Addressing Shingled Magnetic Recording drives with Linear Tape File System

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SNIA



Host Managed SMR

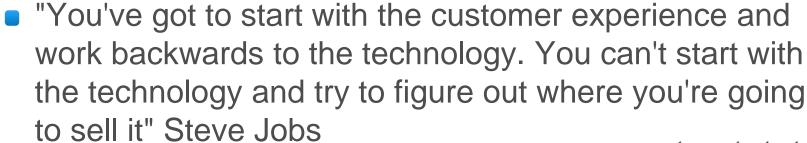
Agenda

- 1 Host Managed SMR Philosphy
- 2 Problem Statement Opportunity
- 3 Approaches
- 4 Hardware: Archive / Cloud drive
- 5 Software: Filesystem
- 6 Development
- 7 Validation: Testing & Performance
- 8 Lesson's Learned / Feedback



Host Managed SMR Philosophy

- Shingle Magnetic Recording is a disruptive technology – Innovation that results in variable performance but adds new value proposition. Fringe customers may benefit and over time performance may increase.
- Identify markets ripe to benefit from SMR
 - Able to modify their storage stack
 - Write 100% sequentially (tape like)
 - Benefit from lower \$/GB



Most demanding use

Time

Host Managed SMR

Problem Statement / Opportunity

Background:

- Archive / Cold storage is a new market.
 - Write once and forget. (read seldom, if ever).
 - Primary care about: TCO
- Shingle Magnetic Recording is a new technology
 - Sequential write (only)
 - Path to lowest \$/GB if host can abide by one simple rule
- Industry has discussed two approaches for SMR
 - Drive does data management similar to FTL
 - Host does data management via a SMR file system

Problem Statement

 Solve the new market opportunity, Archive / cold storage with new recording technology that provides lowest \$/GB





Host Managed SMR Approaches

- Drive managed SMR
 - Similar to Flash Translation layers, data management is complicated.
 - Metadata, garbage collection, over provisioning, write amplification, electronics resources, variable performance, validation.
- Host Managed SMR
 - Similar to Flash File Systems, data management is complicated but there may be opportunity to leverage mature and popular file systems that write sequentially.
 - Drive will not be burdened by the resource cost of a drive managed system



Host Managed SMR

Hardware: Archive cloud drive

- Host Managed SMR, Single Zone, Host writes 100% sequentially – Simplest SMR implementation
- Intended for "Big Data", Cold/Archive, Tape Replacement
- Use Case:
 - Write once, read seldom
 - Garbage Collect from Device to Device
- We built it
 - Proved AD gain, Build & test process
 - Proved Device Architecture, Reliability & Error handling



Software

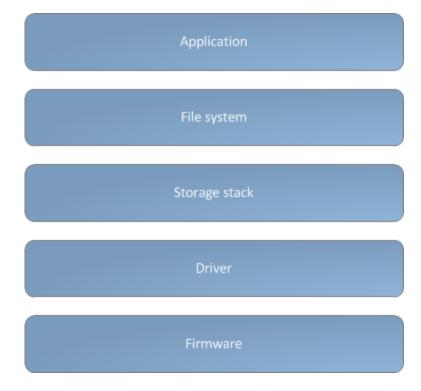


Host I/O

- Sequential write only.
- No overwrite allowed.
- Garbage collection.
- Handle out of bounds access failure.



Abstraction layer





File system

- Industry usage/support.
- Open source/specification.
- User vs kernel space
 - Minimize risk
 - Ease of development/debug



LogFS

"LogFS is a Linux log-structured and scalable flash file system, intended for use on large devices of flash memory. It is written by Jörn Engel and in part sponsored by the CE Linux Forum.

LogFS is included in the mainline Linux kernel and was introduced in version 2.6.34, released on May 16, 2010."

