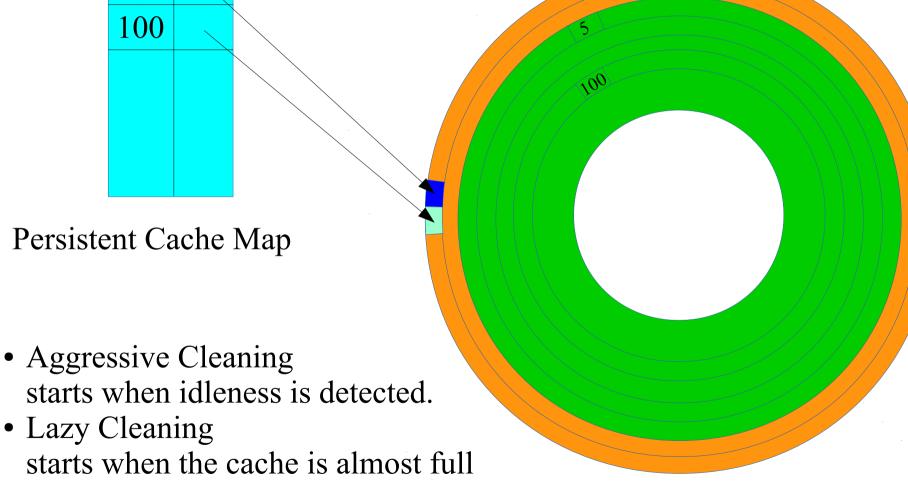


Bands are shown in green. Persistent Cache is shown in orange.



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# Outline

- Introduction to SMR
- Characterization goals and test setup
- Test results

#### Characterization Goals

- Drive Type
- Persistent Cache Type
- Cache Location and Layout
- Cache Size
- Cache Map Size
- Band Size

- Block Mapping
- Cleaning Type
- Cleaning Algorithm
- Band Cleaning Time
- Zone Structure
- Shingling Direction

# Skylight Components

- Software part:
  - Launch crafted I/O operations using fio.
  - Disable kernel read-ahead, drive look-ahead, on-board volatile cache.
  - Use latency to infer drive properties.
- Hardware part:
  - Install a transparent window on the drive.
  - Track the head movements using a high-speed camera.
  - Convert movements to head position graphs.



# **Emulation Strategy**

• STLs from the literature implemented as Linux device-mapper targets.





Drive-Managed SMR with persistent disk cache

Drive-Managed SMR with persistent flash cache

# Tested Drives

• Emulated Drives

Drive Name	Cache Type (Size)	Cache Location	Band Size	Capacity
Emulated-SMR-1	Disk (37.2 GB)	Single at ID	40 MiB	3.9 TB
Emulated-SMR-2	Flash (9.9 GB)	N/A	25 MiB	3.9 TB
Emulated-SMR-3	Disk (37.2 GB)	Multiple	20 MiB	3.9 TB

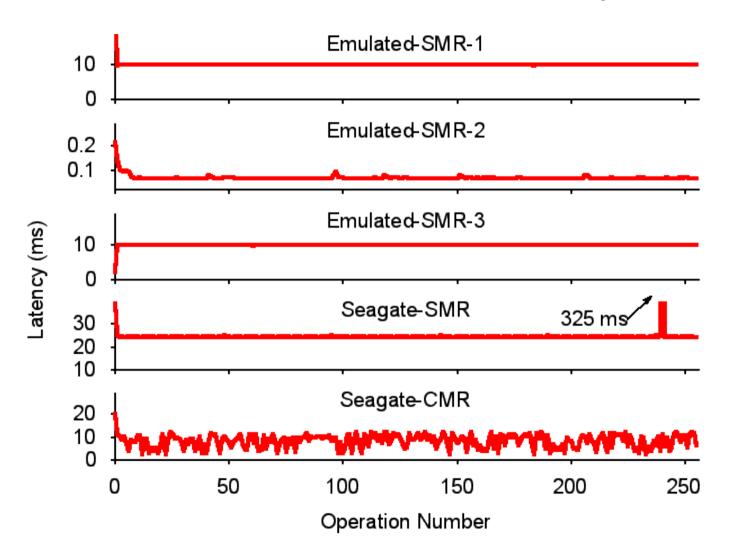
- All were emulated using a 4TB conventional Seagate drive.
- Real Drives
  - 5TB and 8TB Seagate drive-managed SMR drives.
  - We only show 5TB results labeled as Seagate-CMR.
- All disk drives are 5900RPM =>  $\sim 10$  ms rotation time.
- Write cache, read-ahead both disabled

