COMPUTE + MEMORY S + STORAGE SUMMIT

Architectures, Solutions, and Community VIRTUAL EVENT, APRIL 11-12, 2023

Improving Storage
Systems for Simulation
Science with
Computational Storage

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Background: HPC Scientific Simulation Systems

Trinity – circa 2016

- Haswell and KNL
- 20,000 Nodes
- Few Million Cores
- 2 PByte DRAM
- 4 PByte NAND Burst Buffer ~ 4 TByte/sec
- 100 Pbyte Scratch PMR Disk File system ~1.2 TByte/sec
- 60PByte/year Sitewide Campaign Store ~ 50 GByte/sec
- 60 PByte Sitewide Parallel Tape Archive ~ 3 Gbyte/sec



I know its not Tier1 sized but at LANL its for one job for several years. 10 PB files and 200 PB Campaigns For a single user/small user team

Topics: Crawl, Walk, Run - with much help from our partners!

- ABOF 1.0 (Eideticom, Aeon, Nvidia, SK hynix)
 - Format agnostic operations (compression, erasure, encoding)
- DeltaFS->Ordered KV-CSD (CMU and SK hynix)
 - Format aware, record-oriented applications with a single-dimension, easily shard-able indexing
- ABOF 2.0 plans (Eideticom, Aeon, Nvidia, SK hynix, others?)
 - Format aware, column-oriented applications, multi-dimension, difficult to shard indexing
- Data-analysis in disk tier (Seagate)



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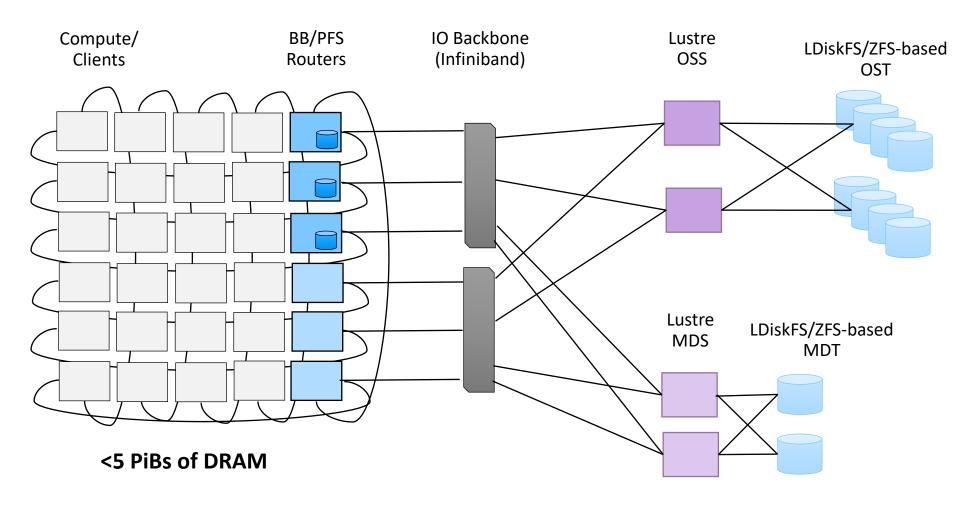
ABOF 1.0

- Format agnostic operations: compression, erasure, encoding
 - Eideticom, Aeon, Nvidia, SK hynix