- Wikipedia (http://en.wikipedia.org/wiki/LogFS)



LTFS



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Abstract—While there are many financial and practical reasons drive. A tape library can host dozens or even hundreds of to prefer tape storage over disk for various applications, the drives operating in parallel. A case-study comparison [3] found difficultly of using tape in a general way is a major inhibitor to the

its wider usage. We present a I a new generation of tape has to tape using standard, famil The Linear Tape File System flexible, portable, and intuitive sharable media, such as a USI

I. Mo

In today's digital world, in of their data, including their form. The broadcast and me changing to all-digital, file-bar a major transition. In what is Media Transformation, tradit being replaced with file-base

This transformation helps storage capacity. More impofor storage technologies. A de Motion Pictures Art and Scimission to keep and presen hundred years. It states that software exists that can reas bility to digital assets. When information technology com

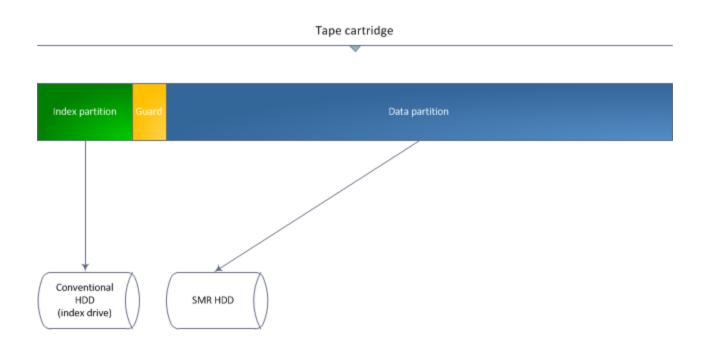
File Mark File Mark File Mark File Mark Index VOL1 LTFS Index Partition Label Label File Mark File Mark File Mark Data VOL1 LTFS Index Index Index Partition Label Label File Mark File Mark File Mark File Mark

year old film can still be projected and scanned today. In update-in-place is not possible. Hence tape in used as an