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Test 1: Discovering the drive type and the persistent cache type

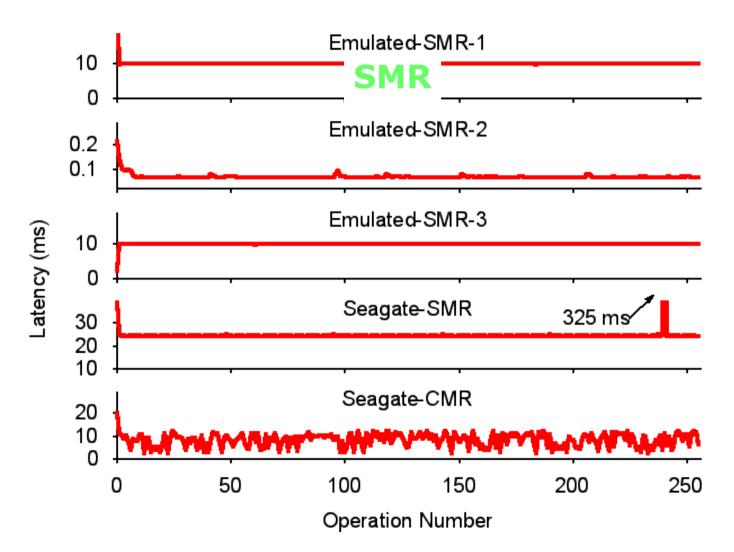
- Test exploits unusual random write behavior in SMR drives.
- Write blocks in the first 1GiB in random order.
- If latency is fixed then the drive is SMR, otherwise it is a conventional magnetic recording (CMR).
- Sub-millisecond latency indicates a drive with a persistent flash cache.



Y-axis varies in each graph.

- Conventional drive (Seagate-CMR) stands out from the rest. Emulated drive with persistent flash cache has sub-ms latency.

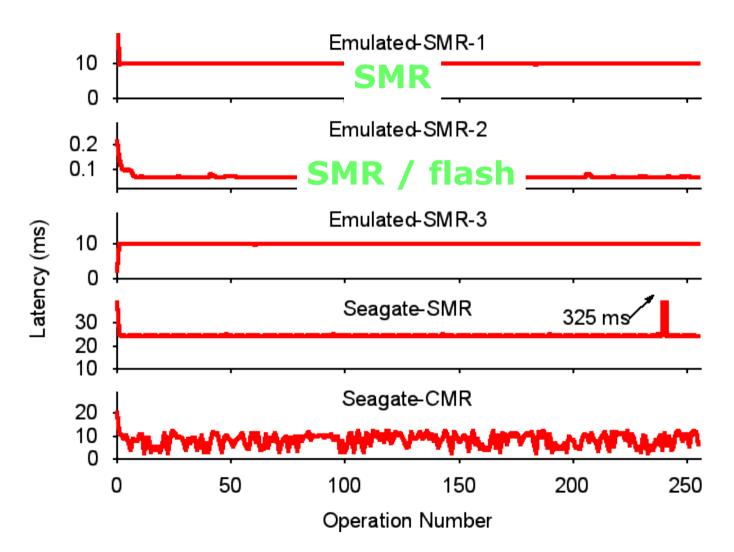
Latency is high for the real SMR drive.