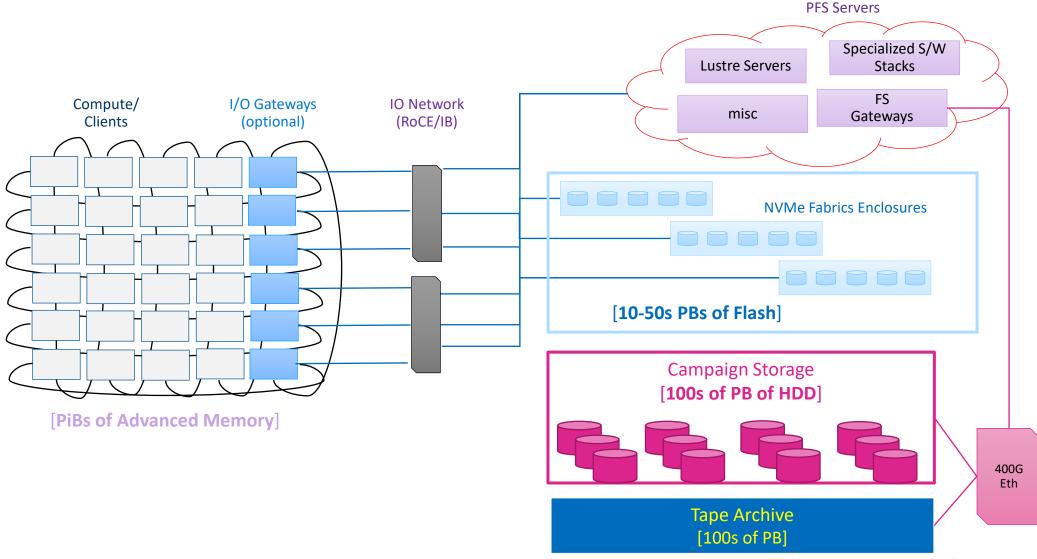
Traditional HPC Storage

<10 of PiBs of Flash

~100 of PiBs of HDD



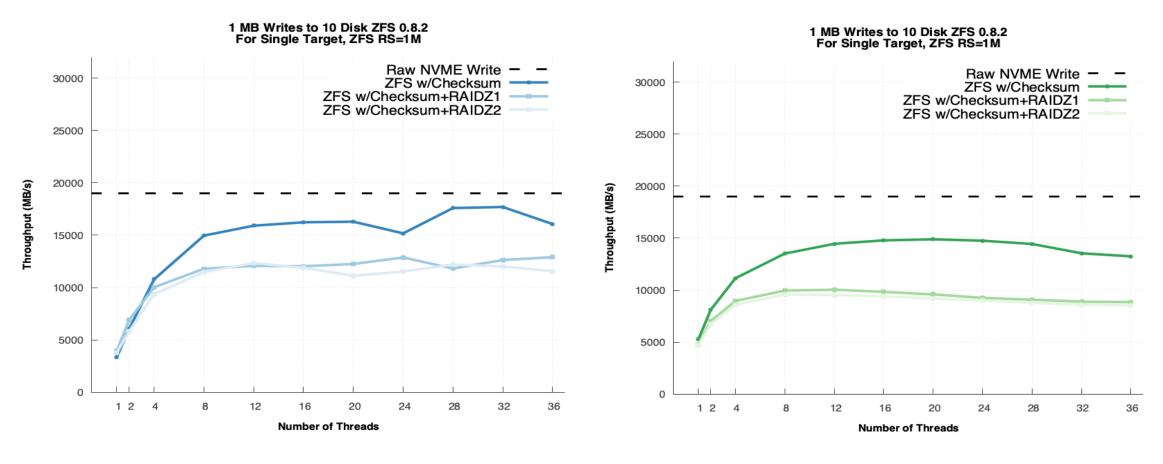
Redesign Opportunity



All Flash File Systems

- Require high performing storage server endpoints
 - Otherwise disaggregated isn't as important cost wise
- Current generation server memory bandwidth limitations observed relatively quickly
- With a budget, buying BW often doesn't result in high capacity
 - Compression is important
 - Compressing simulation data is hard!