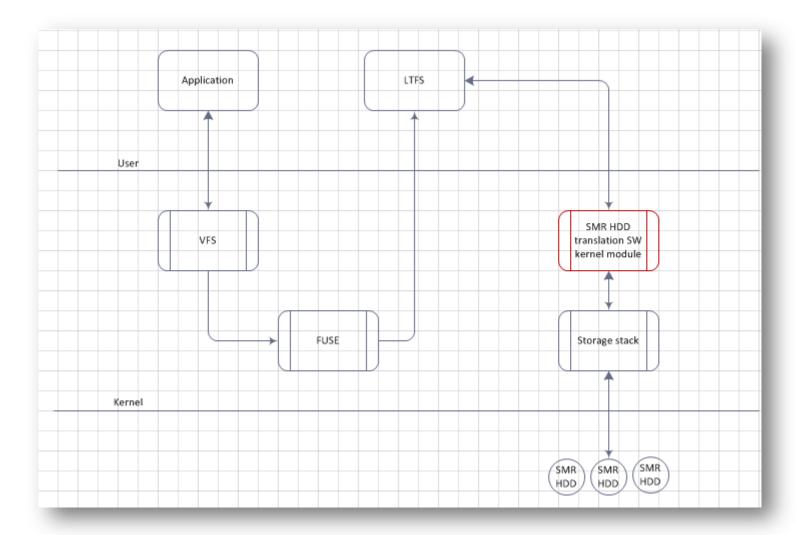


Partition



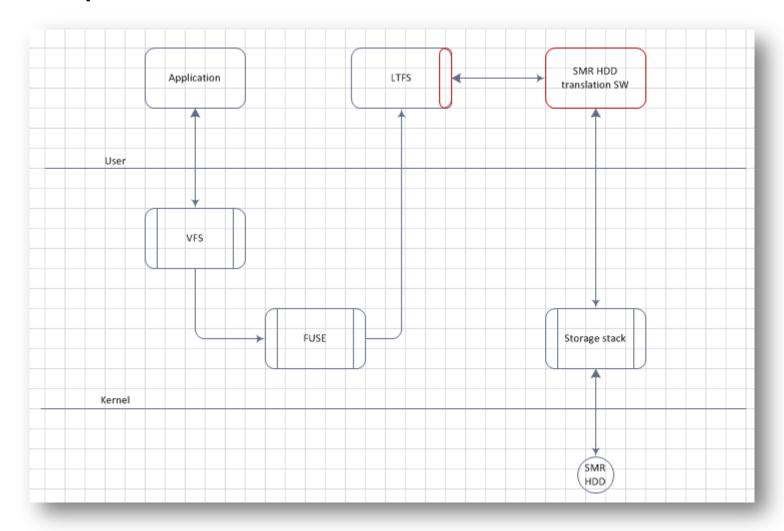


Kernel module



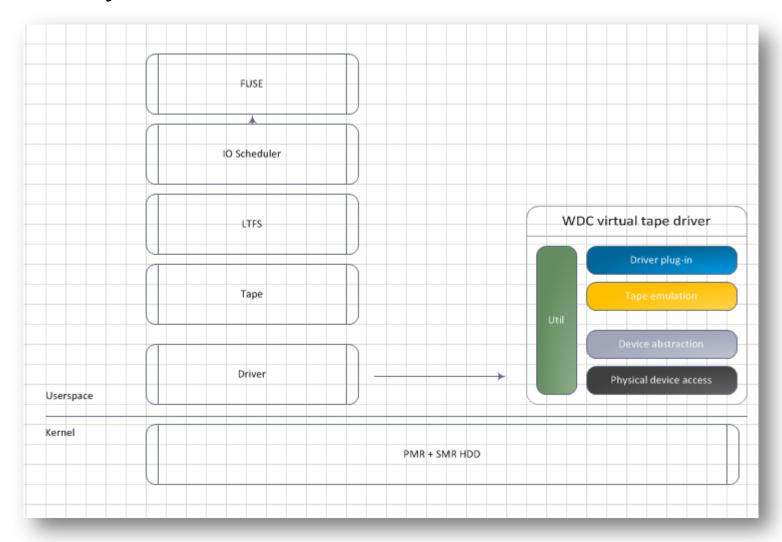


User space





Code layout





Development



LTFS 1.2.5 with Ubuntu 10.04.4 LTS Lucid Lynx

Coding

- No changes to LTFS source
- Backend tape driver
- ~5000 LOC in 25 files
- 2 month POC
- Work in progress