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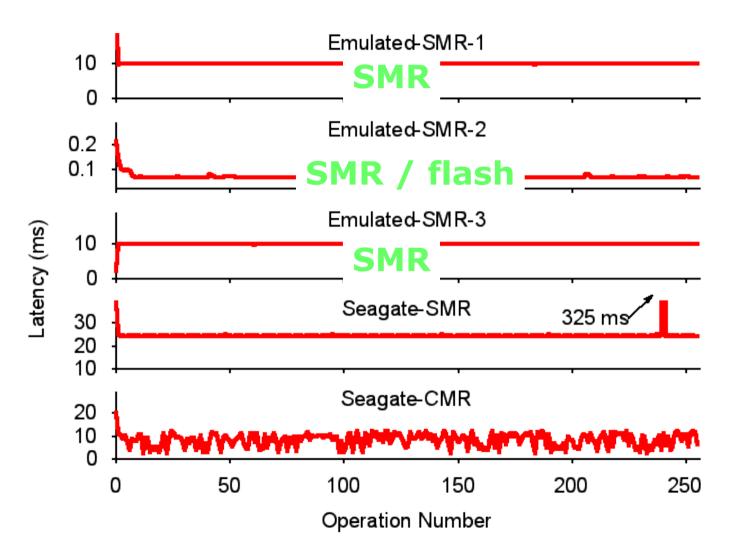
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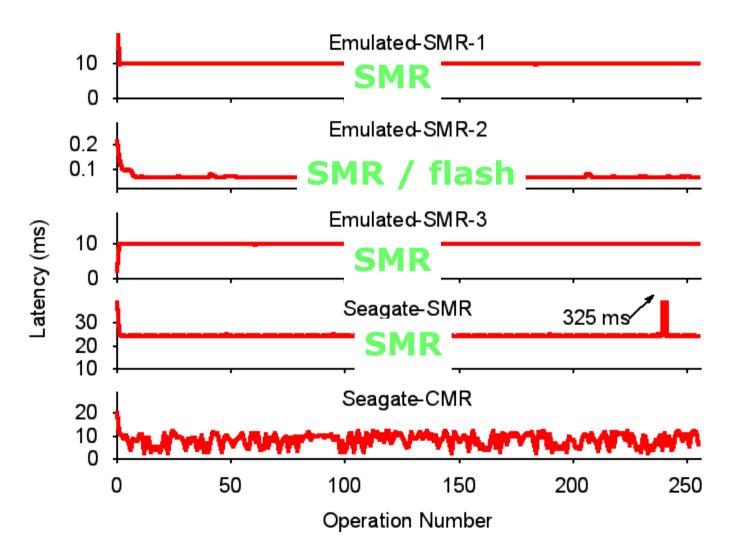
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#### Random Write latency

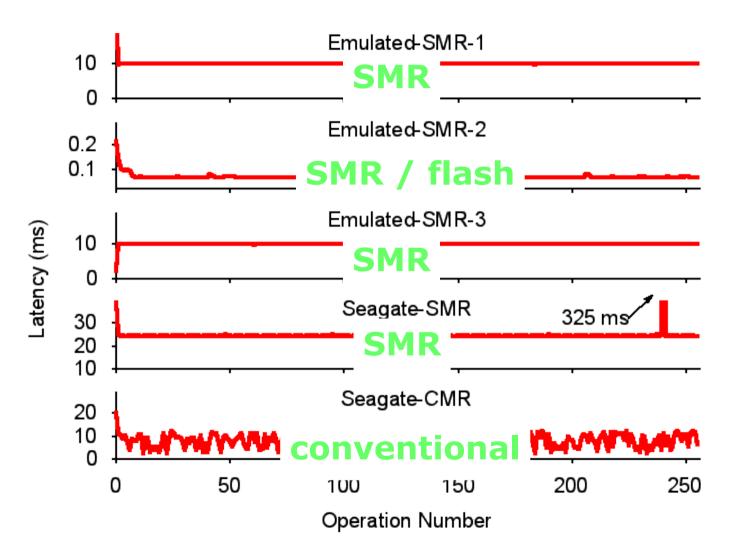


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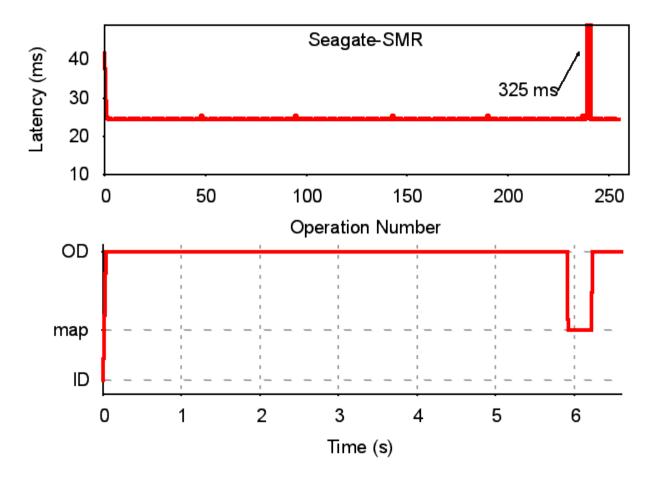