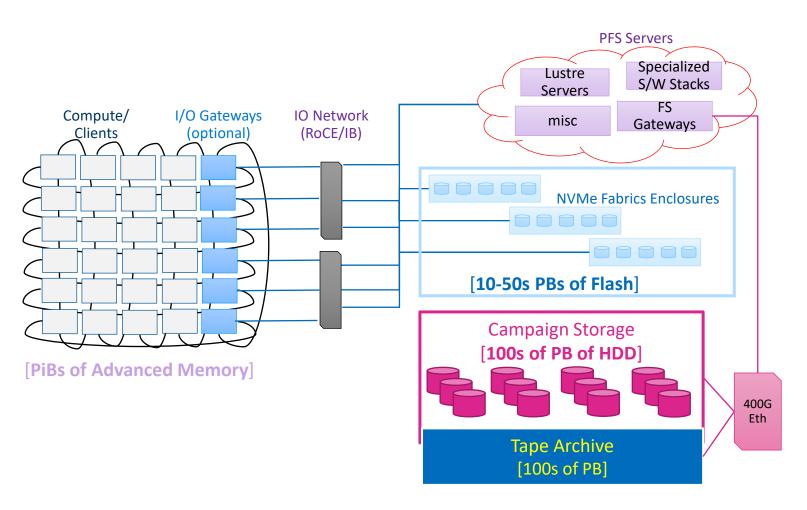
Why Offload? ZFS Checksums, Erasure, Compressive Server memory bandwidth is problematic and expensive

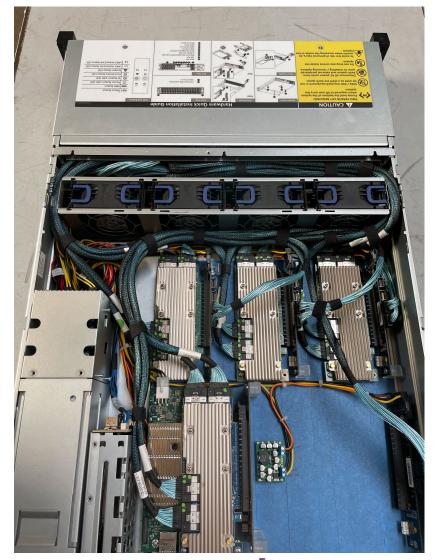


Intel Platinum (Dual Socket)

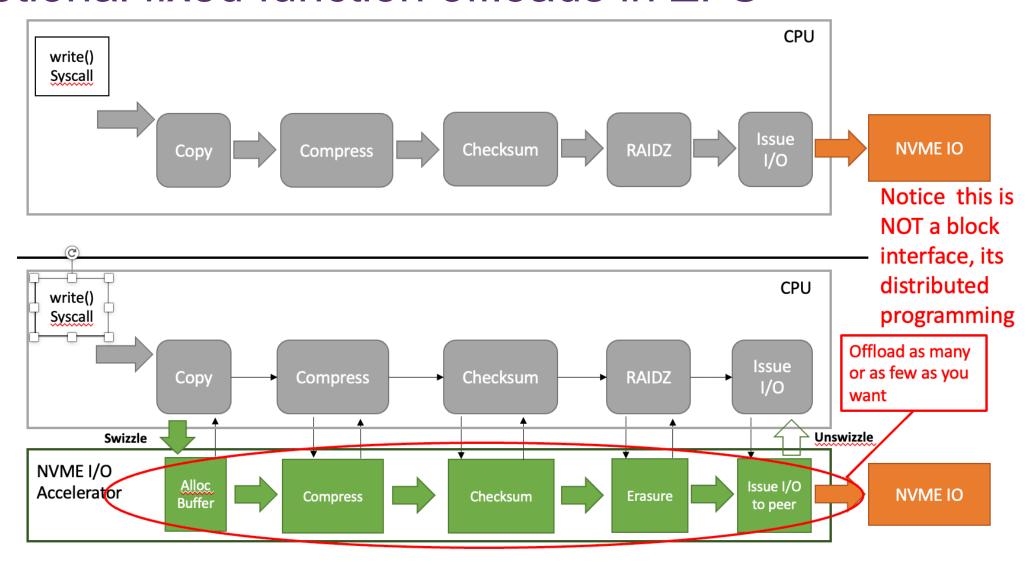
AMD EPYC (2nd Gen)

How to Consume: File System Services Offload





Notional fixed function offloads in ZFS





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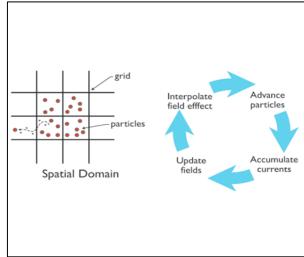