Validation: Testing & Development

Validation at each level

- Directory
- File
- Device



POC setup







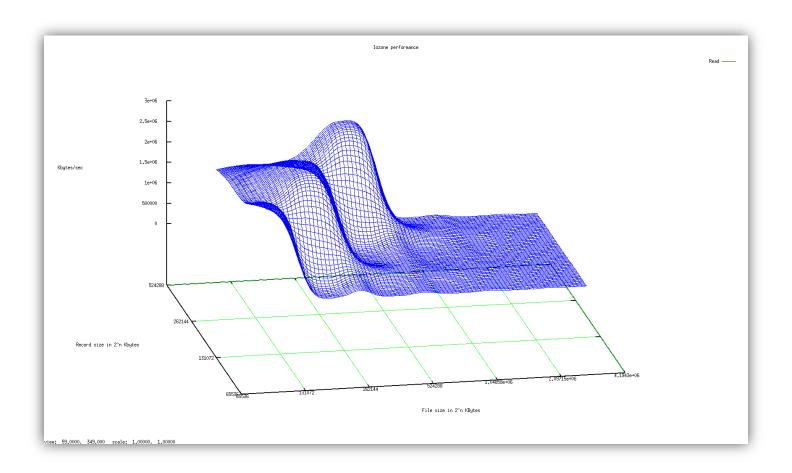


Parameters

- File size: 500MB to 4GB
- Record size: 100MB to 500MB
- O_SYNC
- Purge processor cache before each file operation
- Unmount LTFS between each test
- Intel Corporation 5 Series/3400 Series Chipset 6
 port SATA AHCI Controller (rev 06)