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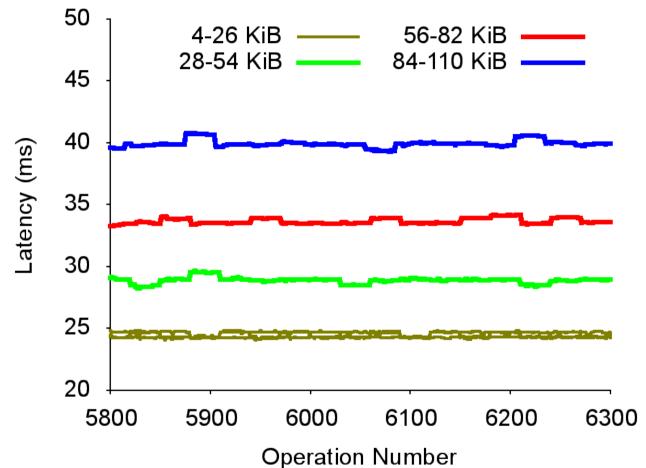
Latency is high for the real SMR drive.

## Random Write Latency + Head Position



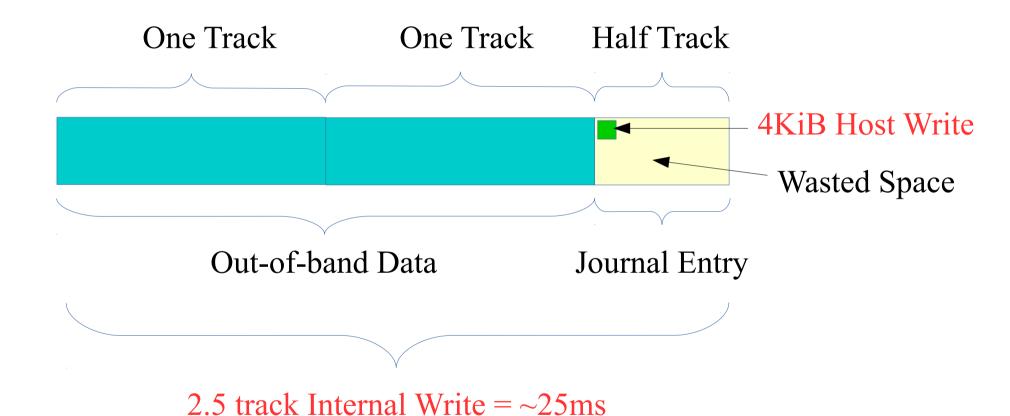
- There is a persistent cache at the outer diameter (OD).
- Writes are (likely) piggy backed with out-of-band data.
- There is (likely) a persistent map stored at the middle diameter.

# Random Writes with Max Queue Depth



- Different write sizes produce equal latencies.
- Latency increases in ~5ms jumps.
- Given  $\sim 10$ ms rotation time,  $\sim 5$ ms is  $\sim$  half-track increase in write size.