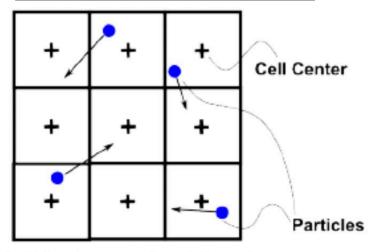
DeltaFS->Ordered KV-CSD

- Format aware, record-oriented applications with a single-dimension, easily shard-able indexing
 - CMU, SK hynix

Vector Particle in Cell (VPIC) our Record Based/Single Dimensional Index Application

- Particle-in-cell MPI code (scales to ~100K processes)
 - Fixed mesh range assigned to each process
 - Record: 32 64 Byte particles (id, cell id, energy, ...)
 - Particles move frequently between processes
 - Million particles per node (Trillions of particles in target simulation)
 - Interesting particles identified at simulation end (say 1000 interesting particles)





Grid for Computational Domain

Emerging Trends: Analysis Increasingly Selective

 Analysis used to require seeing all data records

 Today: queries tend only to hit a small subset of data

Problem: how to retrieve just interesting rows?

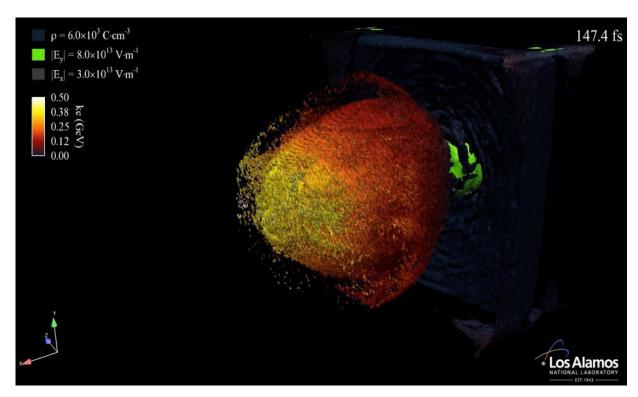


Image from LANL VPIC simulation done by L. Yin, et al at SC10

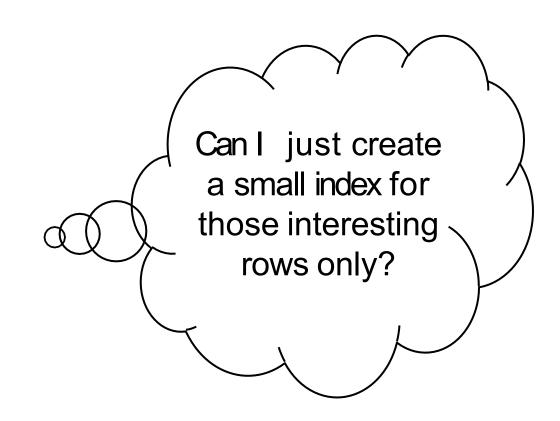
Example: SELECT X, Y, Z FROM particles WHERE E >= 1.5

Less than 0.1% needs to be read from storage

Reading Back Just Interesting Data is Non-Trivial

- Data known to be interesting only at simulation end
- Indexing only works when all rows are indexed at all timesteps
- Compute node resources are limited