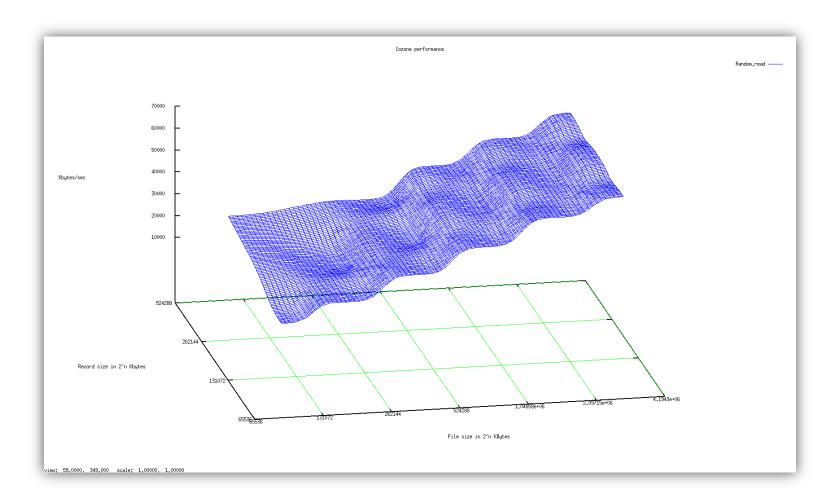


Performance – Sequential Read



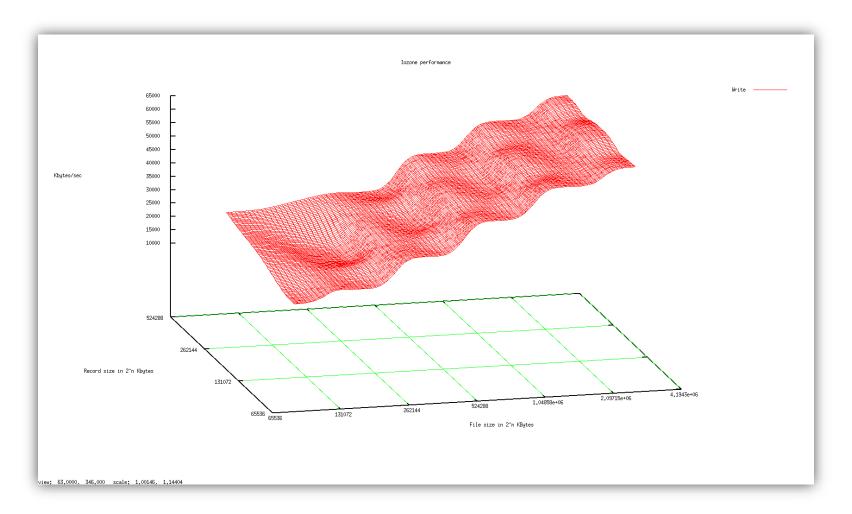


Performance – Sequential Write



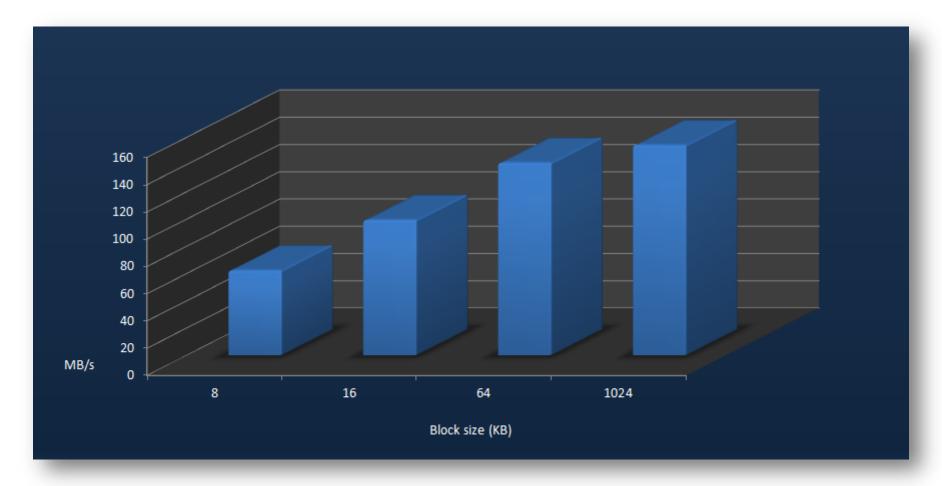


Performance – Sequential Write



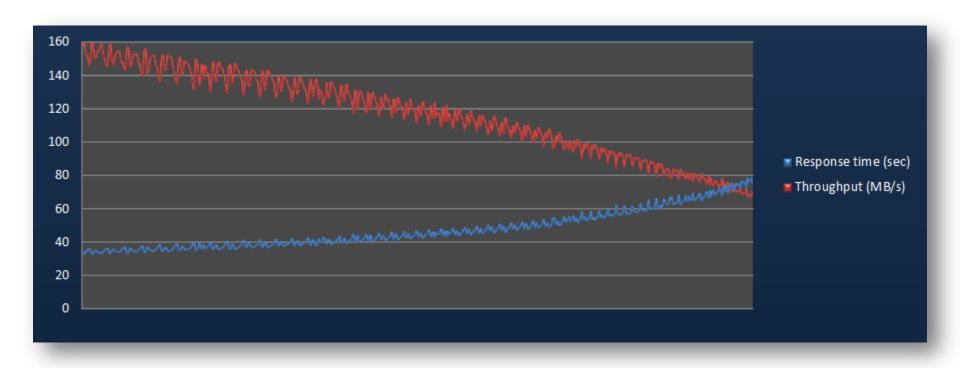


Throughput f(block size)





Throughput f(LBA)