#### Host Write vs Internal Write

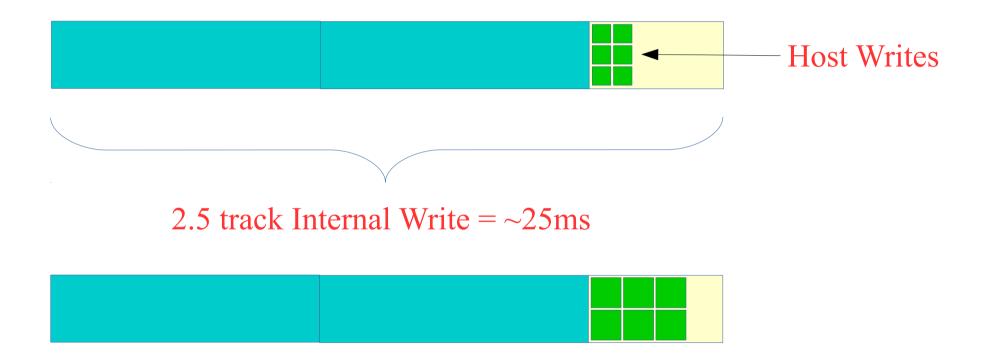


## Journal Entries with Quantized Sizes

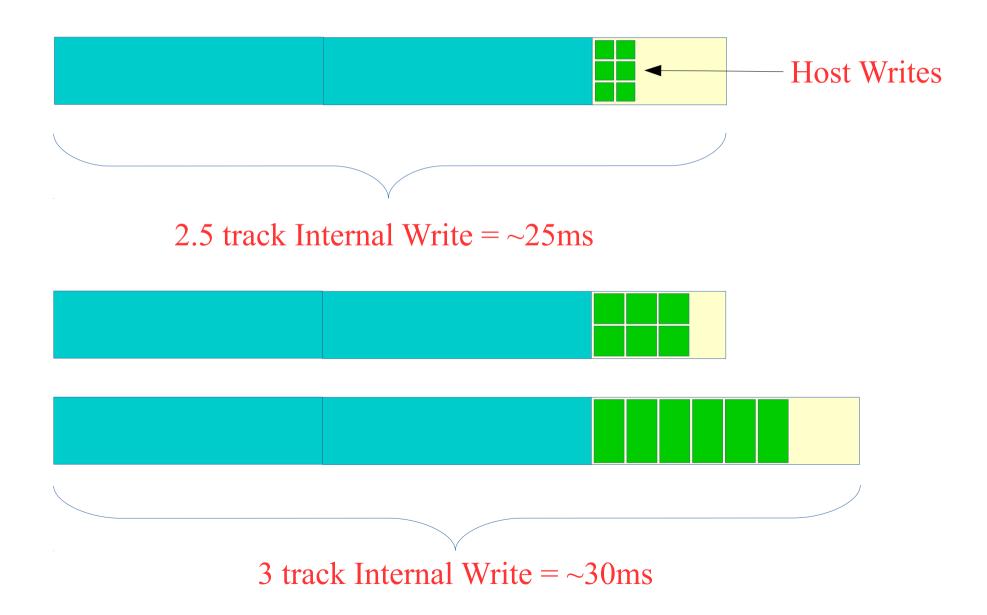


2.5 track Internal Write =  $\sim 25$ ms

## Journal Entries with Quantized Sizes



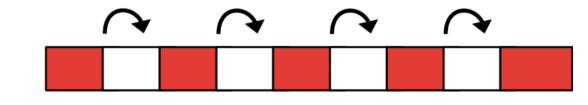
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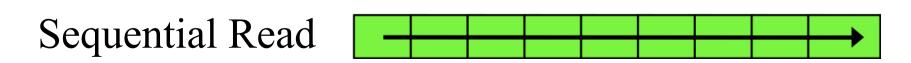


# Test 2: Discovering Disk Cache Location and Structure

- Test exploits a phenomenon called "fragmented reads".
- Fragmented read: during sequential read, seek to the persistent cache and back to read an updated block.
- Force fragmented reads at different offsets to infer persistent cache location based on seek time.