Sorting only helps one query

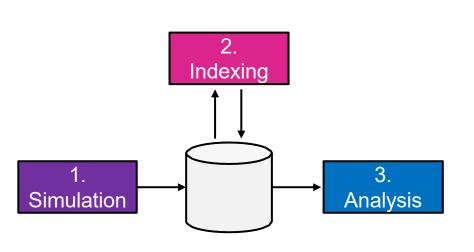


Existing Solutions Fall Short in Different Ways

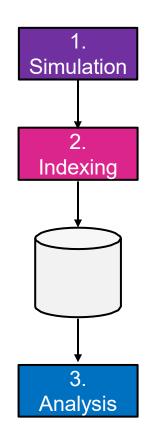
Post-processing

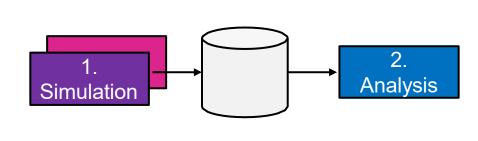
In-transit processing

In-situ processing



Excessive data movement





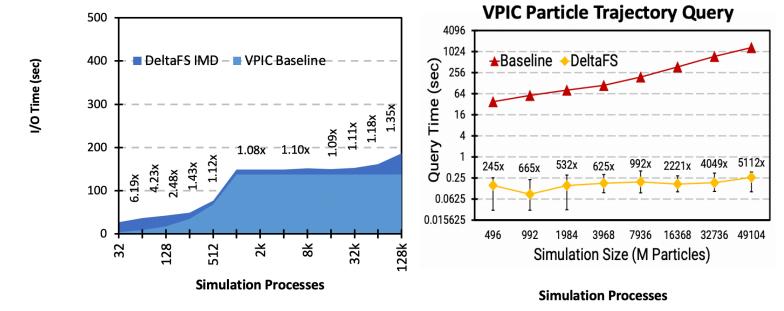
May only produce indexes on 1 or few columns

Requires additional compute nodes than the job Does not work for larger jobs

DeltaFS - Near-device Indexing and Analytics

- Requirements
 - Simulations run under intense memory pressure (app may use 90%)
- Computational Storage Benefits/Opportunities
 - Speedups for post-hoc analysis (1000x speedup demonstrated)
 - Less reliance on massive compute tier as a large merge sort space

Get efficiency and lower time to solution (1000X)



(papers at PDSW 15, PDSW 17, SC19 (Best Student Paper)

HPC-Driven KV Storage API

- Data insertion:
 - Bulk KV put operations
- Reads:
 - Range queries
 - Secondary indexes
 - Histogram construction
- Management:
 - Compaction control
 - Per key space data export



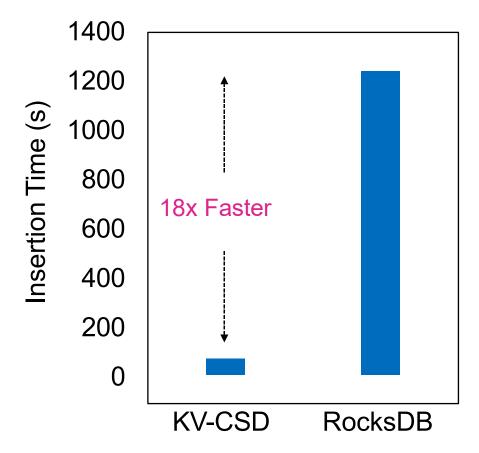
LANL is collaborating with industry for accelerated KV storage that speeds up scientific discovery

SK hynix Ordered KV-CSD Prototype Revealed at FMS '22

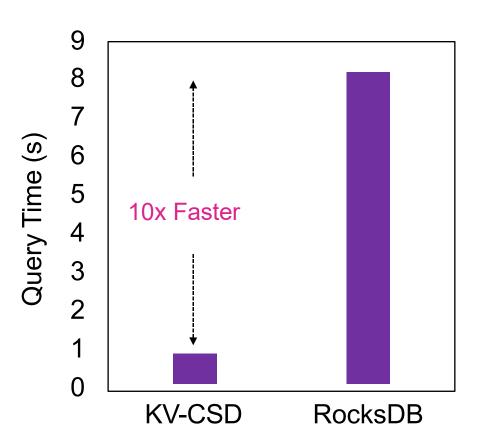
- Fully offloaded ordered key value with point and range query capability (put, get, mput, mget, etc.)
- Extensions control of compaction, and more
- Competitive performance



Preliminary Results: SK KV-CSD vs RocksDB



Data Insertion: Up to 18x faster



Queries: Up to 10x faster



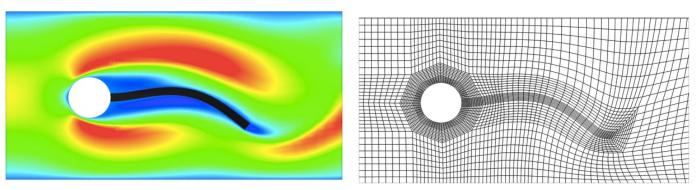
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Next Steps

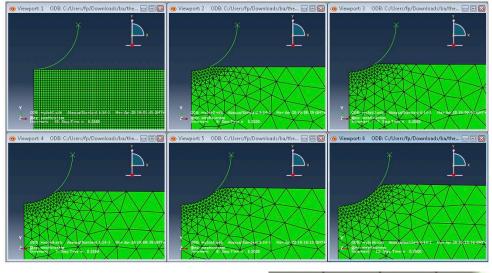
- Format aware, column-oriented applications, multi-dimension, difficult to shard indexing
 - Eideticom, Aeon, Nvidia, SK hynix, Seagate, others?)