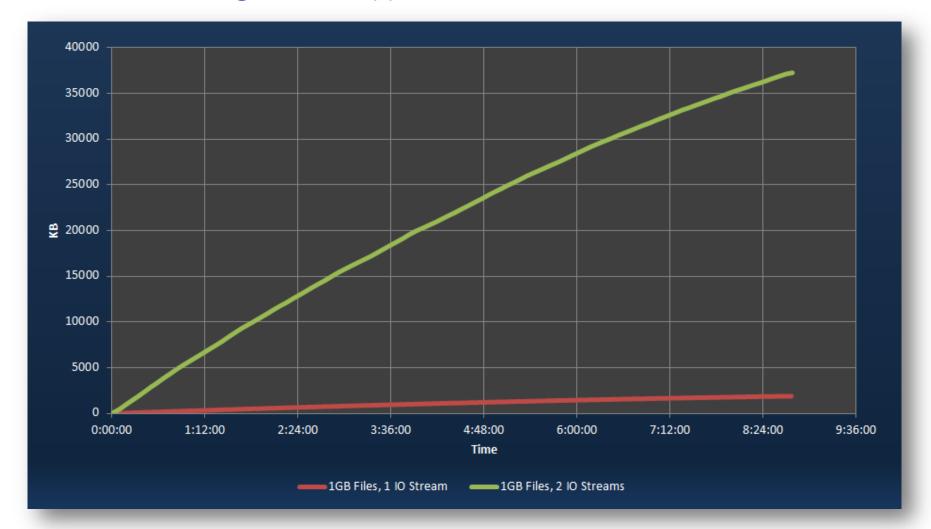
Write sequential: ~116MB/s



Metadata

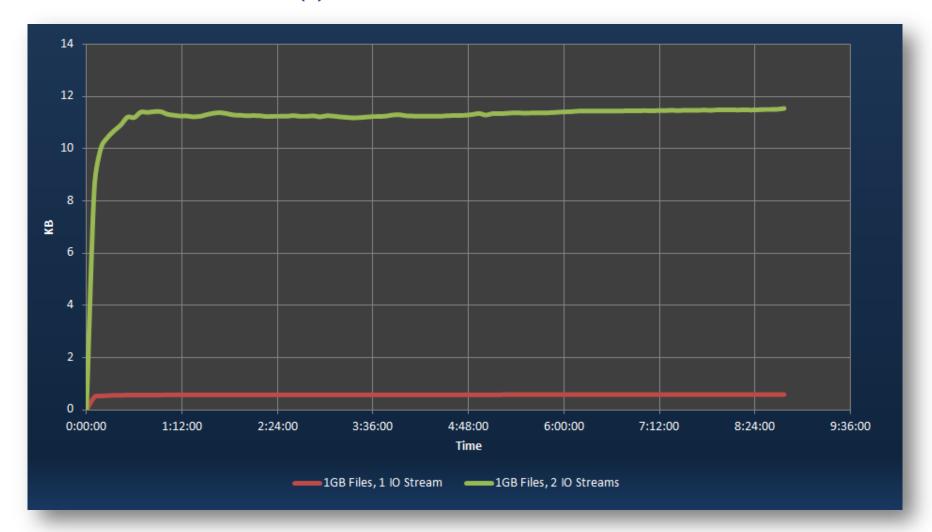


Index size growth f(t)



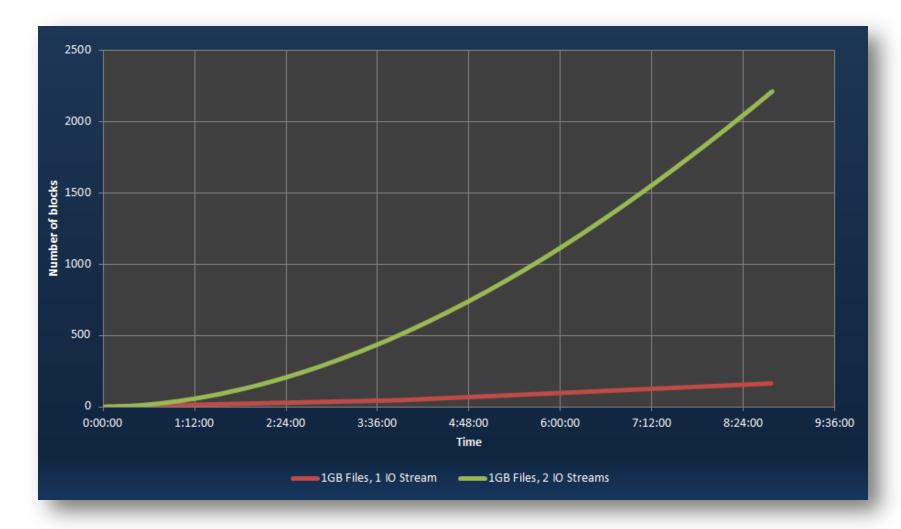


Index size/file f(t)





Index block allocation (1MB) f(t)



System information

- Dell PowerEdge T110
- 2GB 1333 MHz DRAM

Intel Xeon E3-1220 CPU

LSI SAS9207-8e HBA



CPU and memory usage f(t)





What have we learned?



- 1.LTFS interoperability
- 2. Performance
- 3. Testing



Host managed SMR

Feedback

- Customers "dig" it easy to use
- Writes are 100% sequential fastest way to store data on HDD
- Read access at full HDD speeds (random and sequential)
- Smallest overprovisioned solution lowest \$/GB
- Simple design lower cost higher reliability