Skip Write



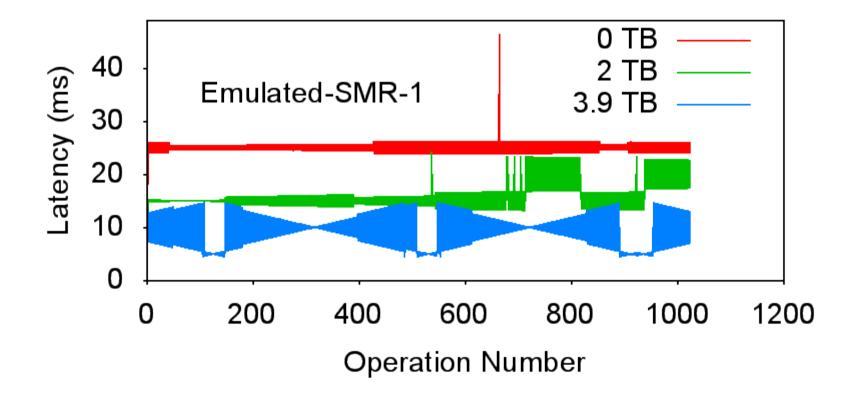


## Fragmented Read at 5TB Offset



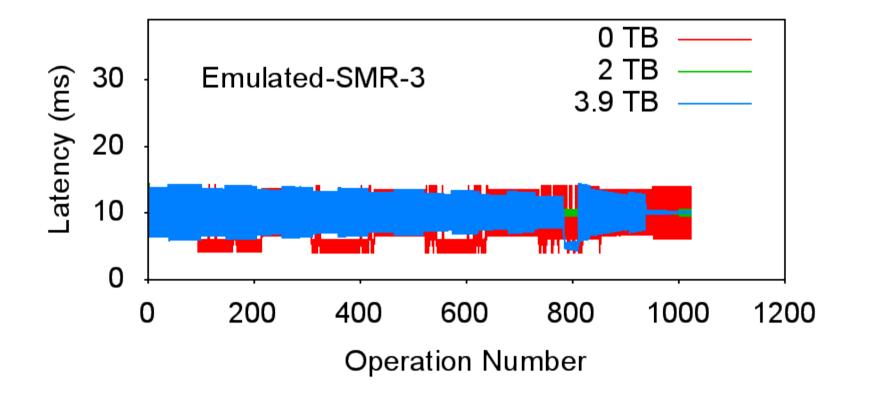
- Head seeks back and forth between a track and persistent cache.
- Persistent Cache is at OD, therefore, 5TB offset is at ID.
- Block numbering convention starts at OD proceeds towards ID.

## Fragmented Read Latency at Low, Middle, and High Offsets



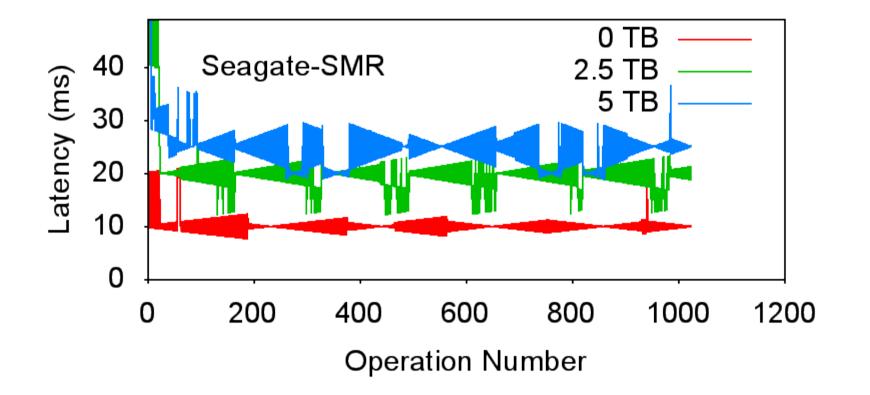
• Average latency high at low offset => cache at ID..

## Fragmented Read Latency at Low, Middle, and High Offsets



• Average latency is roughly fixed => distributed cache.

## Fragmented Read Latency at Low, Middle, and High Offsets



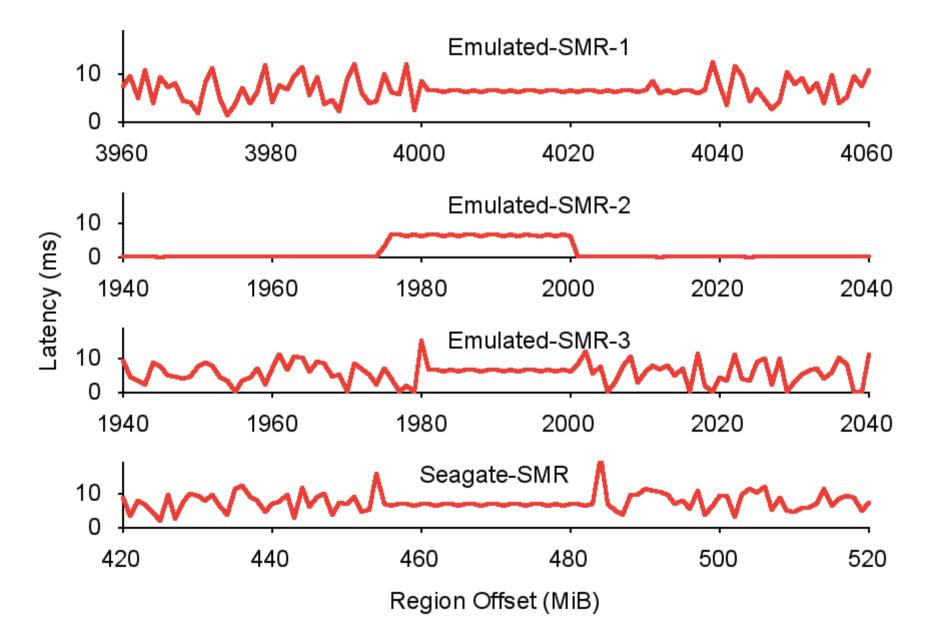
• Average latency is high at high offset => cache at OD.