What's a Grid Method and an Adaptive Mesh Refinement (AMR)?

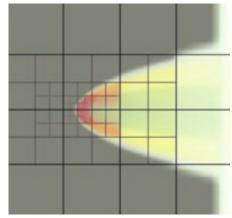


ALE – Advanced Lagrangian Eulerian
http://web.cs.ucdavis.edu/~ma/VolVis/amr_mesh.jpg

- Lagrangian (mesh deforms)
- Eulerian (mesh doesn't deform)
- AMR mesh adapts (refines where the action is)
- Why? to fit a problem that is way to big for your RAM
- AMR eliminates compression, copy on write, other low hanging fruit

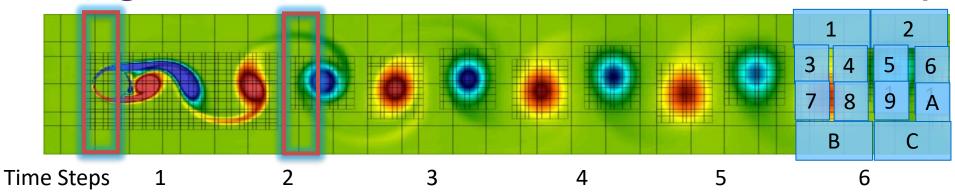






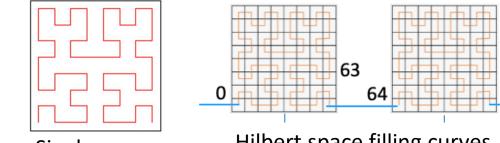
Eulerian AMR

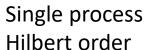
Indexing Multi-Dimensional Unstructured Adaptive Meshes

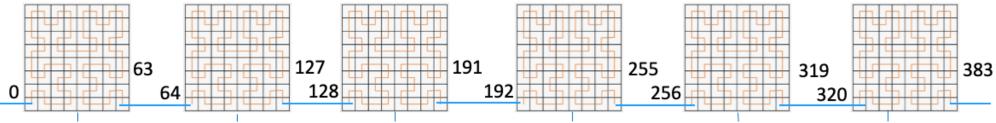


Processes have roughly same number of cells for comp/mem balance but must shuffle cells for AMR

- Time is explicit (a "file" for every time step) and that "file" contains all the state (for restart) (think 1 PB)
- Inside each mesh cell there is 10-100 state variables (64float) (temp, pressure, energy, momentum, ...)
- 2D and often 3D the other dimensions but how do you specify the geometric dimensions?