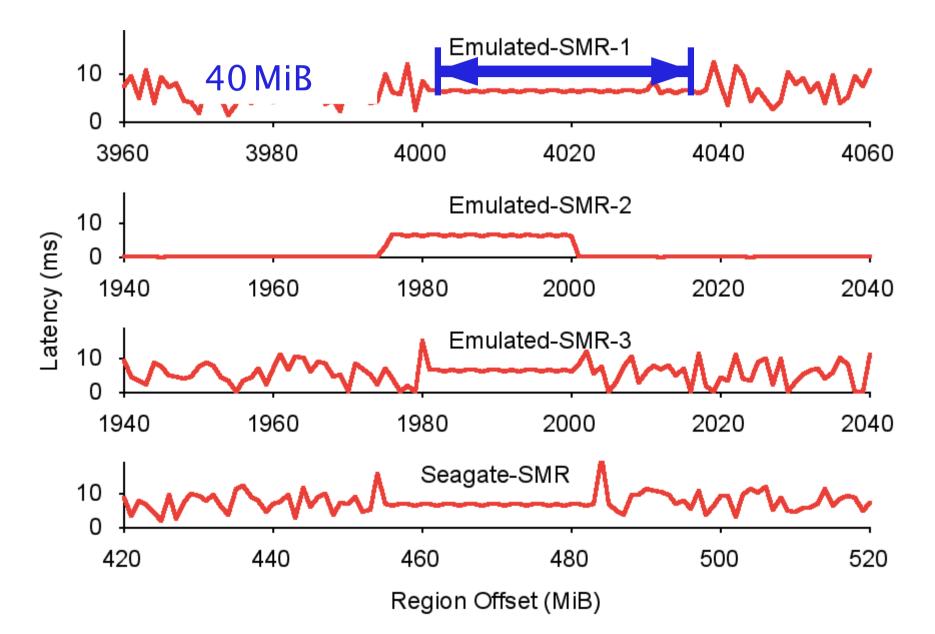
# Test 3: Discovering the Band Size

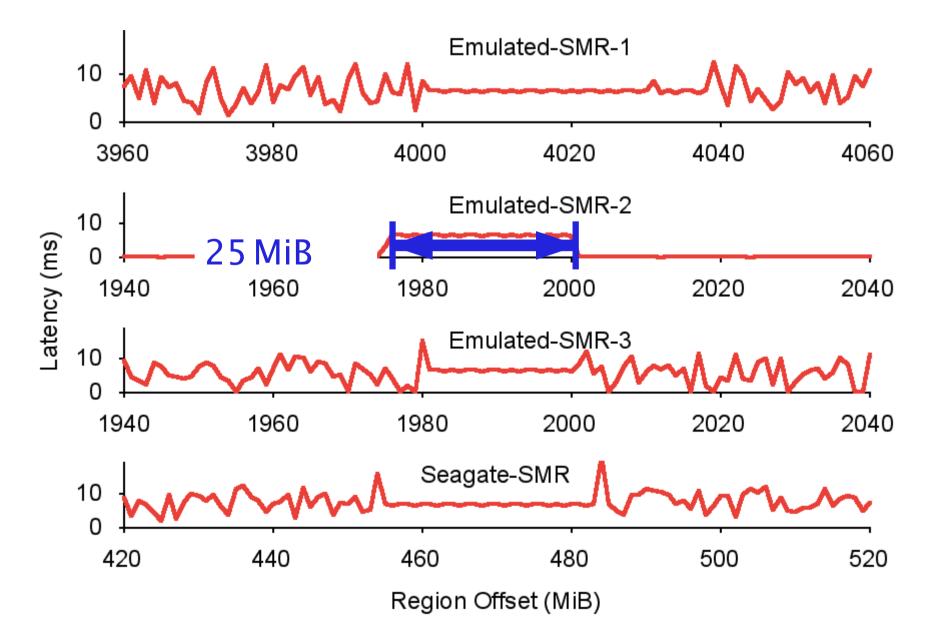
- Test relies on the fact that cleaning proceeds at a band granularity.
- Choose a small region (~1GiB) and write blocks in random order.
- Pause for a short (~3-5s) period, letting the cleaner to clean a few bands.
- Sequentially read the blocks in the region.
- Most latencies will be random a streak of flat latencies will identify a band.

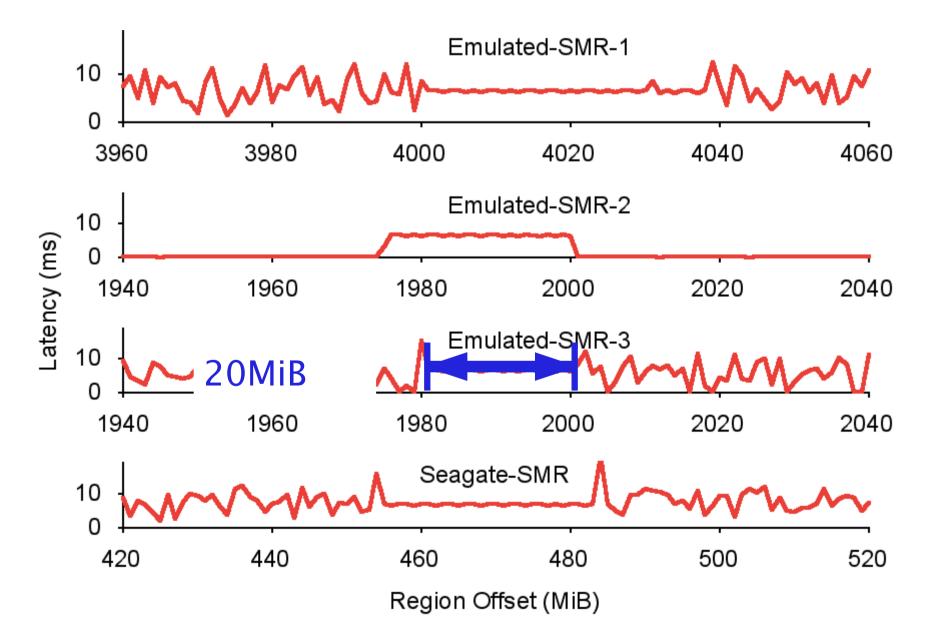
#### Band Size Detection

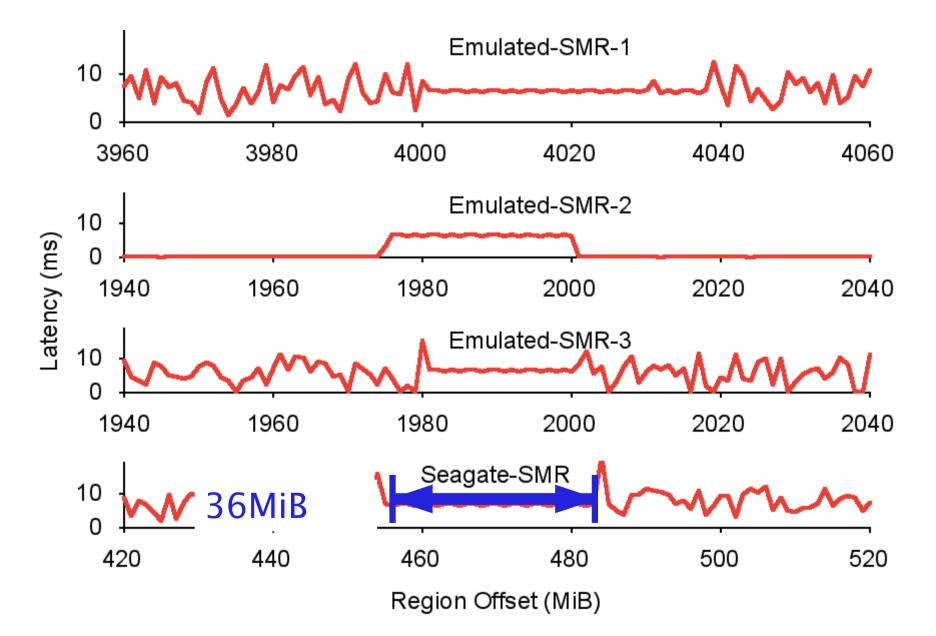












## Conclusion

- Drive-Managed SMR drives have different performance characteristics.
- Using them efficiently will require changes to software stack.
- Skylight aims to guide these changes.
- We aim for generality, more work may be needed.
- Tests, STL source code, video clips are available at http://sssl.ccs.neu.edu/skylight