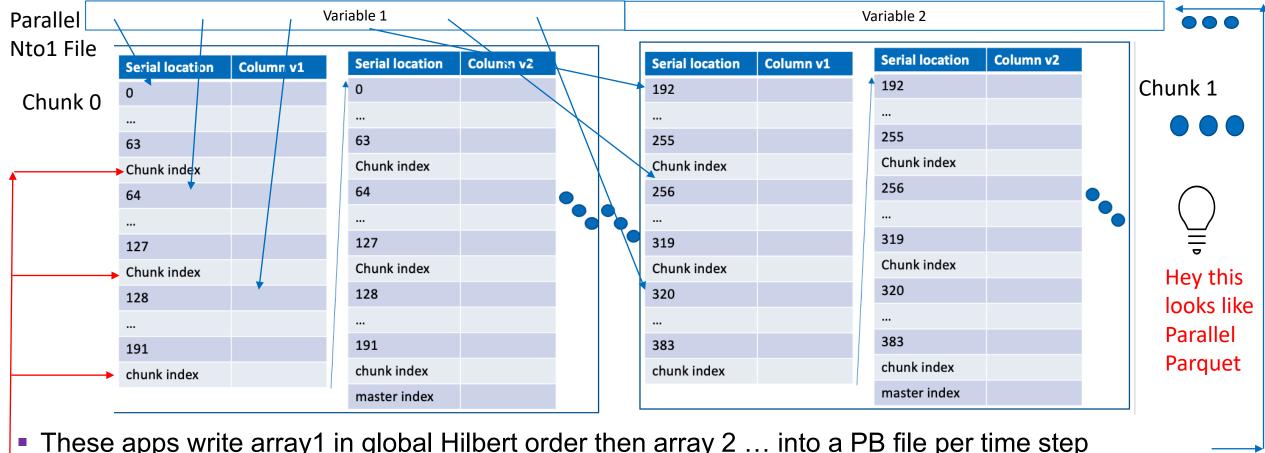






Hilbert space filling curves, we will call this the global Hilbert order, it serializes a value in the cells into distributed array. So we really have 10-100 distributed arrays in Hilbert order ©

Find the outer edge of the eddy's (light blue and yellow) <1/100th of the total data, usually less. Can light weight indexing yield 1000X less data and can it be done very near the storage device to save transmission? Adding index and offloading columnar analytics into Computational Storage, how would it work?

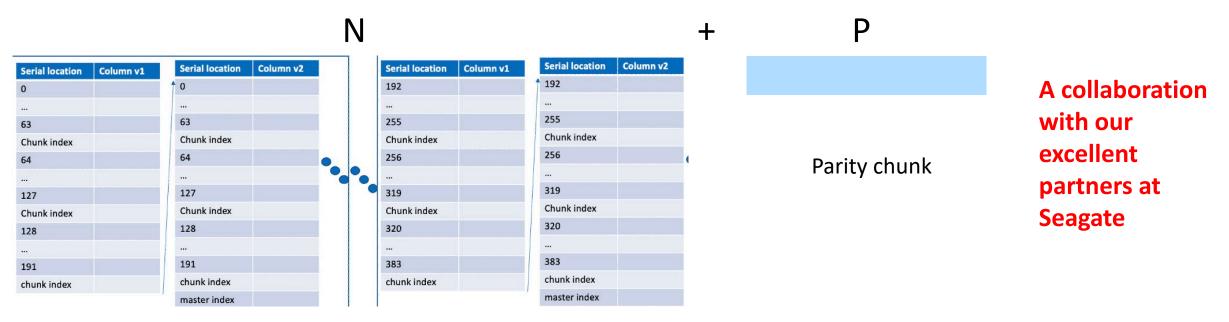


- These apps write array1 in global Hilbert order then array 2 ... into a PB file per time step
- Adding light weight indexes for every chunk of every variable is doable
- Use standard analytics with things like Duckdb or Apache Drill and have the power of SQL and joins on columns and the simple indexes do massive reduction in parallel

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Can this Hilbert Inspired Chunked Parquet Concept Extend to On-Disk Processing, Even with Erasure?

Parquet ZFS File with Erasure and On-Kinetic Disk Analytics in parallel



 Use standard analytics with things like Duckdb or Apache Drill and have the power of SQL and joins on columns and the simple indexes do massive reduction in parallel

Prototype Implementation

Multi-level Erasure Coding

Intel ISA-L

C2 Gateway

■LibZDB (LANL's stripped-down version of ZDB) for mapping ZFS

C2 OSD C2 OSD **ZFS RAID ZFS RAID**

C2 OSD **ZFS RAID**

filenames to disk LBAs

Drive exposed as an NVMeOF block device to ZFS

Analytics Path Data path Analytics Path Data path LibZDB ZFS LibZDB ZFS /dev/sda Kinetic Block custom client /dev/nvme1n1 libkinetic **Device (Kernel)** NVMeOF KV Kinetic Kinetic **NVMe Target HDD HDD** KV

Custom C++ code for indrive analytics

Current Prototype





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